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Example of production problem.

$$H=3, n=5, c(x) = 18x - x^2.$$

$$d_i = 4, \forall i$$

$$f_r(i) = \min_x \left[c(x) + f_{r+1}(i+x-d_r) \right]$$

	f_1	x_1	f_2	x_2	f_3	x_3	f_4	x_4	f_5	x_5
0	244	4	188	7	150	4	112 110 104 * 94 82	4 5 6 7 2	56	4
1	233	3	183	6	139	3	101 101 97 * 89	3 4 5 6	45	3
2	220	2	176	5	126	2	94	5	32	2
3	205	1	167	1	111	1	89	1	17	1

Variations

① Add a smoothing cost $\sigma(x_1, x_2)$ for making x_1, x_2 in successive periods

e.g. $\sigma(x_1, x_2) = k(x_1 - x_2)^2$

$$f_i(i, y) = \min_x \left[c(x) + \sigma(y, x) + f_{i+1}(i+x-d, y_c) \right]$$

prod. last period

111) Machine replacement.

Suppose cost of producing x depends
on the age of the machine : $C(x, t)$
for a machine age t .

Cost of new machine is A .

Must acquire new machine once machine
reaches T , otherwise it is optional.

$$f_r(i, t) = \min_{\substack{\text{keep} \\ \text{replace}}} \begin{cases} \min_x (c(x, t) + f_{r+1}(i+x-d_r, t+1)) \\ A + \min_x (c(x, 0) + f_{r+1}(i+x-d_r, 1)) \end{cases}$$

↑
age of
machine

$$f_r(i, T) =$$

Knapsack problem

Scout X is going to camp.

Has a knapsack.

X can carry at most weight W

n possible items to pack. Each item of type j has weight w_j and value c_j .

Problem: choose items
to maximize
value

Integer program: Maximize $C_1X_1 + C_2X_2 + \dots + C_nX_n$
 $X_i = \# \text{ items of type } i \text{ that } X \text{ takes with them.}$
 $W_1X_1 + W_2X_2 + \dots + W_nX_n \leq W$
 $X_i \in \{0, 1, 2, \dots\}$

Dynamic Programming: $\text{Opt}[1, 2, 3, \dots, n; W] =$
Sequence of decisions
 $X_1, X_2, \dots$
 $\max_{X_1} [C_1X_1 + \text{Opt}[2, 3, \dots, n; W - W_1X_1]]$

$f_r(w)$ = max. obtainable using items of
type $1, 2, \dots, r$ and knapsack of
size w

$$= \max_{x_r} \left[c_r x_r + f_{r-1}(w - w_r x_r) \right]$$