Theorem 6.62. Consider the binomial model with 0 < d < 1 + r < u, and an American option with intrinsic value G. Define

$$V_N = \underbrace{G_N}_{N}, \quad \underbrace{V_n}_{N=1} = \max\left\{\frac{1}{D_n} \underbrace{\tilde{E}_n}_{n+1} \underbrace{V_{n+1}}_{N+1}, \underbrace{G_n}_{n+1}\right\}, \quad \underbrace{\sigma^*}_{n+1} = \min\left\{\underline{n \leq N} \mid \underbrace{V_n}_{n+1} = \underbrace{G_n}_{n}\right\}.$$

Then V_n is the arbitrage free price, and σ^* is the minimal optimal exercise time. Moreover, this option can be replicated. Remark 6.63. The above is true in any complete, arbitrage free market.

Remark 6.64. In the Binomial model the above simplifies to:

$$V_{n,\text{M}}(\omega) = \max\left\{\frac{1}{1+r}\left(\underline{\tilde{p}}V_{n+1}(\omega',\underline{1}) + \underline{\tilde{q}}V_{n+1}(\omega',-\underline{1})\right), G_{n}(\omega)\right\}, \quad \text{where } \omega = (\underline{\omega}', \underline{\omega}_{n+1}, \omega''), \quad \omega' = (\omega_{1}, \dots, \omega_{n}).$$

Prod i $V_{M} = AFP$ at fine α (fine) of 1
 $IOU : V^{*} = \text{wind}$ of 1 (fine) of 1
 $IOU : V^{*} = \text{wind}$ of 1 (fine) of 1
 $IOU : P^{*} = \text{wind}$ of 1 (fine) of 1

Ma is a mantingale (under P) if $E_n M_{n+1} = M_n$ Pef: @ We say a process Ma is a safeer - ong (under \$PP) $E_{M}M_{M+1} \leq M_{N}$ (b) M is a sub-ma if $E_n M_{nn} \gg M_n$. Pf of Obs 1: NTS DuVn is a super mg mor P. i.e. NTS $\widetilde{E}_{n}(D_{n+1}, V_{n+1}) \leq D_{n} V_{n}$

> I a preditable line process A & a my M + $R_{M}V_{M} = (M_{M}) - A_{M} (Choose A = 0)$ Set $X_{n} = \frac{M_{n}}{D_{n}} \longrightarrow D_{n}V_{n} = (D_{n}X_{n} - A_{n})$ (On HW today -Explicit famle p A). 3) => X_M = wealth of some set fin faitabio $k \quad (z) \quad (v_n \leq u < \tau^* \leq v_n > G_n)$

Note D+(2) >> X is the weath of the vep fort of american often. (2 Self fin) (& v* is an appinal exercise time). 3 => 5 tie the minimal optimal exercise time. ("," Suy T is any storping time $+ T \leq t^*$. $\lambda P(T < t^*) > 0$ Then $\widetilde{E}(P_{\mathcal{F}}X_{\mathcal{F}}) = E(P_{\mathcal{F}}^{*}G_{\mathcal{F}}^{*})$ will return to this sheatly. tor

 $P_{f} = P_{f} = \frac{1}{2} \operatorname{daim}_{i} \quad (P \times X_{n} \geq G_{n}).$ Note $D_n V_n = D_n X_n - A_n \rightarrow X_n = V_n + \begin{pmatrix} A_n \\ D_n \end{pmatrix}$ Note An is increasing & A = 0 >> An > 0 $\gg \chi_n \geq V_n = \max \{ G_n \} =$ $\Rightarrow X_n \ge G_n$. $\neq QED.$ (lam 2: X = VG

 \gg Hun $M \neq T^*$, $D_M V_N = \tilde{E}_M (D_{N+1} V_{N+1})$ Note $D_n X_n = \tilde{E}_n (D_{n+1} X_{n+1}) \quad \forall n.$ $L D_n V_n = D_n X_n + A_n$ $1_{\{m \notin \tau^*\}} \tilde{E}_n(D_{n+1} \vee_{n+1}) = \tilde{E}_n(1_{\{m \notin \tau^*\}} D_{n+1} \vee_{n+1}) =$ JARGER DUVN.

 $\frac{1}{\{m \in \mathcal{T}^{\text{res}}\}} \stackrel{\sim}{\in} \mathcal{L}_{m} \left(\begin{array}{c} \mathcal{D}_{m+1} \\ \mathcal{M}_{m+1} \\ \mathcal{M}_{m+1} \end{array} \right)$ (": A is pred). $\frac{1}{2} \left(\begin{array}{c} D_{u} X_{n} - A_{n+1} \end{array} \right) = \frac{1}{2} \left(\begin{array}{c} D_{u} X_{n} - A_{n+1} \end{array} \right) = \frac{1}{2} \left(\begin{array}{c} D_{u} X_{n} - A_{n+1} \end{array} \right)$ $= \int_{\{n \in V^*\}} \left(\frac{D_n X_n - A_n + A_n - A_{n+1}}{D_n V_n} \right) = \int_{\{n \notin V^*\}} \frac{D_n V_n}{D_n V_n}$

 $\Rightarrow 1_{\xi_{M} \notin \tau^{*}\xi} (A_{nn} - A_{n}) = ()$ $\Rightarrow A_{m} = 0 \quad \forall m \leq \nabla^{*},$ (Will finish he wet arest time)