# Department of Mathematical Sciences Carnegie Mellon University <br> 21-393 Operations Research II <br> Test 1 

Name:

| Problem | Points | Score |
| :--- | :--- | :--- |
| 1 | 30 |  |
| 2 | 30 |  |
| 3 | 40 |  |
| Total | 100 |  |

Q1: (30pts)
(a) Fill in the last column of the table below for solving the following knapsack problem:

$$
\begin{array}{ll}
\operatorname{maximise} & 3 x_{1}+7 x_{2}+17 x_{3} \\
\text { subject to } \\
& 2 x_{1}+3 x_{2}+6 x_{3} \leq 12
\end{array}
$$

$$
x_{1}, x_{2}, x_{3} \geq 0 \text { and integer. }
$$

What is the optimal solution?

| $w$ | $f_{1}\left(x_{1}\right)$ | $b_{1}$ | $f_{2}\left(x_{2}\right)$ | $b_{2}$ | $f_{3}\left(x_{3}\right)$ | $b_{3}$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 0 | 0 |  |  |
| 1 | 0 | 0 | 0 | 0 |  |  |
| 2 | 3 | 1 | 3 | 0 |  |  |
| 3 | 3 | 1 | 7 | 1 |  |  |
| 4 | 6 | 1 | 7 | 1 |  |  |
| 5 | 6 | 1 | 10 | 1 |  |  |
| 6 | 9 | 1 | 14 | 1 |  |  |
| 7 | 9 | 1 | 14 | 1 |  |  |
| 8 | 12 | 1 | 17 | 1 |  |  |
| 9 | 12 | 1 | 21 | 1 |  |  |
| 10 | 15 | 1 | 21 | 1 |  |  |
| 11 | 15 | 1 | 24 | 1 |  |  |
| 12 | 18 | 1 | 28 | 1 |  |  |

## Q2: (30pts)

A factory uses a single machine to manufacture two distinct products $A$ and $B$. It costs $c_{A}(x)$ to make $x$ units of $A$ and $c_{B}(x)$ to manufacture $x$ units of $B$. The demand for $A$ in period $j$ is $d_{j}(A)$ and the demand for $B$ in period $j$ is $d_{j}(B)$. The factory can store at most $H$ units altogether at any one time. Demand for $A$ must be met in the period that it occurs, either from inventory or from production that period. Demand for $B$ can be met in the period that it occurs, or in the following period.
Design a dynamic programming algorithm for finding the cheapest way of meeting demand for the next $n$ periods.

## Q3: (40pts)

(a) Find a minimum spanning tree in the following weighted graph.

(b) Find a path from vertex 1 to all other vertices of the digraph below that minimises the path function $\ell$. Here, if the edges have length $\ell(e), e \in E$ then a path $P=\left(e_{1}, e_{2}, \ldots, e_{k}\right)$ has length

$$
\ell(P)=\ell\left(e_{1}\right)+2 \ell\left(e_{2}\right)+3 \ell\left(e_{3}\right)+\cdots+k \ell\left(e_{k}\right) .
$$

In the figure below, an edge $(i, j)$ with $i<j$ is directed from $i$ to $j$ and its length is given in the middle of the edge.


