

Axiomatizing AECs and Applications

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Overview

- 1 Main results
- 2 Abstract elementary classes (AECs)
- 3 Proof idea
- 4 Other results
- 5 Open questions

Main results

Theorem (Shelah)

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- 1 $K = PC(T, \Gamma, L(\mathbf{K}))$;
- 2 $|T| \leq \lambda$;
- 3 $|\Gamma| \leq 2^\lambda$.

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Theorem (4.1, 4.10)

Let \mathbf{K} be an AEC and $\lambda = \text{LS}(\mathbf{K})$. Then there is χ depending on \mathbf{K} such that $\lambda \leq \chi \leq 2^\lambda$ and K is $PC_{\chi, \chi}$.

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Theorem (4.1,4.10)

Let \mathbf{K} be an AEC and $\lambda = \text{LS}(\mathbf{K})$. Then there is χ depending on \mathbf{K} such that $\lambda \leq \chi \leq 2^\lambda$ and K is $PC_{\chi, \chi}$.

Moreover, under $2^\lambda < 2^{\lambda^+}$, if \mathbf{K} is categorical in λ, λ^+ and stable in λ , then \mathbf{K} is $PC_{\lambda, \lambda}$.

Fact (Shelah)

Let \mathbf{K} be an AEC, $\theta \geq \text{LS}(\mathbf{K})$. Suppose the following hold:

- 1 $K, K^{<}$ are both $PC_{\theta, \theta}$;
- 2 \mathbf{K} is categorical in both θ and θ^+ ;
- 3 $\delta(\theta, 1) = \theta^+$. (Threshold cardinal for an infinite decreasing chain to exist in a $PC_{\theta, 1}$ -class.)

Then $K_{\theta^{++}} \neq \emptyset$.

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Corollary (4.8)

(1) is true for $\theta \geq \chi$. Moreover, under $2^\lambda < 2^{\lambda^+}$ and stability in λ , (2) already implies (1) for $\theta = \lambda$.

Abstract elementary classes (AECs)

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Definition

Let L be a finitary language. An abstract elementary class $\mathbf{K} = \langle K, \leq_{\mathbf{K}} \rangle$ in $L = L(\mathbf{K})$ satisfies the following axioms:

- 1 K is a class of L -structures and $\leq_{\mathbf{K}}$ is a partial order on $K \times K$.
- 2 For $M_1, M_2 \in K$, $M_1 \leq_{\mathbf{K}} M_2$ implies $M_1 \subseteq M_2$ (as L -substructures).

Abstract elementary classes (AECs)

Definition (Continued)

3 Isomorphism axioms:

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- a If $M \in K$, N is an L -structure, $M \cong N$, then $N \in K$.
- b Let $M_1, M_2, N_1, N_2 \in K$. If $f : M_1 \cong M_2$, $g : N_1 \cong N_2$, $g \supseteq f$ and $M_1 \leq_K N_1$, then $M_2 \leq_K N_2$.

$$\begin{array}{ccc} N_1 & \xrightarrow{g} & N_2 \\ \leq_K \uparrow & & \leq_K \uparrow \cdots \\ M_1 & \xrightarrow{f} & M_2 \end{array}$$

Abstract elementary classes (AECs)

Definition (Continued)

- ④ Coherence: Let $M_1, M_2, M_3 \in K$. If $M_1 \leq_K M_3$, $M_2 \leq_K M_3$ and $M_1 \subseteq M_2$, then $M_1 \leq_K M_2$.

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- 4 Coherence: Let $M_1, M_2, M_3 \in K$. If $M_1 \leq_K M_3$, $M_2 \leq_K M_3$ and $M_1 \subseteq M_2$, then $M_1 \leq_K M_2$.
- 5 Löwenheim-Skolem axiom: There exists an infinite cardinal $\lambda \geq |L(\mathbf{K})|$ such that: for any $M \in K$, $A \subseteq |M|$, there is some $N \in K$ with $A \subseteq |N|$, $N \leq_K M$ and $\|N\| \leq \lambda + |A|$. We call the minimum such λ the Löwenheim-Skolem number $LS(\mathbf{K})$.

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- 6 Chain axioms: Let α be an ordinal and $\langle M_i : i < \alpha \rangle \subseteq K$ such that for $i < j < \alpha$, $M_i \leq_K M_j$.
 - 1 Then $M = \bigcup_{i < \alpha} M_i$ is in K and for all $i < \alpha$, $M_i \leq_K M$.
 - 2 Let $N \in K$. If in addition for all $i < \alpha$, $M_i \leq_K N$, then $M \leq_K N$.

Abstract elementary classes (AECs)

Definition

Let \mathbf{K} be an AEC and $\lambda \geq \text{LS}(\mathbf{K})$.

$$I(\lambda, \mathbf{K}) = |\{M/\cong : M \in K_\lambda\}|$$

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$$I_2(\lambda, \mathbf{K}) = |\{(M, N)/\cong : M \leq_{\mathbf{K}} N \text{ both in } K_\lambda\}|$$

where $(M_1, N_1) \cong (M_2, N_2)$ iff $M_1 \leq_{\mathbf{K}} N_1$, $M_2 \leq_{\mathbf{K}} N_2$ and there is $g : N_1 \cong N_2$ such that $g \upharpoonright M_1 : M_1 \cong M_2$.

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Fact

$$I_2(\lambda, \mathbf{K}) \leq 2^\lambda.$$

In the main results, we had $\chi = \lambda + I_2(\lambda, \mathbf{K})$, and bypassed I_2 under more assumptions.

Proof idea

Definition

Let I be an index set. A directed system $\langle M_i : i \in I \rangle \subseteq K$ indexed by I satisfies the following: for any $i, j \in I$, there is $k \in I$ such that $M_i \leq_K M_k$ and $M_j \leq_K M_k$.

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Given $M \in K$, we index the system by finite tuples of elements in M . Namely, $I = |M|^{<\omega}$ (ordered by inclusion).

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Given $M \in K$, we index the system by finite tuples of elements in M . Namely, $I = |M|^{<\omega}$ (ordered by inclusion). Conversely,

Fact

Let $\langle M_i : i \in I \rangle \subseteq K$ be a directed system. Then

- 1 $M = \bigcup_{i \in I} M_i \in K$;
- 2 For all $i \in I$, $M_i \leq_K M$;
- 3 Let $N \in K$. If in addition for all $i \in I$, $M_i \leq_K N$, then $M \leq_K N$.

Proof idea

Fact

Let \mathbf{K} and \mathbf{K}' be two AECs with $L(\mathbf{K}) = L(\mathbf{K}')$. If there exists a cardinal λ such that

- 1 $\lambda \geq \text{LS}(\mathbf{K}) + \text{LS}(\mathbf{K}')$;
- 2 $\mathbf{K}_\lambda = \mathbf{K}'_\lambda$ (both the models and the ordering),

then $\mathbf{K}_{\geq \lambda} = \mathbf{K}'_{\geq \lambda}$.

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then $\mathbf{K}_{\geq \lambda} = \mathbf{K}'_{\geq \lambda}$.

Hence, given an AEC \mathbf{K} , it suffices to encode the models and the ordering in $\mathbf{K}_{\text{LS}(\mathbf{K})}$.

Proof idea

Let M be an $L(\mathbf{K})$ -structure, $\lambda = \text{LS}(\mathbf{K})$.

- 1 Expand $L(\mathbf{K})$ by adding λ -many functions. They map a finite tuple to a \mathbf{K} -structure containing it:

$$a \in |M|^{<\omega} \mapsto \{f_i(a) : i < \lambda\} = |M_a| \text{ with } M_a \in K.$$

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- 2 Stipulate that $\langle M_a : a \in |M|^{<\omega} \rangle$ forms a directed system:
If $a \cup b \subseteq c$ then $M_a \leq_{\mathbf{K}} M_c$ and $M_b \leq_{\mathbf{K}} M_c$.

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If $a \cup b \subseteq c$ then $M_a \leq_{\mathbf{K}} M_c$ and $M_b \leq_{\mathbf{K}} M_c$.
- ▶ “ $\leq_{\mathbf{K}}$ ” is by listing all the isomorphism types of pairs $N_0 \leq_{\mathbf{K}} N_1$ in K_λ , there are $I_2(\lambda, \mathbf{K})$ -many.

Proof idea

We have axiomatized \mathbf{K} in an $L'_{\chi^+, \omega}$ -sentence σ where L' contains extra functions than $L(\mathbf{K})$ and $\chi = \lambda + I_2(\lambda, \mathbf{K})$. To convert this to a PC-class, we use the following fact:

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Chang's presentation theorem

Let θ be an infinite cardinal, L be a language of size $\leq \theta$, T be an $L_{\theta^+, \omega}$ -theory contained in a fragment of size $\leq \theta$, then the models of T is a $PC_{\theta, \theta}$ -class.

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Taking $T = \{\sigma\}$, we obtain \mathbf{K} as a $PC_{\chi, \chi}$ -class.

Proof idea

Theorem

Let \mathbf{K} be an AEC and $\lambda = \text{LS}(\mathbf{K})$. Under $2^\lambda < 2^{\lambda^+}$, if \mathbf{K} is categorical in λ, λ^+ and stable in λ , then \mathbf{K} is $PC_{\lambda, \lambda}$.

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By categoricity in λ , we avoid listing (individual) models indexed by $I(\lambda, \mathbf{K})$. The assumptions also imply that we can build limit models that are isomorphic over the same base:

$$\begin{array}{ccc} M_1 & \overset{\cong}{\dashrightarrow} & M_2 \\ & \swarrow (\lambda, \omega) & \nearrow (\lambda, \omega) \\ & M_0 & \end{array}$$

(In other words, if $I(\lambda, \mathbf{K}) = 1$ and $\leq_{\mathbf{K}}$ is replaced by “ (λ, ω) -limit”, then $I_2(\lambda, \mathbf{K}) = 1!$)

Proof idea

Instead of encoding $M_a \leq_{\mathbf{K}} M_b$, we use coherence axiom and encode $M_a \subseteq M_b$ and the existence of a common limit model over them.

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$$\begin{array}{ccc} M_b & & M_b \xrightarrow{(\lambda, \omega)} M^* \\ \uparrow \leq_{\mathbf{K}} & \text{is the same as} & \uparrow \subseteq \\ M_a & & M_a \end{array}$$

The diagram illustrates the equivalence between two ways of representing the relationship between models M_a and M_b . On the left, M_a is related to M_b via the relation $\leq_{\mathbf{K}}$, shown as a vertical arrow pointing up. On the right, M_a is related to M_b via the relation \subseteq , also shown as a vertical arrow pointing up. Additionally, there is a diagonal arrow from M_a to a common limit model M^* , labeled with (λ, ω) . A horizontal arrow from M_b to M^* is also labeled with (λ, ω) .

Other results

One can also axiomatize an AEC \mathbf{K} within the same language $L(\mathbf{K})$.

Fact (Shelah-Villaveces)

Let \mathbf{K} be an AEC, $L = L(\mathbf{K})$ and $\lambda = LS(\mathbf{K})$. Then \mathbf{K} can be axiomatized by a sentence in $L_{(2^{2^{\lambda^+}})+++}$.

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Theorem (3.7)

Let \mathbf{K} be an AEC, $L = L(\mathbf{K})$, $\lambda = \text{LS}(\mathbf{K})$ and $\chi = \lambda + I_2(\lambda, \mathbf{K})$. Then \mathbf{K} can be axiomatized by a sentence in $L_{\chi^+, \lambda^+}(\omega \cdot \omega)$ (game quantification of $\omega \cdot \omega$ steps).

Other results

Let $\mu \geq \aleph_0$. Define PC^μ as before except that the underlying languages L, L' are $(< \mu)$ -ary.

K is	K is axiomatizable in	K is
An AEC	$L_{\chi^+, \lambda^+}(\omega \cdot \omega)$	$PC_{\chi, \chi}$
A μ -AEC	$L_{\chi^+, \lambda^+}(\mu \cdot \mu)$	$PC_{\chi, \chi}^\mu$

Open questions

Question

Let \mathbf{K} be an AEC and $\lambda = \text{LS}(\mathbf{K})$.

- Under extra assumptions (categoricity, stability, etc), is it possible to bound $I_2(\lambda, \mathbf{K})$ strictly below 2^λ ?
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 - ▶ This might remove the game quantification in our axiomatization.
- Does the Hanf number exist for μ -AECs?
 - ▶ Hanf number is the threshold cardinal for arbitrarily large models. In AECs, the Hanf number is $\beth_{(2^\lambda)^+}$.