

9/28/12

All pairs shortest path problem.

Given a matrix $D = \|d_{ij}\|$ where

d_{ij} = length of edge (i, j) .

Find a shortest ^{walk} path between every pair of vertices.

Assume there are no negative cycles

Floyd's Algorithm

For $k = 1$ to n do

For $i = 1$ to n do

For $j = 1$ to n do

$$D_{ij} := \min \{ D_{ij}, D_{ik} + D_{kj} \}$$

od

od

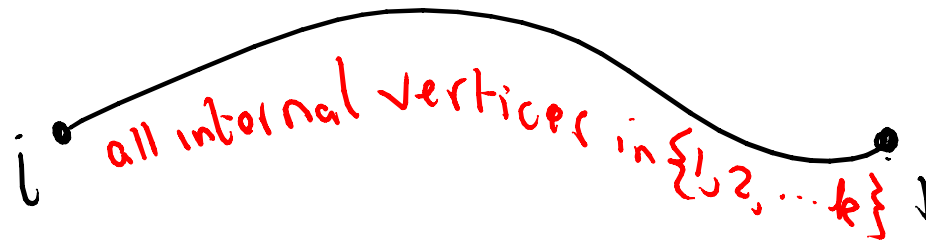
od

$O(n^3)$

Claim

At the end of round k ,

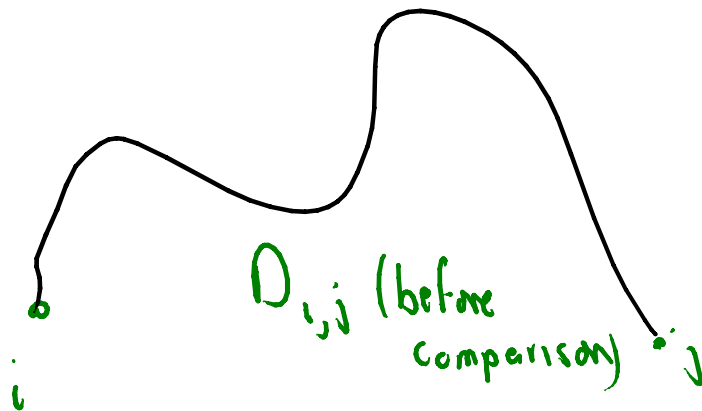
D_{ij} = min. length of walk from
 i to j



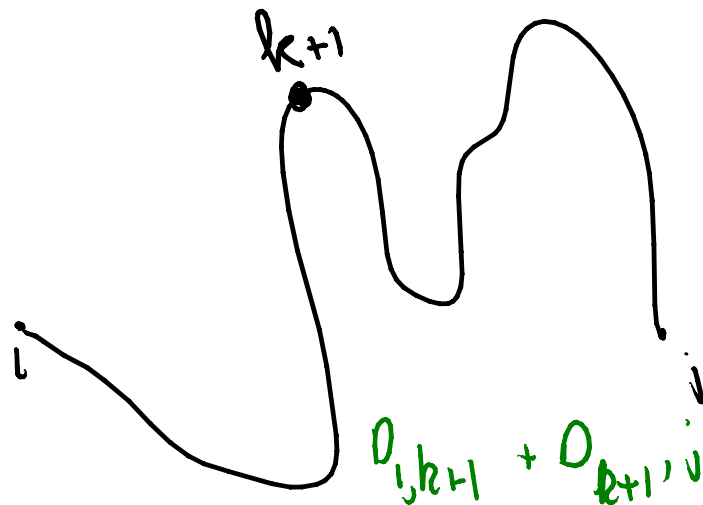
Proof by induction on k :

True for $k=0$

True for k : only uses internal vertices from
 $1, 2, \dots, k, k+1$



Don't use $k+1$
as an internal vertex

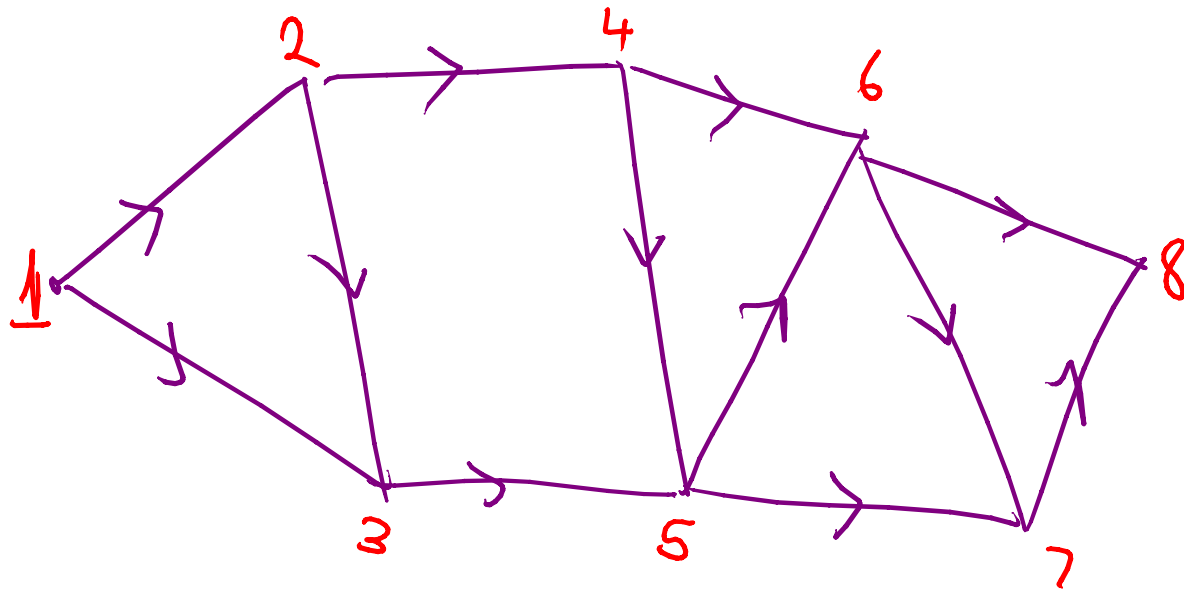


Digraphs without (oriented) cycles: DAG's

Directed Acyclic

Graphs

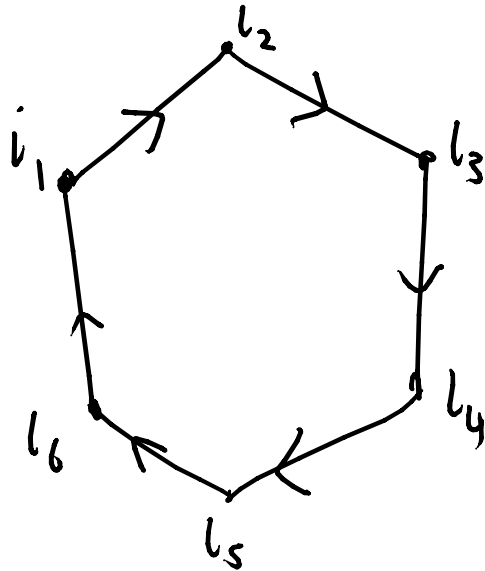
Linear (time) in # of arcs.



Topological ordering: # vertices $1, 2, \dots, n$
s.t. every arc $i \rightarrow j$ has $i < j$

Claim: a digraph has a topological ordering
iff it has no oriented cycles.

→
only if

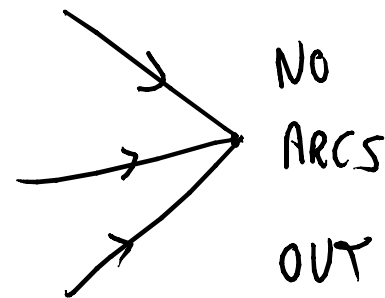


$$i_1 < i_2 < l_3 < \dots < l_6$$

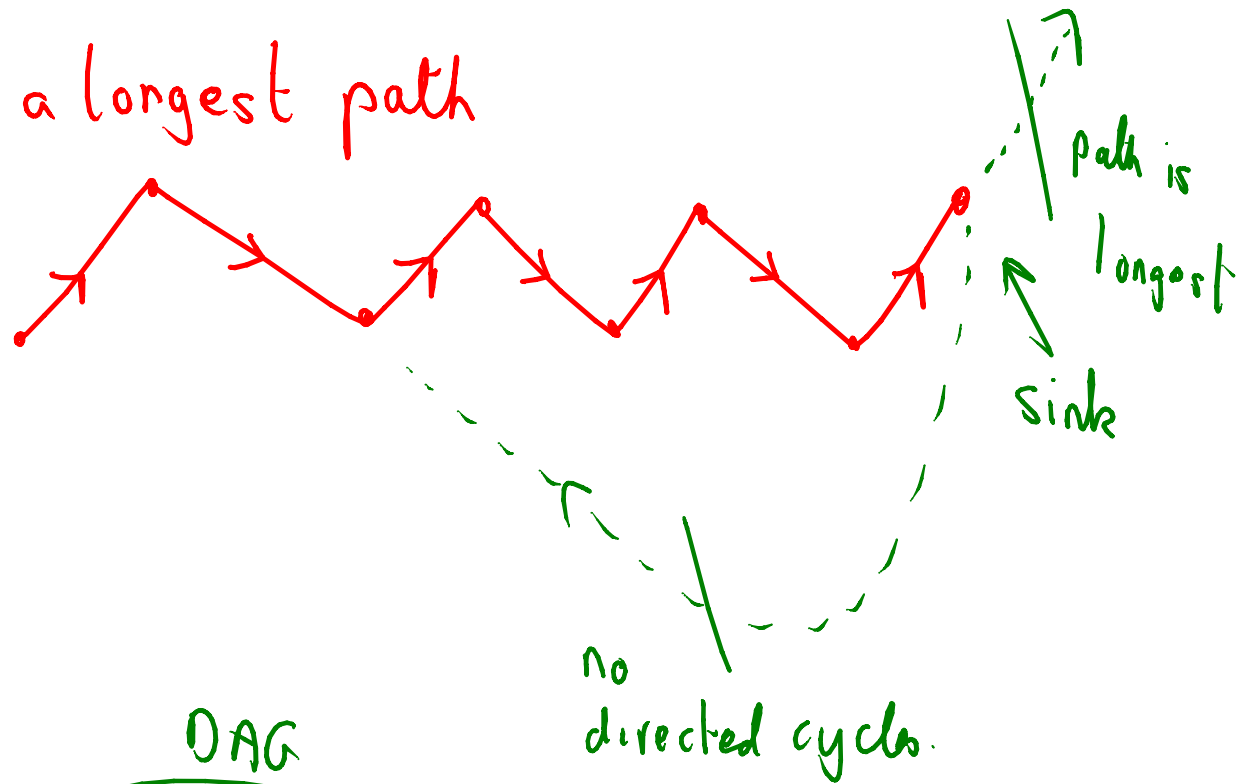
Contradiction

←
DAG $\Rightarrow$
 $\exists$ ordering

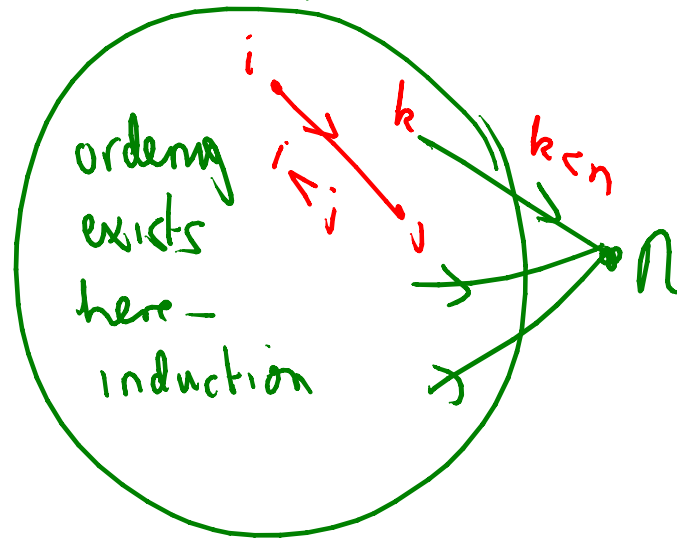
DAG $\Rightarrow \exists$ a "sink"



Take a longest path



DAG

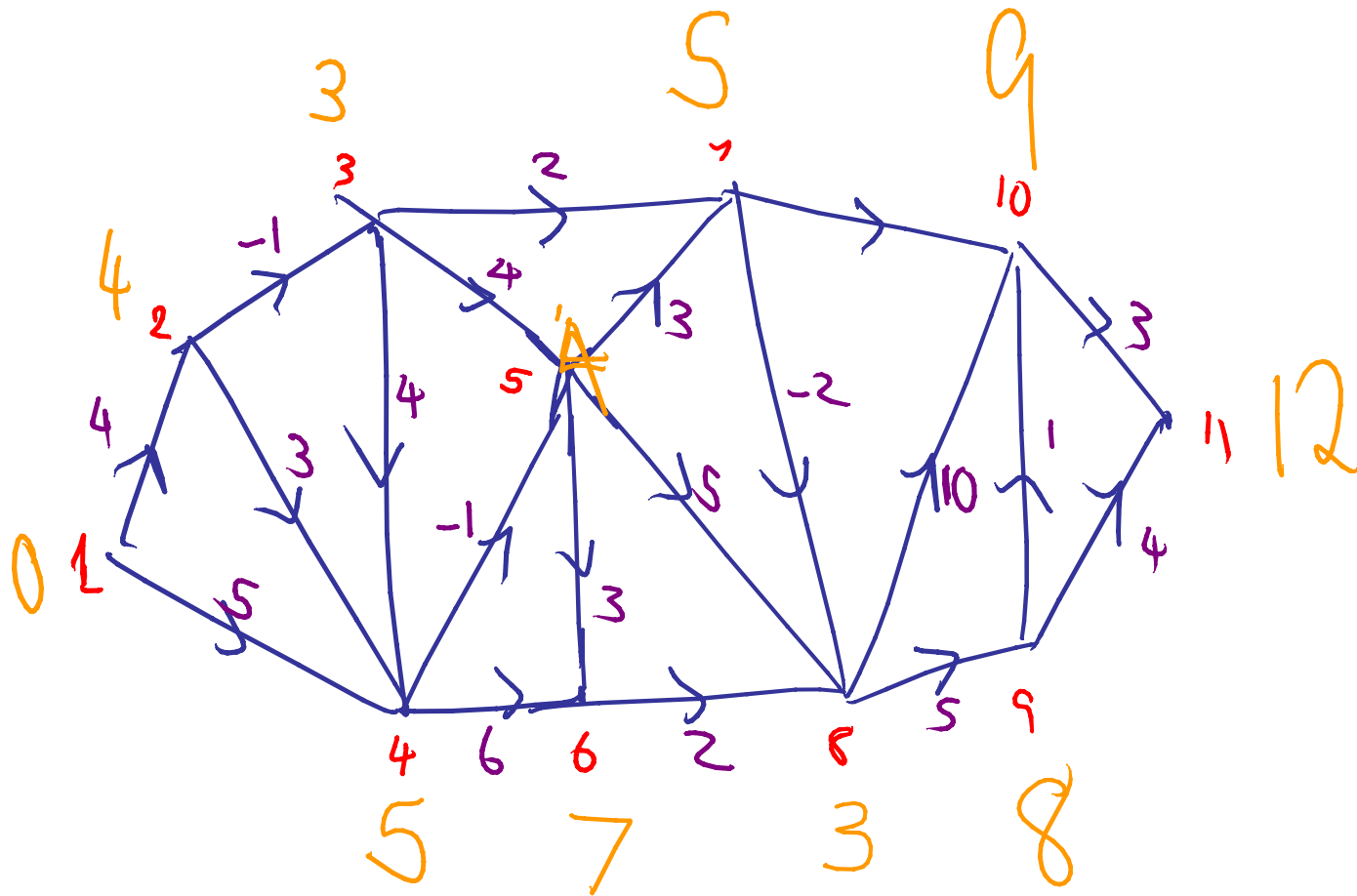


Suppose now that we want to find a shortest path from 1 to every other vertex.

d_i = min. path length to i

$$d_i = \min_{i < j} d_j + l_{ij}$$

$d_1 = 0$ and then I can compute $d_2, d_3, \dots, d_n$



Solve

$$d_j = \min d_i + l_{ij}$$

↑

⊕

↑

⊗

$$d_j = \oplus l_{ij} \otimes d_i$$

DAG
⇒

get a triangular system