# Department of Mathematical Sciences Carnegie Mellon University <br> 21-393 Operations Research II <br> Test 2 

Name: $\qquad$

| Problem | Points | Score |
| :--- | :--- | :--- |
| 1 | 50 |  |
| 2 | 30 |  |
| 3 | 20 |  |
| Total | 100 |  |

Q1: (50pts)
Find a minimum spanning tree in the following weighted graph.


Q2: (30pts) Let $\mathcal{W}$ denote the set of walks in a directed graph $D$. If $W_{1}$ is a walk from $a$ to $b$ and $W_{2}$ is a walk from $b$ to $c$ then $W_{1}+W_{2}$ is the walk from $a$ to $c$ obtained by following $W_{1}$ and then $W_{2}$.
Let $\ell: \mathcal{W} \rightarrow \mathbb{R}$ be a real valued function defined on $\mathcal{W}$. Suppose that it has the following properties:

1. $\ell(C) \geq 0$ for any closed walk $C$. (A walk is closed if it begins and ends at the same vertex).
2. If $W_{1}, W_{1}^{\prime}$ are walks from $a$ to $b$ and $W_{2}, W_{2}^{\prime}$ are walks from $b$ to $c$ and $\ell\left(W_{i}^{\prime}\right) \geq \ell\left(W_{i}\right)$ for $i=1,2$ then $\ell\left(W_{1}^{\prime}+W_{2}^{\prime}\right) \geq \ell\left(W_{1}+W_{2}\right)$.

Consider the following algorithm: $n$ is the number of vertices in $D$.
Initialise $W_{i, j}=(i, j)$ and $D_{i, j}=\ell\left(W_{i, j}\right)$ for $i, j=1,2, \ldots, n$.
For $k=1$ to $n$ Do
For $i=1$ to $n$ Do
For $j=1$ to $n$ Do
$D_{i, j} \leftarrow \min \left\{D_{i, j}, \ell\left(W_{i, k}+W_{k, j}\right)\right\}$
oD
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Prove that when the algorithm finishes,

$$
D_{i, j}=\min \{\ell(P): P \text { is a path from } i \text { to } j\}
$$

Q3: (20pts) Give an algorithm to solve the following scheduling problem. There are $n$ jobs labelled $1,2, \ldots, n$ that have to be processed one at a time on a single machine. There is an acyclic digraph $D=(V, A)$ such that if $(i, j) \in A$ then job $j$ cannot be started until job $i$ has been completed. The problem is to minimise $\max _{j} f_{j}\left(C_{j}\right)$ where for all $j, f_{j}$ is a monotone increasing. As usual, $C_{j}$ is the completion time of job $j$. This is distinct from its processing time $p_{j}$.

