

EDGES AND TRIANGLES

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Joint work with Jacob Fox

OBSERVATION

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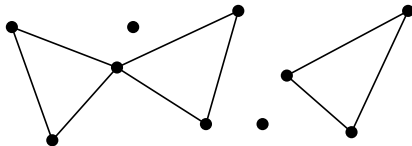
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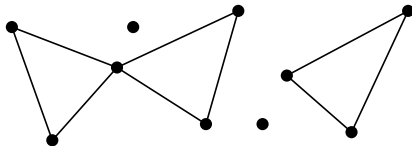


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QUESTION (ERDŐS-ROTHSCHILD)

What if the total number of edges must be at least $0.001n^2$?
Must some edge be in many triangles?

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For every ϵ , there is M such that every graph can be ϵ -approximated by an object of complexity bounded by M .

TRIANGLE REMOVAL LEMMA

For any ϵ , there is a δ such that every graph with $\leq \delta n^3$ triangles can be made triangle-free by deleting only ϵn^2 edges.

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Dependency between parameters:

- (From Regularity Lemma.) $\frac{1}{\delta}$ is tower of height power of $\frac{1}{\epsilon}$.
- (Fox.) $\frac{1}{\delta}$ is tower of height logarithmic in $\frac{1}{\epsilon}$.

LOWER BOUND FOR ERDŐS-ROTHSCHILD

OBSERVATION

Let c be a constant. Given cn^2 edges, each of which is in a triangle, there is always an edge which is in $\log^* n$ triangles.

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- Either case gives an edge in at least $\min\{\frac{3\delta n}{c}, \frac{c}{3\epsilon}\}$ triangles.
- Take $\frac{1}{\delta} = \sqrt{n}$, and $\frac{1}{\epsilon} = \text{power of } \log^* n$. □

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QUESTION (ERDŐS, 1987)

Given cn^2 edges, each of which is in a triangle, is there always some edge which is in at least n^ϵ triangles, for a constant ϵ ?

THEOREM (FOX, L.)

There are n -vertex graphs with

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edges, each of which is in a triangle, but with no edge in more than $n^{14/\log \log n}$ triangles.

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Remarks:

- Every edge is in under $n^{o(1)}$ triangles.
- The edge density approaches $\frac{1}{4}$ from below.
- Sharp transition: after edge density $\frac{1}{4}$, some edge is in a linear number of triangles.

THEOREM (HOEFFDING-AZUMA)

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$$\mathbb{P} [|X - \mathbb{E}[X]| > t] \leq 2e^{-\frac{t^2}{2L^2n}} .$$

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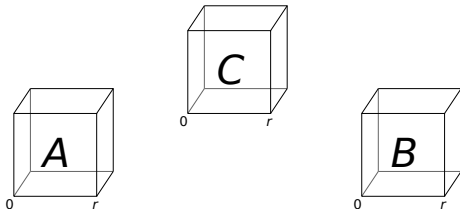
CLASSICAL RESULT

In even dimensions d , the Euclidean ball of radius r has

$$\text{Vol} \left(B_r^{(d)} \right) = \frac{\pi^{d/2} r^d}{(d/2)!} .$$

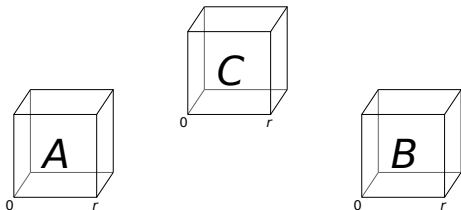
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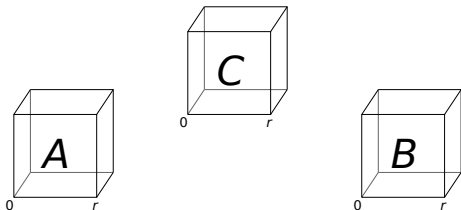
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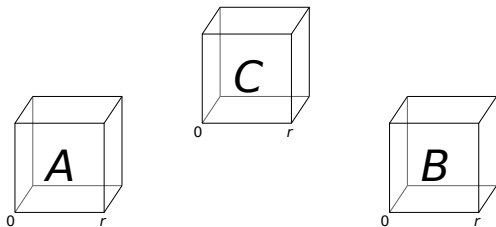
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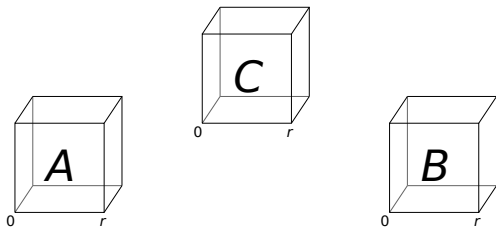


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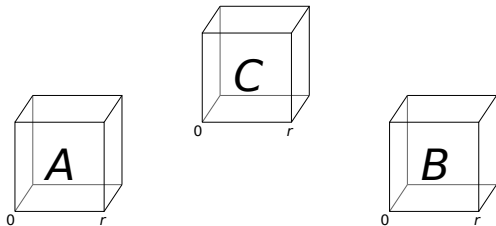
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- The edge density between A and B is the probability that two random points in the cube have squared-distance within $\mu \pm d$.



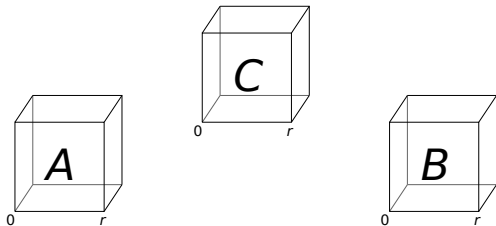
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- The typical deviation from μ is $r^2\sqrt{d} = r^{4.5} \ll d$, since $d = r^5$, so the A - B edge density approaches 1!



Every A–B edge is in a triangle:

- A – B endpoints have squared-distance $\mu \pm d$.
- Their integer-rounded midpoint has squared-distance $\frac{\mu}{4} \pm 2d$ from each endpoint.

Every A–B edge is in few triangles:

- Given $0 = (0, \dots, 0)$ and $z = (z_1, \dots, z_d)$ with $\|z\|^2 = \mu \pm d$.
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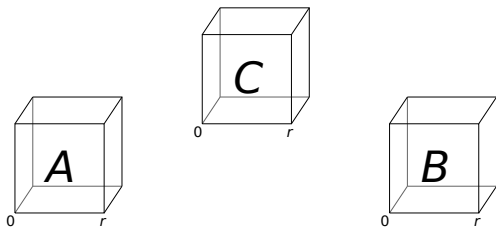
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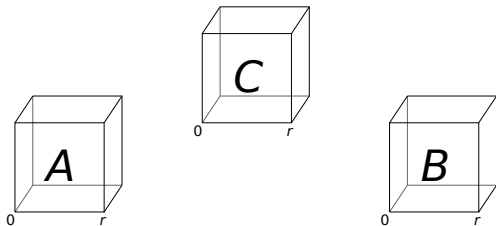
- The number of lattice points in $B_{3\sqrt{d}}^{(d)}$ is at most $15^d \lll r^d$.

FINAL CONSTRUCTION



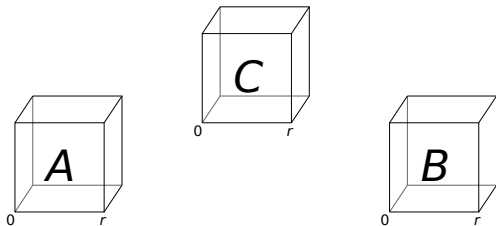
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Blow up (simplification by Alon):

- Replace every vertex in A and B with 2^d copies of itself.
- Now every edge is in at most $30^d \lll (2r)^d$ triangles.