

# Infinity

①

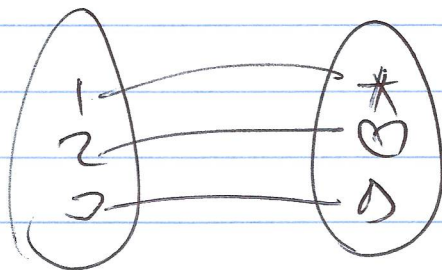
## The concept of size

- we would say  $\{*, \heartsuit, \Delta\}$  has 3 elements, or is of size 3

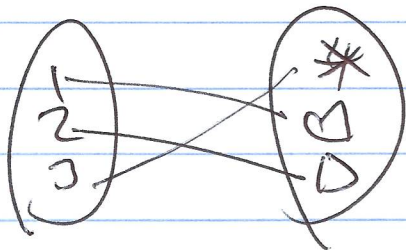
- why? By counting:

$\{*, \heartsuit, \Delta\}$   
1 2 3

- in doing so we are implicitly defining a bijection between  $\{*, \heartsuit, \Delta\}$  and  $\{1, 2, 3\}$



- there is more than one way to do this:



- we will use "existence of a bijection between" as our concept of "of the same size."

Def'n Given sets  $A, B$ , we say  $A, B$  have the same cardinality and write  $A \sim B$  iff  $\exists f: A \rightarrow B$  a bijection.

Note: - in set theory courses one defines, for every set  $A$ , the cardinal number  $|A|$  of  $A$ .  
- then show  $A \sim B$  iff  $|A| = |B|$ .

- defining cardinal numbers beyond our scope: For us,  $|A| = |B|$  is just abbreviation for  $A \sim B$ .

Properties of  $\sim$ :

① For any  $A$ ,  $Id_A: A \rightarrow A$  is a bijection. So  $A \sim A$ , i.e.  $\sim$  is reflexive.

② if  $\exists$  a bijection  $f: A \rightarrow B$  then  $f^{-1}: B \rightarrow A$  is a bijection. So  $A \sim B$  implies  $B \sim A$ ; i.e.  $\sim$  is symmetric.

③ On HW you show: if  $f: A \rightarrow B$  and  $g: B \rightarrow C$  are bijections then  $g \circ f: A \rightarrow C$  is a bijection. So  $A \sim B$  and  $B \sim C \Rightarrow A \sim C$ , i.e.  $\sim$  is transitive.

$\hookrightarrow \sim$  is an equivalence relation on sets!

Def'n Let  $A, B$  be sets.

- ① We write  $A \lesssim B$  (or  $|A| \leq |B|$ ) iff  $\exists$  an injection  $f: A \rightarrow B$
- ② We write  $A \gtrsim B$  (or  $|A| \geq |B|$ ) iff  $\exists$  a surjection  $f: A \rightarrow B$

$\hookrightarrow$  we'd also write  $A < B$  to mean  $A \lesssim B$  and  $A \not\gtrsim B$ .

Notice: -  $A \gtrsim B$  is not "reverse of"  $A \lesssim B$ , i.e. is not asserting an injection from  $B$  to  $A$ .  
- but this follows:

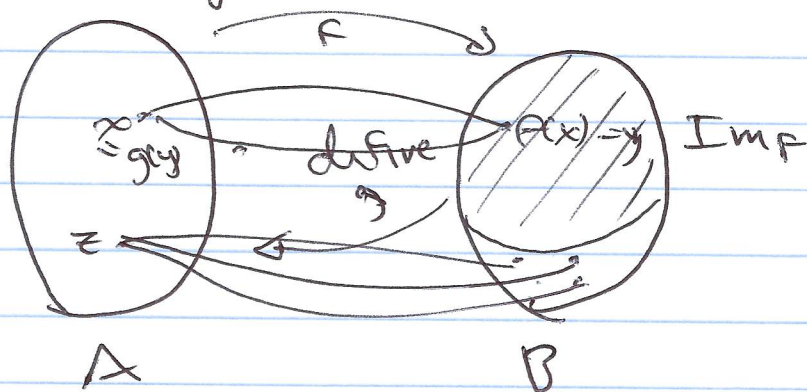
Theorem For all sets  $A, B$  we have  $A \lesssim B$  iff  $B \gtrsim A$ .



PF: (Sketch)

( $\Rightarrow$ ) - Suppose  $A \approx B$   
 - then  $\exists f: A \rightarrow B$  an injection  
 - want to construct  $g: B \rightarrow A$   
 a surjection.

Idea Let  $g$  be "inverse of  $f$ " on  $\text{Im}_f(A)$ , and for anything left over, map to a single element  $z \in A$ .



Explicitly, fix a  $z \in A$ .

Define

$$g = \{ (b, a) \mid (a, b) \in f \} \cup \{ (b, z) \mid b \notin \text{Im}_f \}$$

this says:

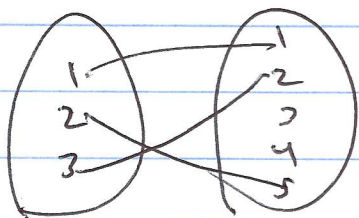
$$g(b) = \begin{cases} a & \text{if } f(a) = b \\ z & \text{if } b \notin \text{Im}_f \end{cases}$$



(5)

then  $g$  is clearly a surjection from  $B$  to  $A$ ,  
i.e.  $B \succeq A$ . ✓

ex: Let  $f: [3] \rightarrow [5]$  be  
defined by  $f(1) = 1$   
 $f(2) = 5$   
 $f(3) = 2$   $\hookrightarrow f$  is an injection



then define  
 $g: [5] \rightarrow [3]$   
by  $g(1) = 1$   
 $g(2) = 3$   
 $g(5) = 2$

and  $g(3) = g(4) = 1$ .

$\hookrightarrow g$  is a surjection ✓

( $\Leftarrow$ ) Suppose  $B \succeq A$ .

- then there is a surjection

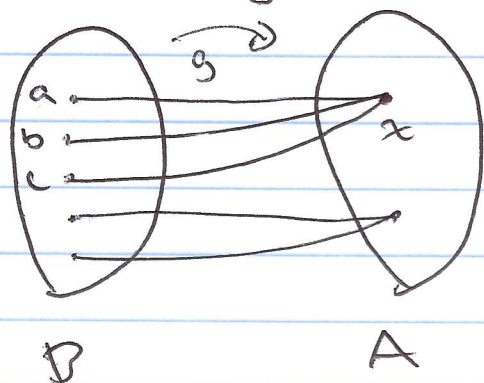
$g: B \rightarrow A$

- want to construct an injection  $f: A \rightarrow B$

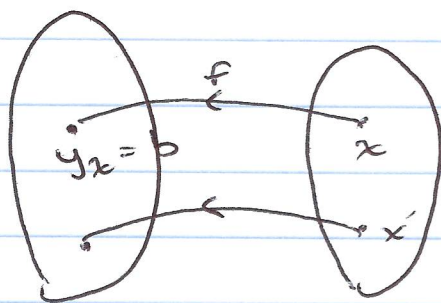
Idea for every  $x \in A$ , pick  
an el<sup>t</sup>  $y_x$  in  $\text{PreIm}_g(x)$

6

and let  ~~$f$~~   $F(x) = y_x$



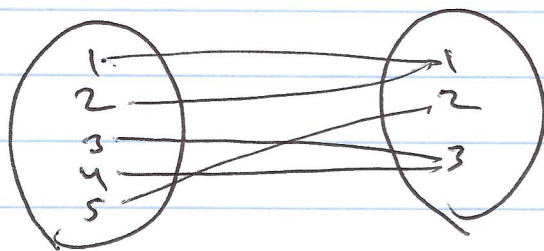
$g(a) = g(b) = g(c) = x$   
~~pick~~ pick one  
say  $y_x = b$   
let  $F(x) = y_x$



then  $F: A \rightarrow B$  will be injective

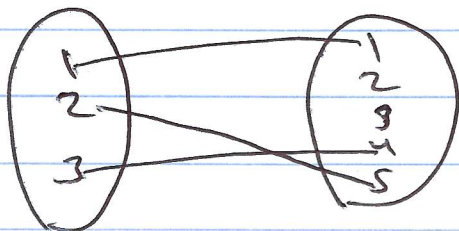
- hence  $A \lesssim B$ . ✓

e.g. consider the scheme for  
 $g: [5] \rightarrow [3]$  below



②

Can derive an injection  $f: [3] \rightarrow [5]$



For infinite sets, to guarantee such an injection exists need to include the axiom of choice.

### Properties of $\approx$

①  $A \approx A$  and  $A \approx A$  since  $\text{id}_A$  is both an injection and surjection so these relations are reflexive. we have injections

② if  $f: A \rightarrow B$  and  $g: B \rightarrow C$  then  $g \circ f: A \rightarrow C$  is an injection, so  $A \approx B$  and  $B \approx C \Rightarrow A \approx C$ .

- Similarly  $A \approx B, B \approx C \Rightarrow A \approx C$ .

- so relations are transitive

③ Are they antisymmetric?



Not literally: if  $A \approx B$  and  $B \approx C$  then  
 we don't have  $A = C$   
 (consider  $\{1, 2, 3\}$  and  $\{*, \square, \Delta\}$ )

↳ but we'll show that in such a case  $A \approx B$ !

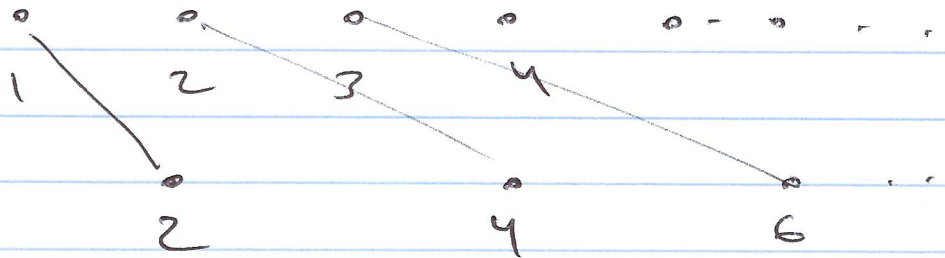
But first on infinite:

Paradoxes of infinity

① "There are as many even ~~whole~~ numbers as whole numbers"

i.e.  $E \approx N$

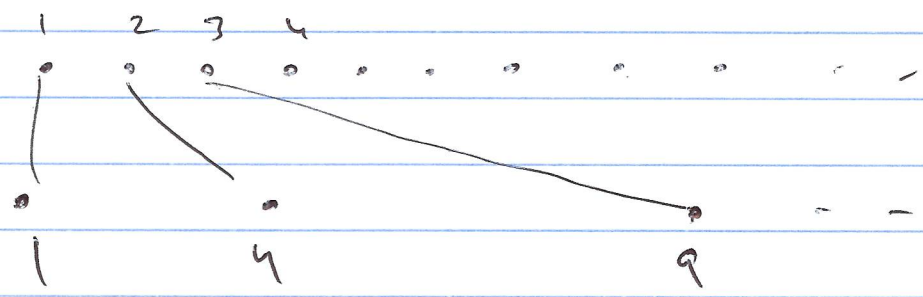
Proof:  $f: N \rightarrow E$  defined  
 by  $f(n) = 2n$  is  
 a bijection.



② (Galileo) "There are as many perfect squares as whole numbers"

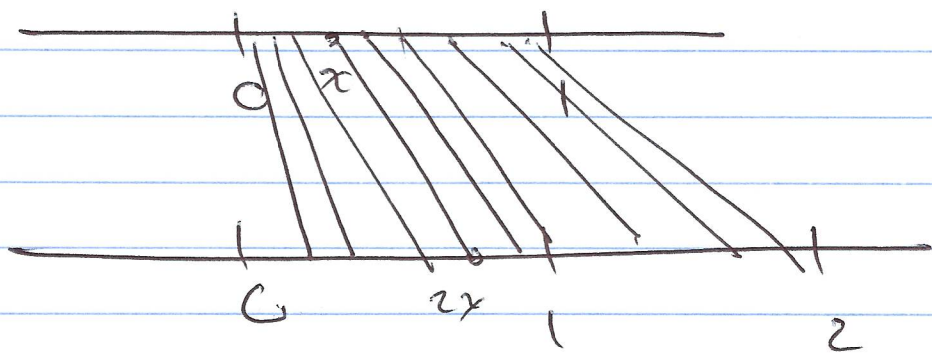
i.e.  $W-S = \{1, 4, 9, 16, \dots\}$   
Then  $S \sim N$

PF:  $F: N \rightarrow S$  defined by  
 $F(n) = n^2$  is a bijection



③ "there are as many points between 0 and 1 as between 0 and 2."

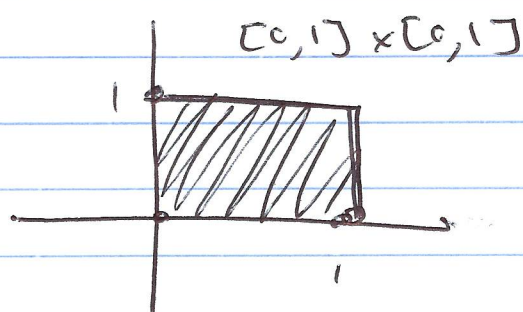
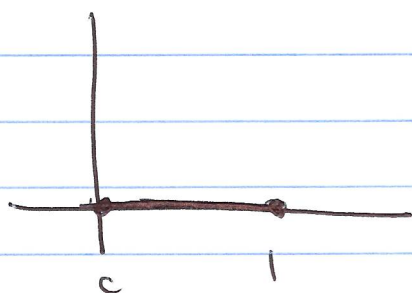
Let  $A = [0, 1]$   
 $B = [0, 2]$   
then  $f: A \rightarrow B$   
 $f(x) = 2x$  is  
a bijection



(10)

(4) "The side is as large as the square"

there is a bijection  $f: [0,1] \rightarrow [0,1] \times [0,1]$



PF. beyond scope.

(8) "There's always a vacancy at Hilbert's Hotel"

- Suppose there is a hotel with infinitely many rooms: #1, 2, 3, 4, ...
- Suppose all are filled and a new guest arrives - how can the staff accommodate her?
- Move everyone down a room and put her in room 1!
- (~~option~~  $f(n) = n+1$  is a bijection from  $\mathbb{N}$  to  $\mathbb{N} \setminus \{1\}$ )



⑥ "A vanishing abundance"

- You find yourself next to an infinite pile of balls and a large bucket.
- At 11 am you take 10 balls from the pile, place them in the bucket, and then remove one ball from the bucket and put it aside (9 balls remain)
- At 11:30 you put 10 more in the bucket, and then remove one and put it aside (19 balls remain)
- At 11:45 you repeat
- And so on: ~~and~~ ~~and~~ whenever the time before noon halves, you repeat.

How many balls are in the bucket at noon?

These examples illustrate the strangeness of ~~of~~ the concept of the size of an infinite set"



Def'n - a set  $X$  is finite iff  $X$  is empty or  $\exists n \in \mathbb{N}$  and a bijection  $f: [n] \rightarrow X$

- a set is infinite iff it is not finite, that is  $X \neq \emptyset$  and  $\forall n \in \mathbb{N}$  there is no bijection  $f: X \rightarrow [n]$ .

The set  $\mathbb{N}$  is the smallest infinite set in the following sense.

Theorem if  $X$  is an infinite set then  $\mathbb{N} \approx X$ .

Pf. - We define an injection  $f: \mathbb{N} \rightarrow X$  "inductively"

- Since  $X$  is not empty  $\exists x_1 \in X$ .

- Let  $f_1 = \{(1, x_1)\}$

- Since  $X$  is not finite,  $f_1$  is not a bijection of  $[1]$  with  $X$ . Hence  $\exists x_2 \in X$ ,  $x_2 \neq x_1$ .

- Let  $f_2 = \{(1, x_1), (2, x_2)\}$



i.e.  $F_n \cup \text{anything}$

- at stage  $n$  suppose we have defined  $F_n = \{(1, x_1), (2, x_2), \dots, (n, x_n)\}$

st.  $\forall i, j \leq n$  if  $i \neq j$  then  $x_i \neq x_j$ .

- then since  $X$  is not finite

$F_n : [n] \rightarrow X$  is not a bijection

- so  $\exists x_{n+1} \in X$  distinct from  $x_1, \dots, x_n$

- let  $F_{n+1} = \{(1, x_1), \dots, (n+1, x_{n+1})\}$

- By induction, the  $N$  we can define on inductor  $F_n : [n] \rightarrow X$

- let  $F = \bigcup_{n \in \mathbb{N}} F_n$

$= \{(1, x_1), (2, x_2), \dots\}$

- then  $F$  is on inductor

$f : \mathbb{N} \rightarrow X$  ✓

- hence  $\mathbb{N} \lesssim X$ .

- This says that among infinite sets,  $\mathbb{N}$  is as small as possible

- hence any infinite sets that appear smaller (e.g.  $\mathbb{E}, \mathbb{Q}$ , etc) are not

- But OTOH,  $\mathbb{N}$  is very large in the sense that many sets  $X$  which seem larger than  $\mathbb{N}$  are not.



(14)

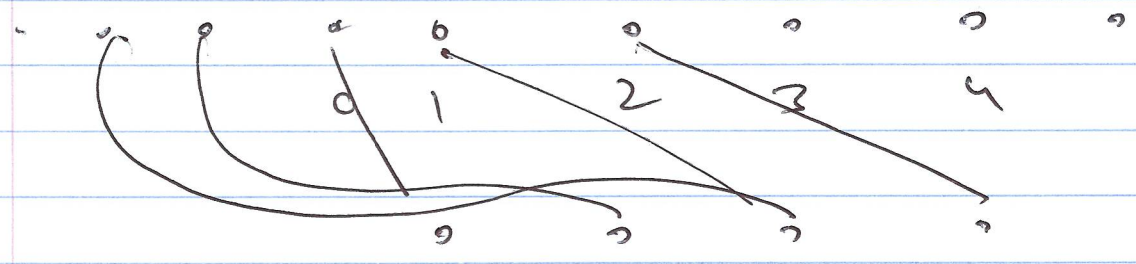
Def'n A set  $X$  is called countable iff  $\mathbb{N} \sim X$ .

Ex's ①  $\mathbb{Z}$  is countable, i.e.  $\mathbb{Z} \sim \mathbb{N}$ .

PF: We showed before that  $f: \mathbb{Z} \rightarrow \mathbb{N}$  defined by

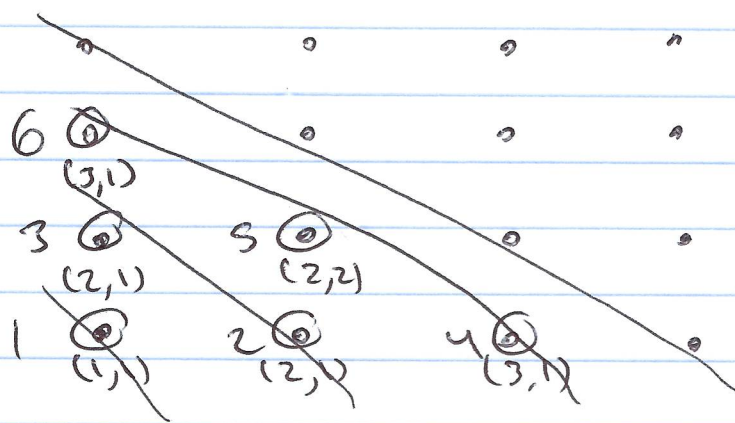
$$f(n) = \begin{cases} 2n & n > 0 \\ 2(-n)+1 & n \leq 0 \end{cases}$$

is a bijection



②  $\mathbb{N} \times \mathbb{N}$  is countable, i.e.  $\mathbb{N} \sim \mathbb{N} \times \mathbb{N}$

PF by picture:



Continuing in this way ("carry along diagonals") we get a bijection

$$f: \mathbb{N} \rightarrow \mathbb{N} \times \mathbb{N}$$

- e.g.  $f(1) = (1,1)$
- $f(3) = (2,2)$
- $f(6) = (3,1)$  etc. ✓

Another approach  
 2.b By a different ~~approach~~ route, we can show  $\mathbb{N} \times \mathbb{N} \lesssim \mathbb{N}$ .

PF: - Day 1 we proved. there are infinitely many primes  
 let  $p_1, p_2, p_3, \dots$   
 be an increasing enumeration

of the primes

So e.g.

$$p_1 = 2$$

$$p_2 = 3$$

$$p_3 = 5$$

$$p_4 = 7$$

$$p_5 = 11$$

etc...

Define:  $f: \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$  by

$$f(n, m) = p_n^m$$

So e.g.

$$f(1, 3) = p_1^3 = 2^3 = 8$$

$$f(2, 2) = p_2^2 = 3^2 = 9$$

$$f(5, 2) = p_5^2 = 11^2 = 121$$

etc...

Claim:  $f$  is an injection

PF:

- Fix  $(n, m), (k, e) \in \mathbb{N} \times \mathbb{N}$
- Suppose  $(n, m) \neq (k, e)$

Case 1:  $n \neq k$

$$\text{- then } f(n, m) = p_n^m \neq p_k^e = f(k, e)$$

different primes

(technically relies on

FTOA, we'll prove that later)



(20)

Case 2:  $n = k$  but  $m \neq l$ .

Then  $f(n, m) = p_n^m \neq p_n^l$   
different powers

$$= p_k^l = f(k, l)$$

In either case  $f(n, m) \neq f(k, l)$   
So  $f$  is injective ✓

Hence  $\mathbb{N} \times \mathbb{N} \approx \mathbb{N}$

We know  $\mathbb{N} \approx \mathbb{N} \times \mathbb{N}$

→ by them or  
 $g: \mathbb{N} \rightarrow \mathbb{N} \times \mathbb{N}$   
 $g(n) = (n, 1)$

From this alone can we  
conclude  $\mathbb{N} \times \mathbb{N} \sim \mathbb{N}$ ?

We need:

Theorem (Cantor - Schroeder - Bernstein)

For any sets  $A, B$ , if  $A \lesssim B$  and  
 $B \lesssim A$  then  $A \sim B$ .

Pf: Fix sets  $A, B$  and suppose  
there are injections

$$f: A \rightarrow B$$

$$g: B \rightarrow A$$

(21)

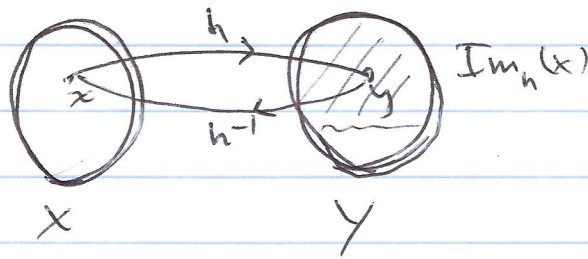
- we need to construct a bijection  $F: A \rightarrow B$ .

- if either  $F$  or  $g$  is a bijection to begin with we are done.

- so assume neither are surjective

Side note - In general if  $h: X \rightarrow Y$  is an injection, then  $h$  is a bijection from  $X$  to  $\text{Im}_h(X)$ .

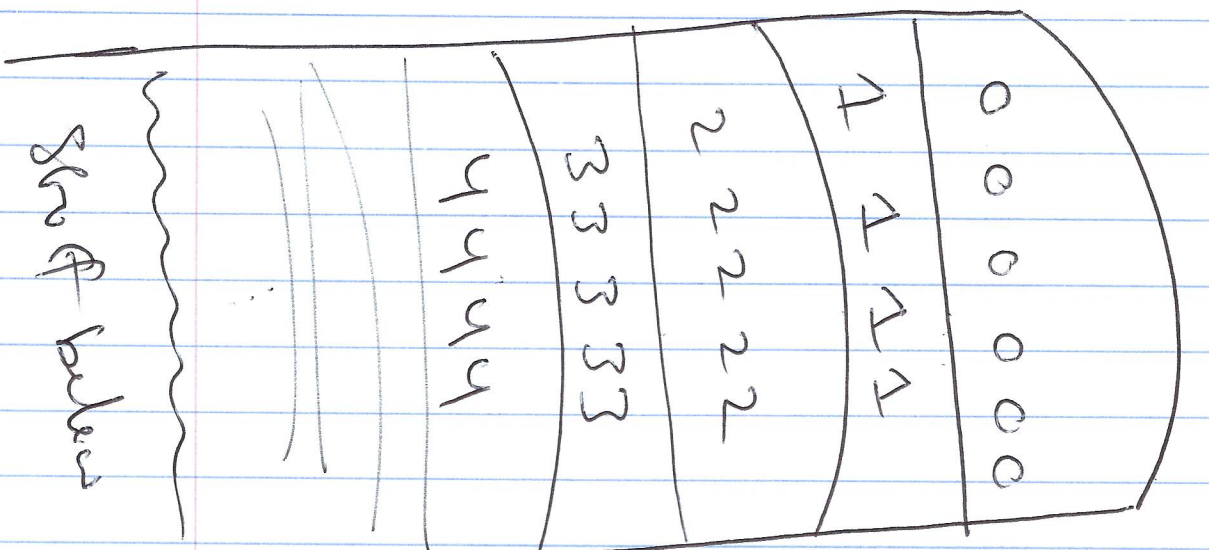
- if  $y \in \text{Im}_h(X)$  we define  $h^{-1}(y) = \text{unique } x \text{ s.t. } h(x) = y$ .



Back to proof:

# THE Picture

$$A \xrightarrow{\Delta} B$$



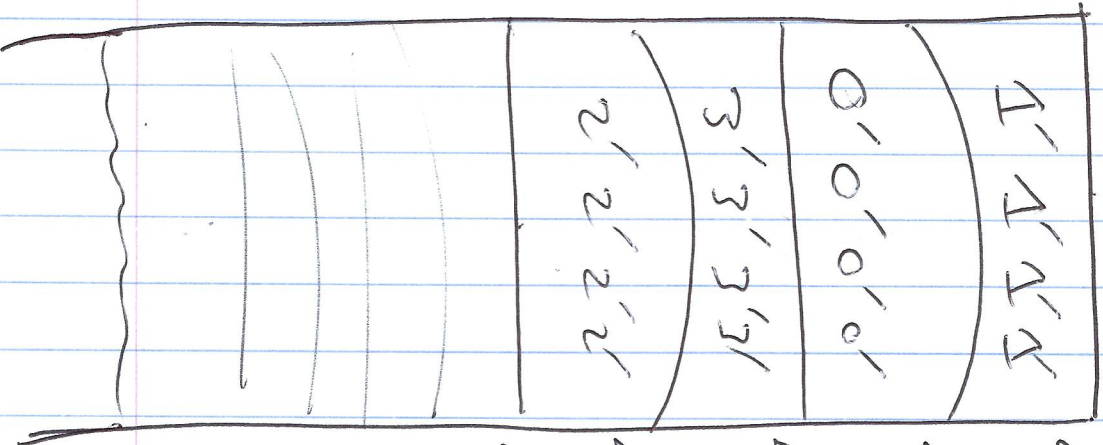
$$A - \text{Im}_g(B)$$

$$\text{Im}_g(B) - \text{Im}_{g^t}(A)$$

$$\text{Im}_{g^t}(A) - \text{Im}_{g^t}(B)$$

$$\text{Im}_{g^t}(B) - \text{Im}_{g^t}(A)$$

$$\text{Im}_{g^t}(A)$$



$$B - \text{Im}_f(A)$$

$$\text{Im}_f(A) - \text{Im}_{f^t}(B)$$

$$\text{Im}_{f^t}(B) - \text{Im}_{f^t}(A)$$

$$\text{Im}_{f^t}(A) - \text{Im}_{f^t}(B)$$

$$\text{Im}_{f^t}(B)$$

Shoof barbas



(23)

$$\begin{aligned} \text{Let } X &= (A - \text{Im}_g(B)) \cup (\text{Im}_{g \circ f}(A) - \text{Im}_{g \circ f \circ g}(B)) \\ &\cup (\text{Im}_{g \circ f \circ g \circ f}(A) - \text{Im}_{g \circ f \circ g \circ f \circ g}(B)) \cup \dots \\ &= (\text{region 0}) \cup (\text{region 2}) \cup (\text{region 4}) \cup \dots \end{aligned}$$

$$\begin{aligned} \text{Let } Y &= \text{everything else in } A \\ &= (\text{region 1}) \cup (\text{region 3}) \cup \dots \\ &\quad \cup (\text{stuff below}) \end{aligned}$$

Define  $F: A \rightarrow B$

$$F(x) = \begin{cases} f(x) & \text{if } x \in X \\ g^{-1}(x) & \text{if } x \in Y \end{cases}$$

Claim  $F$  is a bijection

Pf (Sketch) (i) injectivity: Fix  $x, y \in A$

suppose  $F(x) = F(y)$

Case 1:  $x \in X, y \in X$

$$\begin{aligned} \text{Then } F(x) &= f(x) \\ F(y) &= f(y) \end{aligned}$$

so injectivity of  $f \Rightarrow x = y$

Case 2:  $x \in Y, y \in Y$

$$\text{then } F(x) = g^{-1}(x) \\ F(y) = g^{-1}(y)$$

Since  $g$  is injective,  $g^{-1}$  is injective where it's defined, so  $\Rightarrow x=y$

Cor 3  $x \in X, y \in Y$

Then  $x$  is in an even region  $n$   
 $y$  is in an odd region  $m$   
 (or  $y$  below)

then  $F(x)$  is in an even region  $n'$   
 and  $F(y)$  is in an odd region  $m'$   
 (or  $F(y)$  below)

$\hookrightarrow$  Contradiction as these regions are disjoint and we assumed  $F(x) = F(y)$ .

Cor 4  $x \in Y$  and  $y \in X$   
 Similar to Cor 3.

$\hookrightarrow$  So  $F$  is injective ✓

Surjectivity (even sketcher)

Fix  $y \in B$ .

Cont: -  $y$  is in an ~~even~~ region  $n'$   
 - then there is  $x \in$  region  $n \subseteq X$   
 s.t.  $F(x) = y$   
 hence  $F(x) = y$





(2)

even without producing an explicit bijection  $f: \mathbb{N} \times \mathbb{N} \rightarrow \mathbb{N}$  ourselves.

② We can generalize the above to prove

Theorem: if  $A, B$  are countable sets then  $A \times B$  is countable.

Pf: - Suppose  $A \sim \mathbb{N}$  and  $B \sim \mathbb{N}$ .

- then there are bijections

$$f: A \rightarrow \mathbb{N}$$

$$g: B \rightarrow \mathbb{N}$$

- define

$$F: A \times B \rightarrow \mathbb{N} \times \mathbb{N}$$

$$\text{by } F(a, b) = (f(a), g(b))$$

then  $F$  is a bijection

Pf: you try.

③ Combining ① and ② we can prove that in fact the set of rationals  $\mathbb{Q}$  is countable (amazing!)

Prop'n  $\mathbb{N} \sim \mathbb{Q}$

PF: - Clearly  $\mathbb{N} \lesssim \mathbb{Q}$  since  
 $F: \mathbb{N} \rightarrow \mathbb{Q}$  defined by  $F(n) = n$   
 is an ~~injection~~ injection

- above we proved  $\mathbb{Z} \sim \mathbb{N}$   
 - hence by ②  $\mathbb{Z} \times \mathbb{N} \sim \mathbb{N}$ .

- we now claim:

Claim:  ~~$\mathbb{Z} \times \mathbb{N} \lesssim \mathbb{Q}$~~   $\mathbb{Z} \times \mathbb{N} \gtrsim \mathbb{Q}$ .

PF: - Consider the map  
 $F: \mathbb{Z} \times \mathbb{N} \rightarrow \mathbb{Q}$  defined by  
 $F(m, n) = \frac{m}{n}$ .

- this map is not injective

e.g.  $F(1, 2) = F(3, 6) = \frac{1}{2}$

- but it is clearly surjective

- if  $q \in \mathbb{Q}$  and  $q = \frac{m}{n}$

then  $F(m, n) = q$ .

- So  $\mathbb{Z} \times \mathbb{N} \gtrsim \mathbb{Q}$  as claimed

$\hookrightarrow$  then  $\mathbb{Q} \lesssim \mathbb{Z} \times \mathbb{N} \sim \mathbb{N}$   
 hence  $\mathbb{Q} \lesssim \mathbb{N}$

$\hookrightarrow$  hence by CBS  $\mathbb{N} \sim \mathbb{Q}$ !

(28)

(4) Theorem If  $A_1, A_2, A_3, \dots$  is a sequence of sets indexed by  $\mathbb{N}$ , and  $\forall n \in \mathbb{N}$  we have  $A_n \sim \mathbb{N}$ , then

$$\bigcup_{n \in \mathbb{N}} A_n \sim \mathbb{N}$$

("a countable union of countable sets is countable")

PF - Since each  $A_n$  is ctbl, we have,  $\forall n \in \mathbb{N}$ , a bijection

$$f_n: \mathbb{N} \rightarrow A_n$$

- Can think of  $f_n$  as giving us a way of listing  $A_n$ :

- Let  $a_{n,m}$  denote  $f_n(m)$

- Then:

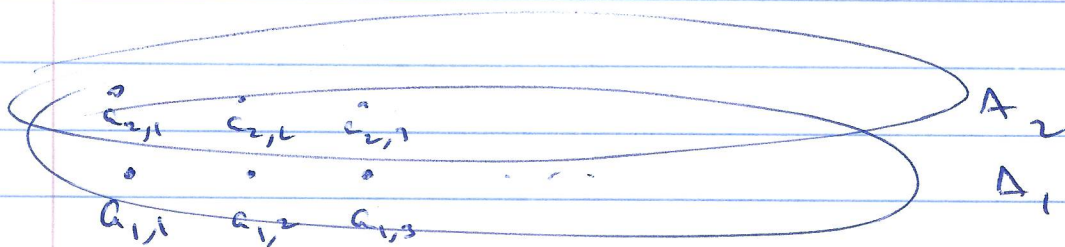
$$A_n: \begin{array}{cccc} f_n(1) & f_n(2) & f_n(3) & \dots \\ a_{n,1} & a_{n,2} & a_{n,3} & \dots \end{array}$$

Now: define  $F: \mathbb{N} \times \mathbb{N} \rightarrow \bigcup_{n \in \mathbb{N}} A_n$   
by  $F(n,m) = f_n(m) = a_{n,m}$

Picture:



(2a)



- Then  $f$  is a surjection:

- Fix  $x \in \bigcup_{n \in \mathbb{N}} A_n$

- Suppose  $x \in A_n$

- Then  $x = f_n(m)$  for some  $m$

- so  $x = f(n, m)$  ✓

↳ Hence  $\mathbb{N} \times \mathbb{N} \supseteq \bigcup_{n \in \mathbb{N}} A_n$

↳ Hence  $\bigcup_{n \in \mathbb{N}} A_n \sim \mathbb{N}$  ✓

- So many examples of sets  $X$  that "appear larger" than  $\mathbb{N}$  but are ctbl

- Q: are there any infinite sets  $X$  which are not countable?

- Yes!

Theorem (Cantor):  $N < P(N)$   
 That is:  $N \lesssim P(N)$  but  $N \not\approx P(N)$

PF: - the proof is by "magic trick"  
 of diagonalization

- we'll show:  $(\forall f: N \rightarrow P(N))$

(1) any function, then there is  
 a set  $T \in P(N)$  that is not in

Im $f$

- hence  $f$  is not a surjection

Claim: Let  $f: N \rightarrow P(N)$  be a  
 function, then  $f$  is not a  
 surjection

PF: Let  $T = \{n \in N \mid n \notin f(n)\}$

Illustration:

Suppose e.g.  $f(1) = \{1, 7, 10\}$

$f(2) = \{1, 3, 5, 7, \dots\}$

$f(3) = \emptyset$

$f(4) = \{2, 4, 6, \dots\}$

$\vdots$

(31)

Then.  $1 \notin T$  because  $1 \in f(1)$   
 $2 \in T$  because  $2 \notin f(2)$   
 $3 \in T$  because  $3 \notin f(3)$   
 $4 \notin T$  because  $4 \in f(4)$   
 $\vdots$

so  $T = \{2, 3, \dots\}$

Claim 1 ( $\forall n \in \mathbb{N}$ )  $f(n) \neq T$

Pf. Fix  $n \in \mathbb{N}$ . - if  $n \in T$  then  
 $n \notin f(n)$  by def'n of  $T$ .  
- hence  $T \neq f(n)$

- if  $n \notin T$ , then  $n \in f(n)$ , hence  
 $T \neq f(n)$  as well.

- Since  $n$  was arbitrary, claim follows

- hence  $f$  is not a surjection  
and original claim follows

- hence there is no bijection  
 $f: \mathbb{N} \rightarrow \mathcal{P}(\mathbb{N})$

i.e.  $\mathbb{N} \not\approx \mathcal{P}(\mathbb{N})$

- Since  $\mathbb{N} \not\approx \mathcal{P}(\mathbb{N})$  we have  $\mathbb{N} < \mathcal{P}(\mathbb{N})$  ✓



The same proof works in general:

Theorem For any set  $A$ , there is no surjection  $f: A \rightarrow P(A)$ .

PF: Fix  $f: A \rightarrow P(A)$

Let  $T = \{ a \in A \mid a \notin f(a) \}$

Then  $(\forall a \in A) f(a) \neq T$ . ✓

- Since  ~~$A \approx P(A)$~~   $f: A \rightarrow P(A)$  defined by

we always have  $A \approx P(A)$   $\cup$  an injection

- So Theorem above says  $A < P(A)$ .

$\rightarrow$  It follows that there are infinitely many levels of infinity

⊗  $\mathbb{N} < P(\mathbb{N}) < P(P(\mathbb{N})) < \dots$

Def'n if  $X$  is infinite and  $\overline{N \times X} = X$  then  $X$  is called uncountable.

↳ above:  $\mathcal{P}(N)$  is uncountable.  
- what about  $\mathbb{R}$ ?

Theorem  $\mathbb{R}$  is uncountable.

Pf: - we'll actually argue that a seemingly "small" subset of  $\mathbb{R}$  is uncountable.

- let  $R = \{x \in (0,1) \mid \text{decimal expansion of } x \text{ contains only 0's and 1's}\}$

- e.g.  $\cdot 11001110101 \dots \in R$
- $\cdot 10101010 \dots \in R$
- $\cdot 111111 \dots \in R$
- $\cdot 123456 \dots \notin R$
- $\cdot 101010 \dots \notin R$

- For  $x \in R$ , define  $x_n$  to be  $n$ th digit in decimal expansion of  $x$

$$x = .x_1 x_2 x_3 \dots$$

So e.g. if  $x = .0010110 \dots$   
then  $x_1 = 0$   
 $x_2 = 0$   
 $x_3 = 1$  etc.

Claim Let  $f: \mathbb{N} \rightarrow \mathbb{R}$  be any function. Then  $f$  is not a surjection.

PF: Fix  $f: \mathbb{N} \rightarrow \mathbb{R}$ .

We define  $r \in \mathbb{R}$  as follows

$$r_n = 0 \quad \text{iff} \quad f(n)_n = 1$$
$$r_n = 1 \quad \text{iff} \quad f(n)_n = 0$$

Picture:

Suppose  $f(1) = .(1)00101.$   
 $f(2) = .0(0)01110$   
 $f(3) = .11(1)10000.$   
 $f(4) = .001(0)0011.$

then  $r_1 = 0$   $r_2 = 0$   $r_3 = 0$   
 $r_4 = 1$  etc  
 $r = .0001 \dots$





Theorem  $F$  is uncountable.

PF: Diagonalize!

Claim: - suppose  $g: \mathbb{N} \rightarrow F$  is a function.

- then  $g$  is not a surjection

PF: define a function

$f: \mathbb{N} \rightarrow \{0,1\}$  as follows

$$f(n) = \begin{cases} 1 & \text{iff } g(n)(n) = 0 \\ 0 & \text{iff } g(n)(n) = 1 \end{cases}$$

then  $\forall n \in \mathbb{N}$   $f \neq g(n)$ . The claim follows

$g(1)$	=	0	1	0	1	0	1	...		
$g(2)$	=	0	0	1	1	0	1	0	1	...
$g(3)$	=	1	1	0	1	1	1	...		
$g(4)$	=	0	0	0	0	1	1	1	...	

$f = 1100...$