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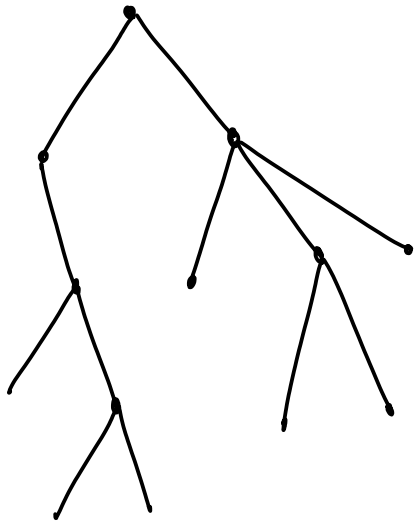
Suppose now that we must cut stick
into exactly k pieces



$$f_r(l) = \max_{x \leq l} [v_{x,l} + f_{r-1}(x)]$$

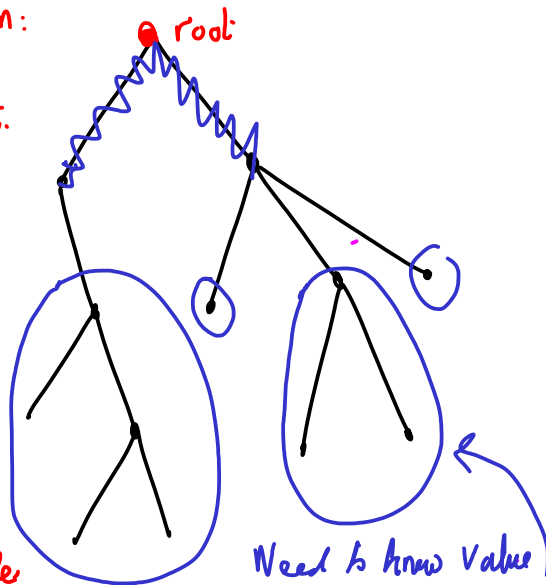
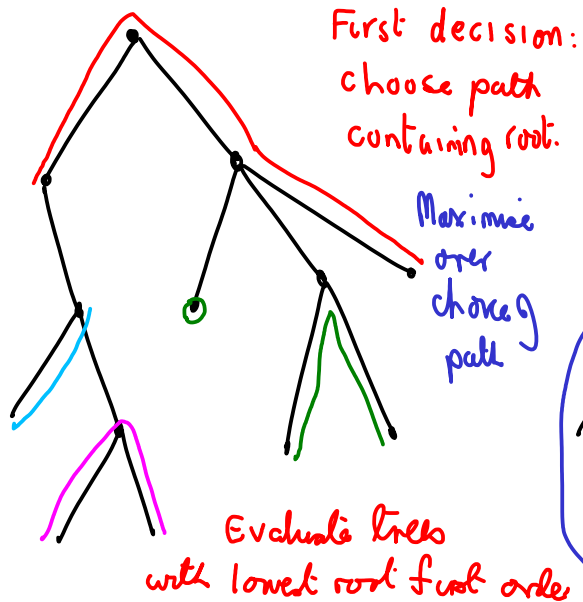
Annotations for the equation:

- Red arrows point to r and l in $f_r(l)$. Below them are the red constraints $1 \leq r \leq k$ and $0 \leq l \leq L$.
- A red arrow points to $x \leq l$ in the maximization term, with the text "place to cut" below it.



Each path in
tree has a value

Break up tree
into vertex disjoint
paths to maximise
total value.





$$f(l) = \max_x \left[v_{x,l} + f(x) \right]$$

Combinatorial Optimization:

- (i) Minimum length spanning tree
- (ii) Shortest path
- (iii) Assignment problem.

Minimum Spanning Tree

Given a ^{connected} graph G with edge weights

Problem: find spanning tree of minimum total weight

Kruskal's Algorithm:

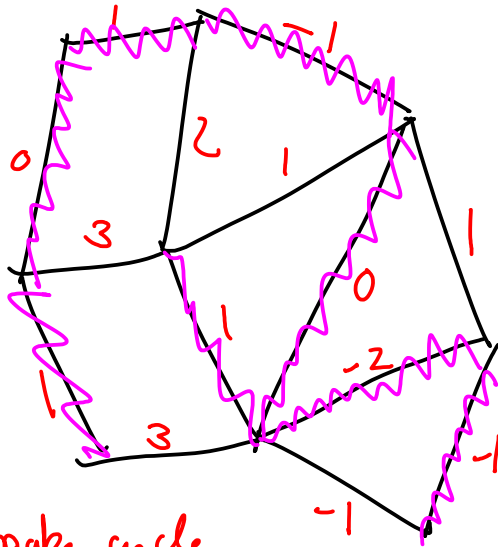
Greedy Algorithm:

Take edges of G in order

$e_1, e_2, \dots$

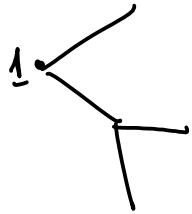
where $w(e_i) \leq w(e_{i+1})$

Include edge if does not make cycle.

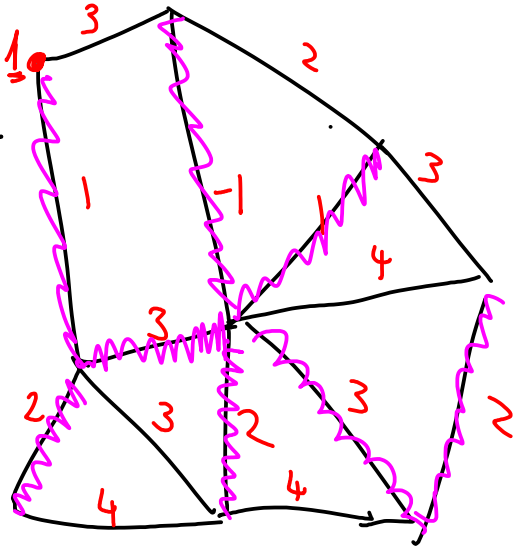


Dijkstra's Algorithm

Having constructed partial tree T



Added shortest edge from T to $\bar{T}$



More general algorithm: Having chosen

$\overline{F}_i = \{e_1, e_2, \dots, e_i\}$ = collection of disjoint
subtrees $T_1, T_2, \dots, T_k$

Choose any T_i and add the shortest edge
from T_i to $\overline{F}_i$.

