Lecture 36 (12/3). Please enable video if you can

- Let $\tau_1 = \min\{n > 0 \mid S_n = 0\}$, be the first time S returns to 0.
- Let τ₂ = min{n > τ₁ | S_n = 0}, be the first time after τ₁ that S returns to 0.
 Let τ_{k+1} = min{n ≥ τ_k | S_n = 0}, be the first time after τ_k that S returns to 0.

Lemma 9.5. S is recurrent at 0 if and only if $P(\tau_{p} < \infty) = 1$.



Elemma 9.6.
$$P(\tau_{p} < \infty) = 1$$
 if and only if $\sum_{N=0}^{\infty} P(S_{n} = 0) = \infty$.
Proof.
(By 9.6, $Z P(S_{n} = 0) = \infty \iff S$ is recurred of D).
P(S_{n} = 0) = $\sum_{N=0}^{\omega} P(S_{n} = 0)$
($: E(\# \text{fines } m + S_{m} = 0) = E(Z \text{ If } S_{m} = 0)$
 $= Z E \text{ If } S_{m} = 0$

 $(2) E(\# fimes m + S_m = 0) = \frac{1}{1 - P(\tau, < p)}$ $P_{1}: F(\#_{1}, m_{g}, m_{g} + S_{m} = 0) = F(\sum_{n=1}^{\infty} \frac{1}{2} \sum_{n=1}^{\infty} \frac{1}{2$ $= \frac{2}{2} P(\tau_{n} < \lambda) = \frac{2}{2} P(\tau_{1} < \lambda)$ $= \frac{P(z_1 < \omega)}{1 - P(z_1 < \omega)} \qquad \left(\begin{array}{c} P(z_2 < \omega) = P(z_1 < \omega) \end{array} \right)$

Equale ()
$$\mathcal{L}(2)$$
: $E(\# \text{ fine, } m + S_m = 0) = \mathbb{Z}P(S_n = 0) = P(\tau_1 < \infty)$
 $I - P(\tau_1 < \omega)$

$$P(s_n = 0) = 0 \quad \langle = \rangle \quad P(\tau_1 \neq \langle h \rangle) = 1$$
QED.

Theorem 9.7. $|\mathbf{P}(S_{2m} = 0) = O(1/m^{d/2})$. Consequently, the random walk is recurrent for $d \leq 2$, and transient for $d \geq 3$. S_m = RW in d dim. Assume $P(S_{2M}=0) = O(\frac{1}{m}d_2).$ Then Sisner of () (=> ZP(S_1=0) = A $(=) \sum P(S_{2n}=0) = 10$ $\stackrel{\scriptstyle ()}{=} \frac{1}{2} \stackrel{\scriptstyle ()}{=} \frac{1}{\sqrt{2}} = 1$

Note lampoison tect : Z 1 K K (=> K > 1 . Sis we at 0 (=> d/ < 1 (=> d < 2 QED.

 $lowerke P(S_{2h} = D)$. d = 1: $P(S_{2N} = O) = {\binom{2N}{N}} \cdot {\binom{1}{2}}^{n} {\binom{1}{2}}^{n}$ $= \frac{1}{2n} \begin{pmatrix} 2n \\ n \end{pmatrix} = \frac{1}{2n} \frac{(2n)!}{(n!)^2} \dots$

Lemma 9.8 (Sterling's formula). For large n, we have

$$\underline{\underline{n!}} \approx \sqrt{2\pi n} \left(\frac{\underline{n}}{\underline{e}}\right)^{\underline{n}} = \sqrt{2\pi} \exp\left(\underline{n \ln n} - n + \frac{\ln n}{2}\right).$$

Proof of Theorem 9.7 for d=1:

 $\sum_{\substack{2h\\ 2\\ 2}}^{1} \sqrt{2} \frac{\sqrt{n}}{2} \frac{\sqrt{n}}{2} \frac{(2n)^{2n}}{2}$

 $\sim \frac{R}{\sqrt{n}} \frac{1}{\sqrt{N}}$ as claimed

 \Rightarrow Thu for d = 1 QED



Theorem 9.10. Consider the Gamblers rule example, with $\tau = \min\{n \mid X_n = 1\}$. Then

$$\underbrace{\boldsymbol{E}\tau = \infty} \quad and \quad \mathbf{P}(\tau = 2\underline{n-1}) = (-1)^{n-1} \binom{1/2}{n} \approx \frac{C}{\underline{n^{3/2}}}$$

Remark 9.11. Let $M_n = \min\{X_{\tau \wedge k} \mid k \leq \underline{n}\}$. Then $\underline{E}M_{\tau} = -\infty$. Thus, this strategy will take (on average) an infinite time before you win \$1. During that time your expected maximum loss is $-\infty$.



 $\tau' = finst time after \tau', X neaches + 1$ dist of T' = dist of T" = dist of T & T" & T are jud $\Rightarrow E x^{T} = \frac{1}{2} x f \frac{1}{2} x E x^{T'} E x^{T'}$ $= \frac{1}{2}x + \frac{\chi}{2}\left(E_{x}x^{T}\right)^{2} \implies F(x) = E_{x}x^{T} = \frac{\chi}{2} + \frac{\chi}{2}F(x)^{2}$ Salve for F(x) & get the fault. F(x) =

Proof of Theorem 9.10

 $() E \tau = p.$ $P_{f}: F(x) = E x^{T} = \frac{1 - \sqrt{1 - x^{2}}}{1 - \sqrt{1 - x^{2}}}$ diff wat a : F'(x) = FT $\operatorname{Put} x=1$: $F'(1) = E_T$ Find F(1) from founde legit ET = + 00 .

(2) Final $P(\tau = n)$ $|-\sqrt{|-\chi^2|}$ $\operatorname{Km} E x^{T} = F(a) =$ $F_{\chi}^{2} = |P(\tau=0) + \chi P(\tau=1) + \chi^{2} P(\tau=2) + \cdots$ $= \sum_{k=1}^{\infty} \chi_{k}^{k} P(\tau = k)$ $F(x) = E x^2 = \frac{1 - \sqrt{1 - x^2}}{x}$. Taylor coires le find a famla.