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## Nonexamples:

①  $\leq$  is not a strict partial order on  $\mathbb{R}$  since  $\leq$  is not ~~not~~ irreflexive (in fact,  $\leq$  is reflexive which is stronger than being not irreflexive). Similarly,  $\leq$  is not a strict p.o.

②  $\text{OTOH}$ ,  $<$  and  $\neq$  are not (non-strict) partial orders: neither are reflexive (in fact: irreflexive).

③ More generally, a relation  $R$  cannot be both a strict and non-strict p.o. (unless  $R = \emptyset$  is the empty order): can't be both reflexive and irreflexive.

④  $\neq$  (e.g. on  $\mathbb{N}$ ) is neither a strict nor non-strict p.o. since  $\neq$  is not transitive.

## Total orders:

Def'n: a relation  $R$  on  $A$  is said to be total iff

$$(\forall x, y \in A) ((x, y) \in R \vee (y, x) \in R \vee x = y)$$

Def'n: - if  $R$  is a partial order on  $A$  that is also total, then  $R$  is called a total order on  $A$ .

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- if  $R$  is a strict partial order on  $A$  that is also total, then  $R$  is called a strict total order on  $A$ .

Ex's ①  $\leq$  is a total order on  $\mathbb{R}$ : we know  $\leq$  is a p.o. and further:

$$(\forall x, y \in \mathbb{R}) (x \leq y \vee y \leq x \vee x = y) \quad \checkmark$$

②  $\subseteq$  is not a total order on  $P(\mathbb{N})$ :  
it is not total, e.g.:

$$\text{If } X = \{1, 2\} \quad Y = \{2, 3\} \\ \text{then } X \not\subseteq Y, Y \not\subseteq X, X \neq Y.$$

③ Similarly,  $|$  is a p.o. on  $\mathbb{N}$  but is not total,  
e.g.

$$3 \times 5, 5 \times 3, 3 \neq 5.$$

④  $<$  is a strict total order on  $\mathbb{R}$ :  
we know it's a strict p.o. and it's total  
since:

$$(\forall x, y \in \mathbb{R}) (x < y \vee y < x \vee x = y).$$

⑤  $\subsetneq$  is not a strict total order on  $P(\mathbb{N})$ :  
it's not total.

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## Functions

- Functions, like relations, are ubiquitous in math. But what "are" functions?

- intuitively: a function is a rule that assigns to each input  $x$  in a domain  $A$  a unique output  $f(x)$  in a codomain  $D$ .

Idea: can define functions ~~as~~ rigorously as a special type of relation.

Def'n: a function (with domain  $A$  and codomain  $B$ ) is a relation  $f \subseteq A \times B$  s.t.

for every  $a \in A$   
there is a unique  $b \in B$   
s.t.  $(a, b) \in f$ .

1-1.

$$(\forall a \in A) (\exists b \in B) [(a, b) \in f \wedge$$

$$(\forall c \in B) ((a, c) \in f \Rightarrow c = b)]$$

Notation: we write

$$f: A \rightarrow B$$

to mean a relation  $f \subseteq A \times B$  is a function.

(iv) We also write:

$$f(a) = b$$

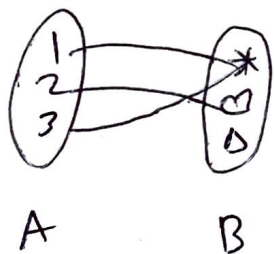
to mean  $(a, b) \in f$ .

Note: - def'n says: every  $a \in A$  assigned an output ~~is~~  $f(a) \in B$ .

- does not insist for every  $b \in B$  there is  $a \in A$  s.t.  $f(a) = b$  (functions w/ this property are called surjective).

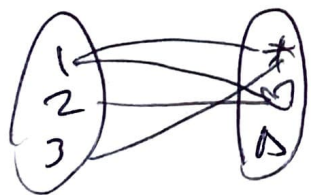
Ex's: ① Let  $A = \{1, 2, 3\}$ ,  $B = \{*, \heartsuit, \Delta\}$ .

Then:  $f = \{(1, *), (2, \heartsuit), (3, *)\}$  is a function from  $A$  to  $B$ :



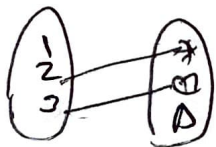
a function

but  $g = \{(1, *), (1, \heartsuit), (2, \heartsuit), (3, *)\}$  is not since 1 does not have a unique output: both  $(1, *)$  and  $(1, \heartsuit) \in g$ .



not a function.

(v) not  $u \mapsto h = \{(2, *), (3, \square)\}$  since 1 is not assigned an output.



not a function

② we'll often define functions by a rule:

e.g.  $f: \mathbb{R} \rightarrow \mathbb{R}$

$$f(x) = x^2$$

or:  $g: \mathbb{R} \rightarrow \mathbb{Z}$

$$g(x) = \lfloor x \rfloor.$$

but behind the scenes these  $f$ 's are still sets of ordered pairs:

e.g. for  $f$  above we have:

$$(2, 4) \in f$$

$$(3, 9) \in f$$

$$\text{but } (7, 12) \notin f$$

(since  $f(7) = 49 \neq 12$ )

Warning: not all rules yield well-defined functions!

e.g. suppose we "define"

$$f: \mathbb{Q} \rightarrow \mathbb{Z}$$

by the rule  $f(m/n) = m+n.$

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then this "function" is not one:

$$f(1/2) = 1+2 = 3 \neq 6 = 2+4 = f(3/4)$$

but  $1/2 = 2/4$ .

- so  $f$  assigns multiple outputs to the same input. what's going on?

- Really: there's an implicit equiv. relation defined on fraction representations ( $1/2 = 2/4 = 3/6 = \dots$ ) and our rule for  $f$  is defined in terms of the representative of an equiv. class, not the class itself.

- in general: when given a rule "defining"  $f \subseteq A \times B$ , to verify  $f$  is a function one must show:

①  $(\forall a \in A)(\exists b \in B) (a, b) \in f$

② if  $a = a'$  then  $f(a) = f(a')$ .

### Equality of Functions:

Q: What does it mean for functions  $f: A \rightarrow B$  and  $g: A \rightarrow B$  to be equal?

A:  $f = g$  iff they're equal as sets (of ordered pairs) i.e. iff  $f \subseteq g$  and  $g \subseteq f$ .

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In practice: often easier to use following criterion:

Thm: If  $f: A \rightarrow B$  and  $g: A \rightarrow B$  are functions then  $f = g$  iff  $(\forall a \in A) (f(a) = g(a))$

PF: you try.

The point: functions can be equal despite being defined by different rules.

e.g.: let  $A = \{1, 2, 3\}$

define  $f: A \rightarrow \mathbb{N}$

$g: A \rightarrow \mathbb{N}$

by  $f(x) = x^3 + 11x$

$g(x) = 6x^2 + 6$

then:  $f(1) = 12 = g(1)$

$f(2) = 30 = g(2)$

$f(3) = 60 = g(3)$

i.e.  $f = \{(1, 12), (2, 30), (3, 60)\} = g$ .

(What's the magic trick?)

$$f - g = x^3 - 6x^2 + 11x - 6 = (x-1)(x-2)(x-3) = 0 \text{ for } x \in \{1, 2, 3\}$$

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## Images

Def'n: Sp's  $f: A \rightarrow B$  is a function and  $X \subseteq A$ .  
The image of  $X$  under  $f$ , denoted  $\text{Imp}(X)$ ,  
is defined as:

$$\text{Imp}(X) = \{b \in B \mid (\exists a \in X) f(a) = b\}$$

informally we could write:

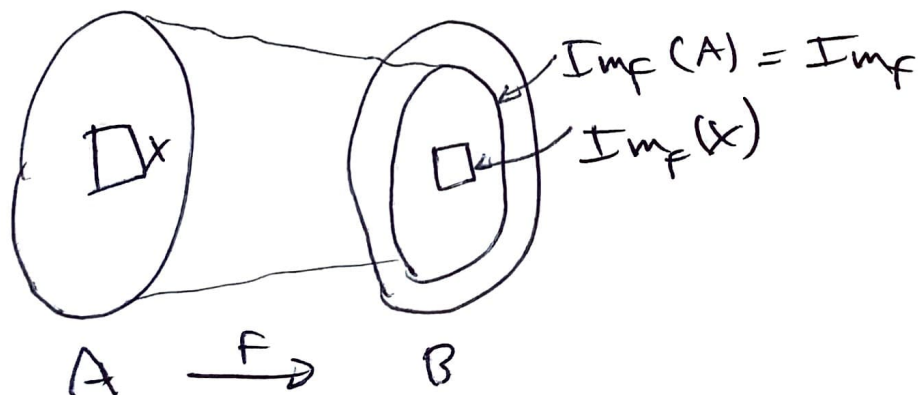
$$= \{f(a) \mid a \in X\}$$

→ the point:  
( $f a \in A$  then  $f(a) \in \text{Imp}(A)$ ).

↳ when  $X = A$  we just say  $\text{Imp}(A)$   
is the image of  $f$  and sometimes just  
write  $\text{Imp}$ .

Def'n says: -  $\text{Imp}_f(X)$  is the "set of outputs  
of el'ts in  $X$ "  
-  $\text{Imp}_f = \text{Imp}_f(A)$  is the "set of  
all outputs".

Picture:





(12) Ex's ① Let  $A = \{1, 2, 3\}$   $B = \{*, \heartsuit, \Delta\}$   
 $f: A \rightarrow B$  be  $\{(1, *), (2, \heartsuit), (3, *)\}$

Then:  $-\text{Im}_f(\{1, 3\}) = \{f(1), f(3)\}$   
 $= \{*, *\}$   
 $= \{*\}$ .

$-\text{Im}_f = \text{Im}_f(\{1, 2, 3\}) = \{f(1), f(2), f(3)\}$   
 $= \{*, \heartsuit, *\}$   
 $= \{*, \heartsuit\}$ .

② Define  $f: \mathbb{R} \rightarrow \mathbb{R}$  by  $f(x) = x^2$ .

Then:  $-\text{Im}_f(\{-1, 0, 1\}) = \{(-1)^2, 0^2, 1^2\}$   
 $= \{0, 1\}$ .

$-\text{Im}_f = \{x \in \mathbb{R} \mid x \geq 0\}$ .

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Functions add a layer of complexity to the basic set theory of  $\cup, \cap, \dots$  that we studied earlier. E.g.:

Prop'n: Sp.  $f: A \rightarrow B$  is a function and  $S, T \subseteq A$ .

Then:  $\text{Im}_f(S \cap T) \subseteq \text{Im}_f(S) \cap \text{Im}_f(T)$ .

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Pf.: - fix  $y \in \text{Imp}(S \cap T)$

- then  $\exists x \in S \cap T$  s.t.  $f(x) = y$

- hence  $x \in S$  and  $x \in T$

- hence  $f(x) \in \text{Imp}(S)$  and  $f(x) \in \text{Imp}(T)$

- i.e.  $y \in \text{Imp}(S)$  and  $y \in \text{Imp}(T)$

- i.e.  $y \in \text{Imp}(S) \cap \text{Imp}(T)$

Since  $y$  was arbitrary, the prop'n is proved. ✓

Note: in general we don't have

$$\text{Imp}(S \cap T) = \text{Imp}(S) \cap \text{Imp}(T)$$

e.g. Consider  $f(x) = x^2$  on  $\mathbb{R}$ .

$$\text{Let } S = \{-1, 0\} \quad T = \{0, 1, 2\}$$

$$\text{then } \text{Imp}(S) = \{f(-1), f(0)\} = \{1, 0\}$$

$$\text{Imp}(T) = \{f(0), f(1), f(2)\} = \{0, 1, 4\}$$

$$\Rightarrow \text{Imp}(S) \cap \text{Imp}(T) = \{0, 1\}$$

$$\text{but: } \text{Imp}(S \cap T) = \text{Imp}(\{0\}) \\ = \{f(0)\} = \{0\}.$$

So in this case:  $\text{Imp}(S \cap T) \subsetneq \text{Imp}(S) \cap \text{Imp}(T)$ .  
the essence of the issue: functions can send multiple inputs to the same output.