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Variations on production problem:

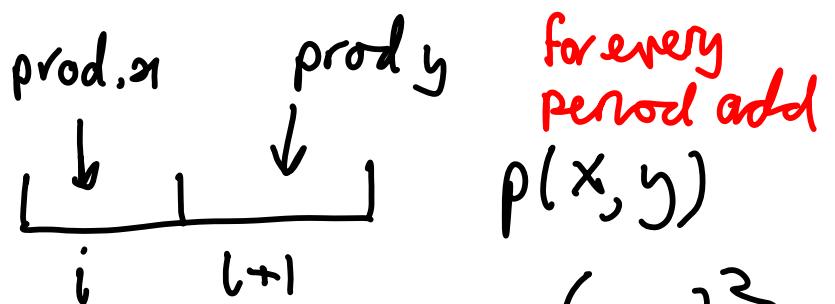


Demand:  $d_1, d_2, \dots, d_n$

Cost of production  $c(x)$

Max. inventory  $H$

Smoothing penalty:



e.g.  $(x - y)^2$

$$f_r(i, y) = \min_x \left[ c(x) + p(x, y) + f_{r+1}(i+n-d_j; x) \right]$$

## Machinery Replacement.

As a machine gets older, it costs more to produce:

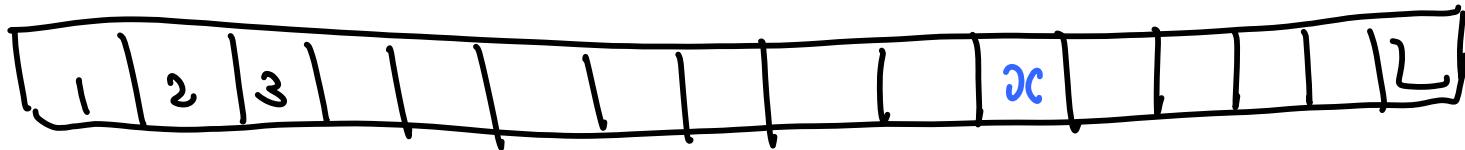
$c_t(x) =$  cost of producing  $x$  on a machine of age  $t$ . (Increase with  $t$ )

$$f_r(i) = \min_x [c_r(x) + f_{r+1}(i+r - d_r)]$$

Now suppose cost of buying a new machine is  $B$

$$f_r(i, t) = \min_x \left\{ \begin{array}{l} \min_x \left[ c_t(x) + f_{r+1}(i + \alpha - d_r, t+1) \right] \\ \min_x \left[ B + c_0(x) + f_{r+1}(i + \alpha - d_r, 1) \right] \end{array} \right.$$

LINGAR TOWN:



DIVIDE STRIP  $1..L$  INTO PIECES  
FOR DEVELOPMENT.

$v[i:j]$  = value of strip  $[i, i+1, \dots, j]$

Problem: find partition into pieces  
that maximizes total value  
 $f(l)$  = maximum obtainable from cutting up  $[1..l]$   
 $= \max_{0 \leq x < l} [v[x, l] + f(x-1)]$   $O(l^2)$   
time

Now suppose we want to make exactly  $k$  cuts.

$f(l, k)$  = max from  $l..k$  using  $k$  cuts

$$= \max_{0 \leq x < l} [v[x, l] + f(x-1, k-1)]$$

$$f(l, k) = -\infty \text{ if } k > l$$