hart time :

RNP former: $V_n = \frac{i}{P_n} \tilde{E}_n (P_N V_N) \ll uater functional to compute with$ (Computation fine -> 2)

Theorem 7.2. Suppose a <u>security</u> pays $V_N = |\underline{g}(S_N)|$ at maturity N for some (non-random) function g. Then the arbitrage free price at time $n \neq N$ is given by $V_n = f_n(S_n)$, where: (1) $f_N(x) = \mathbb{W}_N(x)$ for $\underline{x} \in \text{Range}(\underline{S_N})$. (2) $\underline{f_n}(x) = \frac{1}{1+r} (\tilde{p}f_{n+1}(\underline{u}x) + \underline{\tilde{g}}f_{n+1}(\underline{d}x)) \text{ for } x \in |\text{Range}(S_n).$ *Remark* 7.3. Reduces the computational time from $O(\underline{2^N})$ to $O(\sum_{0}^{N} |\text{Range}(S_n)|) = O(N^2)$ for the Binomial model. Remark 7.4. Can solve this to get $f_n(x) = \sum_{k=0}^{N-n} {\binom{N-n}{k}} f_N(xu^k d^{N-n-k}) \cdot \frac{1}{(1+\gamma)^N - N} \xrightarrow{\gamma k} \gamma N - N - k$ $O\left[\operatorname{conducte} \left\{ f_{A} \right\}_{-1}(x) = \frac{1}{1+r} \left(\stackrel{\sim}{p} \left\{ f_{A} \right\}_{N}(x \times) + \stackrel{\sim}{q} \left\{ f_{N}(A \times) \right) \right)$ $= \frac{1}{4\pi} \left(\frac{1}{p} \left(\frac{1}{4\pi} \left(\frac{1}{p} \right) + \frac{1}{p} \left(\frac{1}{4\pi} \left(\frac{1}{p} \right) + \frac{1}{p} \left(\frac{1}{4\pi} \left(\frac{1}{p} \right) + \frac{1}{2\pi} \left(\frac{1}{4\pi} \left(\frac{1}{p} \right) + \frac{1}{2\pi} \left(\frac{1}{4\pi} \left(\frac{1}{p} \right) + \frac{1}{2\pi} \right) \right) \right) + \frac{1}{2\pi} \left(\frac{1}{2\pi} \left(\frac{1}{p} \right) + \frac{1}{2\pi} \left(\frac{1}{2\pi} \right) \right) \right) \right)$ $+\tilde{q}_{-}(\frac{1}{1+r}(\tilde{b}_{+N}(udx) + \tilde{q}_{+N}(d^{2}x))))$

 $=\frac{1}{(1+r)^2}\left[\frac{7^2}{7^2}\int_{\mathcal{W}}(h^2x) + 2\frac{7}{7^2}\int_{\mathcal{W}}(hdx) + \frac{7^2}{7^2}\int_{\mathcal{W}}(d^2x)\right]$ 3 Iterate this & get Rule 7.4.

Question 7.5. How do we handle other securities? E.g. Asian options (of the form $g(\sum_{k=1}^{N} S_{k})$)? European Lations -> try off (SN-K)⁺ Asim call often ~ bryall (1 ZSk - K) & New (American call -> can exercise at my time < N) & IDU

Definition 7.6. We say a process X is a *Markov process* if $P(X_{n+1} = x_{n+1} | X_1 = x_1, \dots, X_n = x_n) = P(X_{n+1} = x_{n+1} | X_n = x_n)$. **Theorem 7.7.** A process X is Markov if and only if for every (bounded, continuous) function f, there exists a function g such that $\boldsymbol{E}_n f(X_{n+1}) = g(X_n).$ Question 7.8. If X_n represents i.i.d. coin tosses, is X_n Markov? Is $Y_n = \sum_{k=0}^{n} X_k$ Markov? $E_{n} \left\{ \left(X_{n+1} \right) = E \left(\left\{ \left(X_{n+1} \right) \mid Y_{n} \right\} \xrightarrow{Markov} g \left(X_{n} \right) \right\}$ (X_{n+1}) indep $\in (X_{n+1}) \stackrel{\text{ind}}{=} \in ((X_{n}) (\text{some } \#)$ $(\text{chase } g(n) = F_{4}(X) \quad \forall n)$ $X_{n} \rightarrow iid$. $Y_{n} = Z X_{k}$ $(\gamma = 0)$

Q: Is Y marker Shues -> No > finece > Yes!

Note: $Y_{m+1} = Y_m + X_{m+1}$ Compte $E_{M}(Y_{n+1}) = E_{M} \left\{ \left(Y_{m} + X_{n+1} \right) \right\}$ indep leve $\sum \left\{ \left(\begin{array}{c} x + \pi \end{array} \right) \begin{array}{c} P(X = \pi) \\ P(X = \pi) \end{array} \right\}$ = some for of /n. > [Y is Markov] Fi #'s.

Question 7.9. Is
$$S_n$$
 (stock in the Binomial model) Markov under \vec{P} ? Is $A_n = \frac{1}{n} \sum_{0}^{n} S_k$ Markov under \vec{P} ?
 $h_{WSS} S_n$ is markov $h_{Mn} = \begin{cases} M & \omega_{n+1} = -1 \\ M & \omega_{n+1} = -1 \end{cases}$
 $\in_{\mathcal{M}} \left\{ \left(S_{n+1} \right) = \sum_{n} \left\{ \left(X_{n+1} S_n \right) \right\}$
 $indep h_{Mn} \mathcal{P} \left\{ \left(n S_n \right) + \mathcal{P} \left\{ \left(A S_n \right) \right\}$
 $= S_{nn} \int_{\mathcal{M}} \int_{\mathcal{M}} S_n = S_n \quad \text{is Markov}.$

laim : An is NOT Mankow [Intuition -> To find Any weld to know (Smt) depends on Sm Cont ditime Sm in time of Am.)



Definition 7.11. We say a <u>d</u>-dimensional process $Y = (Y^1, \dots, Y^d)$ process is a state process if for any security with maturity $(m) \leq N$, and payoff of the form $V_m = f_m(Y_m)$ for some (non-random) function f_m , the arbitrage free price must also be of the form $V_m = f_n(Y_n)$ for some (non-random) function f_n . $(M \leq M)$ Remark 7.12. For state processes given f_N , we find f_n by backward induction. The number of computations at time n is of order $\operatorname{Range}(Y_n).$ Remark 7.13. The fact that S_n is Markov (under \tilde{P}) implies that it is a state process. (Binned work $O < k < 1 + \gamma < k$) Raf neuk 7.13° Consider a see with payoff tw(SN) M = N - 1; $A F P = V_{N-1} = (1+N) F_{N-1} \delta_N(S_N)$ etate at time d_{-1} $M_{ankor} = \frac{1}{1+r}$ some for (S_{N-1}) (not history)