## 51st International Mathematical Olympiad

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## Problems (Day 1)

1. Determine all functions  $f : \mathbb{R} \to \mathbb{R}$  such that the equality

$$f\left(\lfloor x \rfloor y\right) = f(x) \lfloor f(y) \rfloor$$

holds for all  $x, y \in \mathbb{R}$ . (Here  $\lfloor z \rfloor$  denotes the greatest integer less than or equal to *z*.)

2. Let *I* be the incentre of triangle *ABC* and let  $\Gamma$  be its circumcircle. Let the line *AI* intersect  $\Gamma$  again at *D*. Let *E* be a point on the arc  $\widehat{BDC}$  and *F* a point on the side *BC* such that

$$\angle BAF = \angle CAE < \frac{1}{2} \angle BAC.$$

Finally, let G be the midpoint of the segment IF. Prove that the lines DG and EI intersect on  $\Gamma$ .

3. Let  $\mathbb{N}$  be the set of positive integers. Determine all functions  $g : \mathbb{N} \to \mathbb{N}$  such that

$$(g(m)+n)(m+g(n))$$

is a perfect square for all  $m, n \in \mathbb{N}$ .

## Problems (Day 2)

- 4. Let *P* be a point inside triangle *ABC*. The lines *AP*, *BP*, and *CP* intersect the circumcircle  $\Gamma$  of triangle *ABC* again at the points *K*, *L*, and *M*, respectively. The tangent to  $\Gamma$  at *C* intersects the line *AB* at *S*. Suppose that SC = SP. Prove that MK = ML.
- 5. In each of six boxes  $B_1$ ,  $B_2$ ,  $B_3$ ,  $B_4$ ,  $B_5$ ,  $B_6$  there is initially one coin. There are two types of operation allowed:
  - *Type* 1: Choose a nonempty box  $B_j$  with  $1 \le j \le 5$ . Remove one coin from  $B_j$  and add two coins to  $B_{j+1}$ .
  - *Type* 2: Choose a nonempty box  $B_k$  with  $1 \le k \le 4$ . Remove one coin from  $B_k$  and exchange the contents of (possibly empty) boxes  $B_{k+1}$  and  $B_{k+2}$ .

Determine whether there is a finite sequence of such operations that results in boxes  $B_1$ ,  $B_2$ ,  $B_3$ ,  $B_4$ ,  $B_5$  being empty and box  $B_6$  containing exactly  $2010^{2010^{2010}}$  coins. (Note that  $a^{b^c} = a^{(b^c)}$ .)

6. Let *a*<sub>1</sub>, *a*<sub>2</sub>, *a*<sub>3</sub>, ... be a sequence of positive real numbers. Suppose that for some positive integer *s*, we have

$$a_n = \max\{a_k + a_{n-k} \mid 1 \le k \le n - 1\}$$

for all n > s. Prove that there exist positive integers  $\ell$  and N, with  $\ell \le s$  and such that  $a_n = a_\ell + a_{n-\ell}$  for all  $n \ge N$ .

## **Solutions**

1. The answer is f(x) = c for all x, where c = 0 or  $1 \le c < 2$ . To prove that these are the only possible solutions, consider two cases. First suppose that  $\lfloor f(y) \rfloor = 0$  whenever  $0 \le y < 1$ . Then  $f(\lfloor x \rfloor y) = f(x) \lfloor f(y) \rfloor = 0$  whenever  $0 \le y < 1$ . Since every real number can be represented as a product of the form  $\lfloor x \rfloor y$  with  $x \in \mathbb{R}$  and  $0 \le y < 1$ , in this case f is identically zero.

Otherwise, suppose  $\lfloor f(y_0) \rfloor \neq 0$  for some  $0 \leq y_0 < 1$ . For any  $x_n$  satisfying  $n \leq x_n < n+1$ , set  $y = y_0$  and  $x = x_n$  in the given equality to obtain  $f(ny_0) = f(x_n) \lfloor f(y_0) \rfloor$ . Letting  $c_n = \frac{f(ny_0)}{\lfloor f(y_0) \rfloor}$ , it follows that  $f(x_n) = c_n$  for all  $x_n \in [n, n+1)$ . In particular, we have  $\lfloor c_0 \rfloor = \lfloor f(y_0) \rfloor \neq 0$ , hence  $c_0 \neq 0$ . Now set x = y = 0 in the given equality to obtain  $c_0 = f(0) = f(0) \lfloor f(0) \rfloor = c_0 \lfloor c_0 \rfloor$ , hence  $\lfloor c_0 \rfloor = 1$ . Finally, setting y = 0 and x = n in the given equality, we find  $c_n = f(n) = \frac{f(0)}{\lfloor f(0) \rfloor} = \frac{c_0}{\lfloor c_0 \rfloor} = c_0$ . Therefore, in this case we have  $f(x) = c_0$  for all x, and  $\lfloor c_0 \rfloor = 1$ .

This problem was proposed by Pierre Bornsztein of France.

2. Let *P* be the second intersection of ray *EI* and  $\Gamma$ , and let segments *PD* and *FI* meet at *M*. We wish to show that M = G, or, equivalently, FM = MI. Let *Q* be the intersection of segments *PD* and *AF*. Applying Menelaus's theorem to triangle *AFI* and line QMD gives  $\frac{FQ \cdot AD \cdot M}{QA \cdot DI \cdot MF} = 1$ . Hence it suffices to show that  $\frac{FQ \cdot AD}{QA \cdot DI} = 1$  or equivalently that AD/AQ = (DI + DA)/FA.

Triangles *QAD* and *IAE* are similar, so AD/AQ = EA/AI. Also, triangles *ABF* and *AEC* are similar, so we have AF/AB = AC/AE. Together these imply that  $\frac{AD}{AQ} = \frac{AB \cdot AC}{AF \cdot AI}$ . Now, let *H* be the intersection of *BC* and *AD*; notice that triangles *DHC* and *DCA* are similar, hence  $DC^2 = DH \cdot DA$ . Now because  $\angle DCI = \angle CID$ , we have DC = DI, hence  $DA^2 - DI^2 = DA^2 - DC^2 = DA^2 - DH \cdot DA = DA \cdot HA$ . On the other hand, notice that triangles *ABH* and *ADC* are similar, so  $DA \cdot HA = AB \cdot AC$ . Putting these together, we see that  $\frac{AD}{AQ} = \frac{AB \cdot AC}{AF \cdot AI} = \frac{DA \cdot HA}{AF \cdot AI} = \frac{DI + DA}{FA}$ , as needed. This problem was proposed by Wai Ming Tai of Hong Kong and Chongli Wang of

This problem was proposed by Wai Ming Tai of Hong Kong and Chongli Wang of China.

3. All functions of the form g(n) = n + c for a constant nonnegative integer *c* satisfy the problem conditions. We claim that these are the only such functions.

We first show that g must be injective. Suppose instead that g(a) = g(b) for some  $a \neq b$ . Choose n so that n + g(a) = p is prime and greater than |a - b|. From the hypothesis both p(g(n) + a) and p(g(n) + b) must be perfect squares, meaning that g(n) + a and g(n) + b are both divisible by p. But this is impossible, as p > |a - b|. Therefore, g is injective as claimed.

We now show that |g(k + 1) - g(k)| = 1 for all k. Suppose instead that some prime p divides g(k + 1) - g(k). Now, choose an integer n as follows. If  $p^2 | g(k + 1) - g(k)$ , then take n so that n + g(k + 1) is divisible by p but not  $p^2$ . Otherwise, take n so that n + g(k + 1) is divisible by  $p^3$  but not  $p^4$ . Note that the maximum power of p dividing n + g(k + 1) and n + g(k) is odd. Now, the hypothesis implies that (n + g(k + 1))(g(n) + k + 1) and (n + g(k))(g(n) + k) are both squares, meaning that g(n) + k + 1 and g(n) + k are both divisible by p, a contradiction.

For each k, we now have either g(k + 1) = g(k) + 1 or g(k + 1) = g(k) - 1. But g is injective, so if the latter occurs for some k, then it occurs for all k' > k, an impossi-

bility because g takes positive values. Therefore, we have g(k + 1) = g(k) + 1 for all k, hence g(k) = k + g(1) - 1.

This problem was proposed by Gabriel Carroll of the USA.

4. Without loss of generality, we may assume that *S* is on ray *BA*. Set  $x_1 = \angle PAB$ ,  $y_1 = \angle PBC$ ,  $z_1 = \angle PCA$ ,  $x_2 = \angle PAC$ ,  $y_2 = \angle PBA$ , and  $z_2 = \angle PCB$ . Because *SC* is tangent to  $\Gamma$ , we have  $SC^2 = SA \cdot SB$  by the Power of a Point Theorem, and  $\angle SCP = \angle SCM = \angle ACM + \angle ACS = z_1 + \angle ABC = z_1 + y_1 + y_2$ . Because SP = SC, we have  $SP^2 = SC^2 = SA \cdot SB$ , so triangles SAP and SPB are similar. It follows that  $\angle SPA = \angle SBP = y_2$  and that  $\angle ASP = \angle BAP - \angle SPA = x_1 - y_2$ . Now, SP = SC implies  $\angle SPC = \angle SCP = z_1 + y_1 + y_2$ , so  $\angle PSC = 180^\circ - 2(z_1 + y_1 + y_2) = x_1 + x_2 + z_2 - (z_1 + y_1 + y_2)$ . Notice that  $\angle ASC = \angle BAC - \angle ACS = (x_1 + x_2) - (y_1 + y_2)$ , so we have  $\angle ASP = \angle ASC - \angle PSC = z_1 - z_2$ . Combining our two computations of  $\angle ASP$  yields  $x_1 - y_2 = z_1 - z_2$  or  $x_1 + z_2 = y_2 + z_1$ . That is, we have  $(\widehat{KB} + \widehat{BM})/2 = (\widehat{LA} + \widehat{AM})/2$ , hence  $\widehat{KM}/2 = \widehat{LM}/2$  and MK = ML.

This problem was proposed by Marcin E. Kuczma of Poland.

5. The answer is yes. Although the problem specifies that the number of boxes is n = 6, the operations extend in the obvious way to general values of *n*. Our proof will consider several different values of *n* on the way to the final result. For this, it is convenient to let  $(b_1, \ldots, b_n)$  denote the *n*-box configuration where  $b_1$  balls are in box  $B_1, b_2$  balls are in box  $B_2$ , etc. Write  $(b_1, \ldots, b_n) \rightarrow (b'_1, \ldots, b'_n)$  if we can obtain the configuration  $(b'_1, \ldots, b'_n)$  form  $(b_1, \ldots, b_n)$  following the rules in the *n*-box setting. We begin with two lemmas.

LEMMA 1. Let a be a positive integer. Then  $(a, 0, 0) \rightarrow (0, 2^a, 0)$ .

*Proof.* We will show that  $(a, 0, 0) \rightarrow (a - k, 2^k, 0)$  for every  $1 \le k \le a$ , by inducting on k. For k = 1, applying a Type 1 operation to the first number gives  $(a, 0, 0) \rightarrow (a - 1, 2, 0) = (a - 1, 2^1, 0)$ . Now assume the statement holds for some k < a. Starting from  $(a - k, 2^k, 0)$ , repeatedly applying  $2^k$  many Type 1 operations at the middle box yields  $(a - k, 2^k, 0) \rightarrow \cdots \rightarrow (a - k, 0, 2^{k+1})$ . A final Type 2 operation applied at the first box produces  $(a - k, 0, 2^{k+1}) \rightarrow (a - k - 1, 2^{k+1}, 0)$ , completing the induction.

LEMMA 2. Define 
$$P_n = \underbrace{2^{2^{n^2}}}_{n}$$
. Then  $(a, 0, 0, 0) \rightarrow (0, P_a, 0, 0)$  for every positive

integer a.

*Proof.* We will show that  $(a, 0, 0, 0) \rightarrow (a - k, P_k, 0, 0)$  for  $1 \le k \le a$ , by inducting on k. For k = 1, a Type 1 operation applied at the first box gives  $(a, 0, 0, 0) \rightarrow (a - 1, P_1, 0, 0)$ . Now assume that  $(a, 0, 0, 0) \rightarrow (a - k, P_k, 0, 0)$  for some a < k. Applying Lemma 1 to the last three boxes, we obtain  $(a - k, P_k, 0, 0) \rightarrow (a - k, 0, P_{k+1}, 0)$ . A final Type 2 operation applied at the first box gives  $(a - k, 0, P_{k+1}, 0) \rightarrow (a - k - 1, P_{k+1}, 0, 0)$ , completing our induction.

We now describe the construction for the original 6-box setting. Write  $A = 2010^{2010^{2010}}$ . First, apply a Type 1 operation to  $B_5$ , giving  $(1, 1, 1, 1, 1, 1, 1) \rightarrow (1, 1, 1, 1, 0, 3)$ . Second, apply Type 2 operations to  $B_4$ ,  $B_3$ ,  $B_2$ , and  $B_1$  in this order, obtaining  $(1, 1, 1, 1, 1, 0, 3) \rightarrow (1, 1, 1, 0, 3, 0) \rightarrow (1, 1, 0, 3, 0, 0) \rightarrow (1, 0, 3, 0, 0, 0) \rightarrow (0, 3, 0, 0, 0, 0)$ . Third, apply Lemma 2 twice, giving the sequence  $(0, 3, 0, 0, 0, 0) \rightarrow (0, 0, P_3, 0, 0, 0) \rightarrow (0, 0, 0, P_{16}, 0, 0)$ . It is easy to check that  $P_{16} > A$ , so there are more than  $A = 2010^{2010^{2010}}$  coins in  $B_4$  at this point. Fourth, decrease the number of coins in  $B_4$  by applying Type 2 operations repeatedly to  $B_4$  until its size decreases to  $\frac{A}{4}$ . This gives  $(0, 0, 0, P_{16}, 0, 0) \rightarrow \dots \rightarrow (0, 0, 0, \frac{A}{4}, 0, 0)$ . Finally, apply Type 1

operations repeatedly to first empty  $B_4$  and then  $B_5$ , obtaining  $(0, 0, 0, \frac{A}{4}, 0, 0) \rightarrow \cdots \rightarrow (0, 0, 0, 0, 0, \frac{A}{2}, 0) \rightarrow \cdots \rightarrow (0, 0, 0, 0, 0, 0, A)$ , as desired.

This problem was proposed by Hans Zantema of Netherlands.

**Note.** Following a practice established last year, Fields Medalist (and IMO gold medalist) Terence Tao hosted an online project for others to collaborate in solving this problem, which he identified as the most challenging problem on the exam (<u>http:</u>//polymathprojects.org/2010/07/08/minipolymath2-project-imo-2010-q5/).

6. We generalize to the setting where the  $a_n$  may assume negative values. For any  $r \in \mathbb{R}$ , note that the transformation  $a_n \mapsto a_n + rn$  does not change the problem conditions or the result to be proved. Picking  $\ell \leq s$  such that  $a_{\ell}/\ell$  is maximal, we can thus assume without loss of generality that  $a_{\ell} = 0$ . This means all of  $a_1, \ldots, a_s$  are non-positive, hence all  $a_n$  are non-positive. Let  $b_n = -a_n \geq 0$ . For n > s, we have  $b_n = \min\{b_k + b_{n-k} \mid 1 \leq k \leq n-1\}$  and in particular  $b_n \leq b_{n-\ell} + b_{\ell} = b_{n-\ell}$ .

From this, we draw two conclusions. First, all  $b_n$  must be bounded above by  $M = \max\{b_1, \ldots, b_s\}$ . Second, if we let *S* be the set of all linear combinations of the form  $c_1b_1 + c_2b_2 + \cdots + c_sb_s$ , where the  $c_i$  are nonnegative integers, and let  $T = \{x \le M : x \in S\}$ , then since  $b_n = \min\{b_k + b_{n-k} \mid 1 \le k \le n-1\}$ , it is clear that every  $b_n$  must be in *T*. Crucially, *T* is a finite set.

Now, for each integer *i* satisfying  $\ell i + 1 > s$ , let  $\beta_i$  denote the  $\ell$ -tuple  $(b_{\ell i+1}, b_{\ell i+2}, \ldots, b_{\ell i+\ell})$ . By the previous paragraph, the number of such  $\ell$ -tuples is at most  $|T|^{\ell}$ , a finite number. Further, because  $b_n \leq b_{n-\ell}$  for n > s, the individual indices of these  $\beta_i$  are non-increasing functions of *i*. Thus, there can only be finitely many *i* for which  $\beta_i \neq \beta_{i+1}$ . Let  $i_0$  be greater than the largest such value; then, all  $\ell$ -tuples  $\beta_i$  with  $i \geq i_0$  are identical. Choosing  $N = \ell(i_0 + 1)$  finishes the problem, since any  $n \geq N$  gives  $b_n = b_{n-\ell} = b_{\ell} + b_{n-\ell}$ .

This problem was proposed by Morteza Saghafian of Iran. This solution is by Evan O'Dorney.

**Results.** The IMO was held in Astana, Kazakhstan, on July 7–8, 2010. There were 517 competitors from 96 countries and regions. On each day contestants were given four and a half hours for three problems.

On this challenging exam, a perfect score was achieved by only one student, Zipei Nie (China). The USA team ranked third, behind China and Russia. Although the American team has consistently finished in the top ten at the IMO, this year's performance was particularly impressive because none of the team members were in their final year of high school. The students' individual results were as follows.

- Calvin Deng, who finished 9th grade at William G. Enloe High School in Raleigh, NC, won a silver medal.
- Ben Gunby, who finished 10th grade at Georgetown Day School in Washington, DC, won a gold medal.
- Xiaoyu He, who finished 10th grade at Acton-Boxborough Regional High School in Acton, MA, won a gold medal.
- In-Sung Na, who finished 11th grade at Northern Valley Regional High School in Old Tappan, NJ, won a silver medal.
- Evan O'Dorney from Danville, CA, who finished 11th grade (homeschooled through Venture School), won a gold medal. Furthermore, he placed 2nd overall with a score of 39/42. For his spectacular performance, he received a private congratulatory telephone call from the President of the United States, Barack Obama.
- Allen Yuan, who finished 11th grade at Detroit Country Day School in Beverly Hills, MI, won a silver medal.