

10/1/14

Set Covering

Given $S_1, S_2, \dots, S_n \subseteq \{1, 2, \dots, m\} = [m]$

Each set S_j has a cost c_j

A cover $I \subseteq \{1, 2, \dots, n\}$ satisfies $\bigcup_{i \in I} S_i = [m]$

Cost $c(I) = \sum_{i \in I} c_i$. Problem: find minimum cost cover.

$$X_j = \begin{cases} 0 & \text{do not include } S_j \\ 1 & \text{include } S_j \end{cases}$$

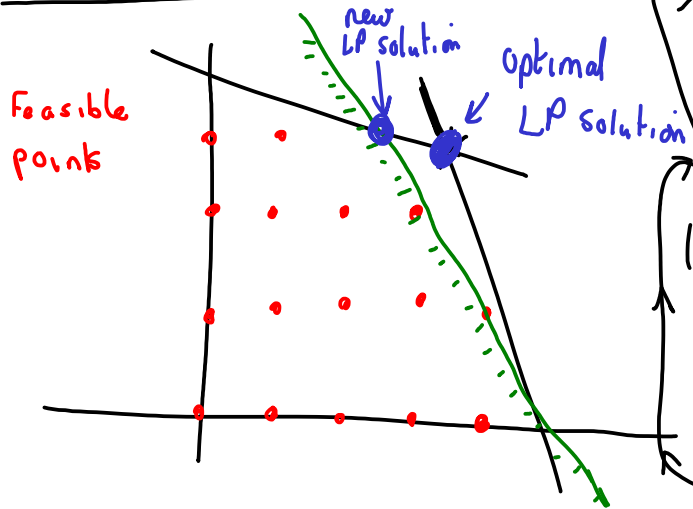
$$\text{Minimise cost} = c_1 x_1 + c_2 x_2 + \dots + c_n x_n$$

$$a_{ij} = \begin{cases} 0 & i \notin S_j \\ 1 & i \in S_j \end{cases}$$

$$a_{i1} x_1 + a_{i2} x_2 + \dots + a_{in} x_n \geq 1, \quad i=1, 2, \dots, m$$

$$x_j = 0 \sim \underline{1}$$

Cutting Plane Algorithm



- Idea:
- (i) Solve the LP relaxation
is ignore integrality
If solution is integer feasible
— done. IP is solved
 - (ii) If not, add a cut —
constraint that cuts off
LP solution, but no
integer solution.

Gomory Cuts: pure integer programs.

Suppose that we know

$$a_1x_1 + a_2x_2 + \dots + a_nx_n = b$$

$x_1, x_2, \dots, x_n$ are non-negative integers

→ this will be an equation from the optimal LP tableau. If we do not have an integer solution

$$a_1x_1 + a_2x_2 + \cdots + a_nx_n = b$$

Write $a_j = \lfloor a_j \rfloor + f_j$

↑
integer part

↑
fractional part

$$b = \lfloor b \rfloor + f$$

$$5\frac{1}{3} = 5 + \frac{1}{3}$$

$$-5\frac{1}{3} = -6 + \frac{2}{3}$$

$$\sum_{j=1}^n (a_j + f_j) x_j = b + f$$



$$\sum_{j=1}^n f_j x_j - f = b - \sum_{j=1}^n a_j x_j$$

Integer



Now remember that: $x_1, x_2, \dots, x_n \geq 0$ & integer.

LHS is an integer & $\geq -f > -1$ so LHS ≥ 0 .

Suppose that

$$\underbrace{x_i}_{\text{basic variable}} + \sum_{j \in J} b_{ij} x_j = b_{i0} \leftarrow \text{not integer}$$

non-basics

is the i th row of Simplex Tableau
Every feasible solution satisfies

$$\sum_{j \in J} f_j x_j \geq f > 0$$

Maximize

$$3x_1 + 4x_2$$

s.t.

$$x_1 + 2x_2 \leq 4$$

$$2x_1 + x_2 \leq 6$$

$$x_1, x_2 \geq 0 \text{ \& integer}$$

b_v	x_1	x_2	x_3	x_4	RHS
x_0	-3	-4			0
x_3	1	2	1		4
x_4	2	1		1	6