

Homework 4: Solutions

3.3.1 Let X denote the number of vertices of degree k in $G_{n,m}$. Then,

$$\begin{aligned} \mathbf{E}(X) &= n \frac{\binom{n-1}{k} \binom{N-n+1}{m-k}}{\binom{N}{m}} = \\ &= n \binom{n-1}{k} \frac{(N-n+1)(N-n)\cdots(N-n-m+k+2)m(m-1)\cdots(m-k+1)}{N(N-1)\cdots(N-m+1)} \approx \\ &= \frac{n^k m^k}{k! N^k} \prod_{i=0}^{m-k+1} \frac{N-n+1-i}{N-i} = \frac{n^k m^k}{k! N^k} \prod_{i=0}^{m-k+1} \left(1 - \frac{n-1}{N-i}\right) \approx \\ &= \frac{n^k m^k}{k! N^k} \exp\left\{-\sum_{i=0}^{m-k+1} \frac{n}{N-i}\right\} = \frac{n^k m^k}{k! N^k} \exp\left\{-\frac{n}{N} \sum_{i=0}^{m-k+1} \left(1 + O\left(\frac{i}{N}\right)\right)\right\} \approx \\ &= \frac{n^k m^k}{k! N^k} \exp\left\{-\frac{nm}{N}\right\} \approx \frac{d^k e^{-d}}{k!}. \end{aligned}$$

After this (a) we use the Chebyshev inequality and (b) see why we prefer $G_{n,p}$.

4.3.3 Fix $v \in [n]$. Let $\mathcal{A}_i, i = 0, \dots, i_0 = \lfloor \frac{2 \log n}{3 \log d} \rfloor$ be the event that $|S_i(v)| \in [(d/2)^i, (2d)^i]$ for all $v \in [n]$. Clearly \mathcal{A}_0 is true. Now if $p = m/N$ so that $d \approx np$ then in $G_{n,p}$,

$$\mathbf{P}(\neg \mathcal{A}_{i+1} \mid S_i(v), \mathcal{A}_j, j \leq i) = \mathbf{P}\left(\text{Bin}\left(n - \sum_{j=0}^i |S_j(v)|, 1 - (1-p)^{|S_i|}\right)\right).$$

Observe that $d^{i_0} \approx n^{2/3}$. Given the conditioning, $|\sum_{j=0}^i |S_j(v)|| = o(n)$ and $1 - (1-p)^{|S_i|} \approx |S_i|p$. Thus,

$$\mathbf{P}(\neg \mathcal{A}_{i+1} \mid S_i(v), \mathcal{A}_j, j \leq i) = \mathbf{P}(\text{Bin}((n - o(n)), (1 - o(1))|S_i|p)).$$

The expected value of the binomial in the above is $\approx d|S_i|$ and applying the Chernoff bounds we get that

$$\mathbf{P}(\neg \mathcal{A}_{i+1} \mid S_i(v), \mathcal{A}_j, j \leq i) \leq e^{-\Omega(d|S_i|)} \text{ and } d|S_i| \gg \log n.$$

This implies that

$$\mathbf{P}(\neg \mathcal{A}_{i+1} \mid \mathcal{A}_j, j \leq i) \leq n^{-\omega} \text{ where } \omega \rightarrow \infty.$$

But, then

$$\mathbf{P}(\neg \mathcal{A}_i) \leq \mathbf{P}(\neg \mathcal{A}_1) + \mathbf{P}(\neg \mathcal{A}_2 \mid \neg \mathcal{A}_1) + \mathbf{P}(\neg \mathcal{A}_3 \mid \mathcal{A}_1, \mathcal{A}_2) \cdots \leq in^{-\omega}.$$

and we can use the union bound to deal with all choices of v, i .

4.3.7 If $G_{n,n,p}$ is not connected then there is a component of size at most n .
So, assuming that $\omega = o(\log n)$,

$$\begin{aligned}
& \mathbf{P}(G_{n,n,p} \text{ is not connected}) \\
& \leq \sum_{k=1}^{n-1} \sum_{\ell=1}^{n-k} \binom{n}{k} \binom{n}{\ell} k^{\ell-1} \ell^{k-1} p^{k+\ell-1} (1-p)^{k(n-\ell)+\ell(n-k)} \\
& \leq \sum_{k=1}^{n-1} \sum_{\ell=1}^{n-k} \left(\frac{ne}{k}\right)^k \left(\frac{ne}{\ell}\right)^\ell k^{\ell-1} \ell^{k-1} p^{k+\ell-1} e^{-(k(n-\ell)+\ell(n-k))p} \\
& = \frac{1}{p} \sum_{k=1}^{n-1} \sum_{\ell=1}^{n-k} \left(\frac{k}{\ell}\right)^{\ell-k} \left(p \exp\left\{1 - \omega + \frac{2k\ell p}{k+\ell}\right\}\right)^{k+\ell} \\
& \leq \frac{1}{p} \sum_{k=1}^{n-1} \sum_{\ell=1}^{n-k} \left(\frac{k}{\ell}\right)^{\ell-k} \left(p \exp\left\{1 - \omega + \frac{(k+\ell)p}{2}\right\}\right)^{k+\ell} \\
& \leq \frac{1}{p} \sum_{s=2}^n s \left(p \exp\left\{1 - \omega + \frac{sp}{2}\right\}\right)^s \\
& = \frac{1}{p} \sum_{s=2}^n u_s.
\end{aligned}$$

Now

$$\frac{1}{p} \sum_{s=2}^{n/2} u_s \leq \frac{1}{p} \sum_{s=2}^{n/2} \left(pe^{1-\omega} n^{1/4+o(1)}\right)^s = O(n^{-1/2+o(1)}) = o(1).$$

$$\frac{1}{p} \sum_{s=n/2}^n u_s \leq \frac{1}{p} \sum_{s=n/2}^n \left(e^{1-\omega} n^{-1/2+o(1)}\right)^s = o(1).$$