

10/13/10

Shortest Path Problem

Digraph $D = (V, A)$

$l: A \rightarrow \mathbb{R}$ [edge a has length $l(a)$]

If P is a path

$$l(P) = \sum_{u \in P} l(u)$$

Problem:

Single Source Problem:

Given $s \in V$, find a shortest path from s to all $v \in V$.

All Pairs Problem:

Find a shortest path between every pair of vertices.

We can assume $l(u) \geq 0, \forall u \in A$

or

we can allow some arcs to have $l(u) < 0$.

① Non-negative arc lengths.

Dijkstra's Algorithm:

Initially $d(s) = 0$ & $d(v) = \infty$, $s \neq v$

$x = s$; $S = \{x\}$,

repeat

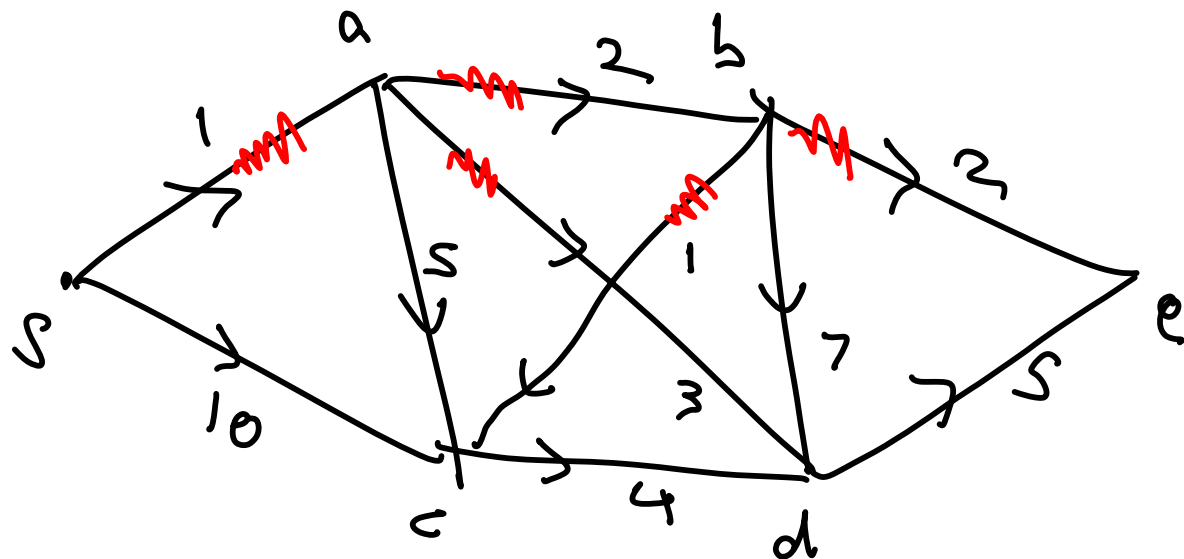
Process x ;

$\forall v: d(v) := \min(d(v), d(x) + l(x, v))$

$S := S \cup \{y\}$, $d(y) := \min_{u \notin S} d(u)$

until $x = y$
 $S = V$

S



	s	a	b	c	d	e
s	0	∞	∞	∞	∞	∞
a	0	1	∞	10	∞	∞
c	0	1	3	6	4	∞
d	0	1	3	4	4	5
e	0	1	3	4	4	5

Time Dependent Arc Lengths

Replace $l: A \rightarrow \mathbb{R}^+$

by $l: A \times \mathbb{R}^+ \rightarrow \mathbb{R}^+$

$u \in A$ then $l(u, t)$ = time taken to
cross u starting
at time t .

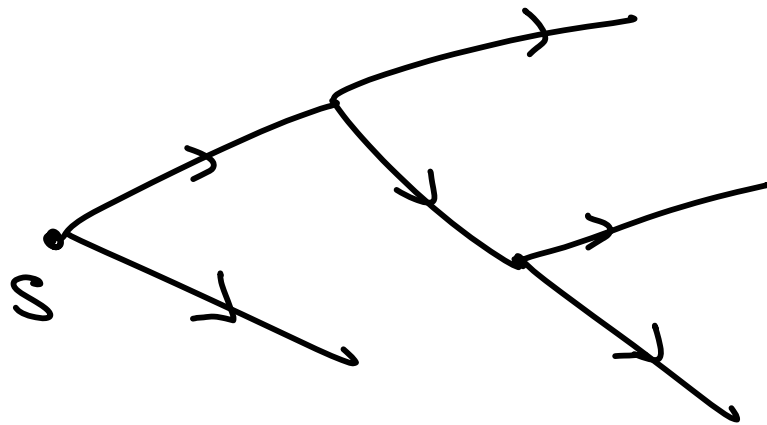
$P = \langle u_0, u_1, u_2, \dots, u_k \rangle$

$$l(P) = \underbrace{l(u_0, u_1)}_{\xi} + l(u_1, u_2, \xi) + \dots$$

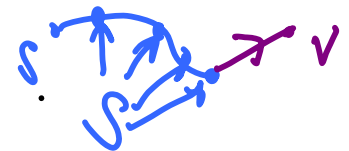
$$P = \{u_0, u_1, u_2, \dots, u_k\}$$

$$l(P) = \underbrace{l(u_0, u_1)}_{\xi_1}, 0 + \underbrace{l(u_1, u_2)}_{\xi_2}, \xi_1 + \dots$$

Dijkstra's Algorithm can be adapted to this situation:



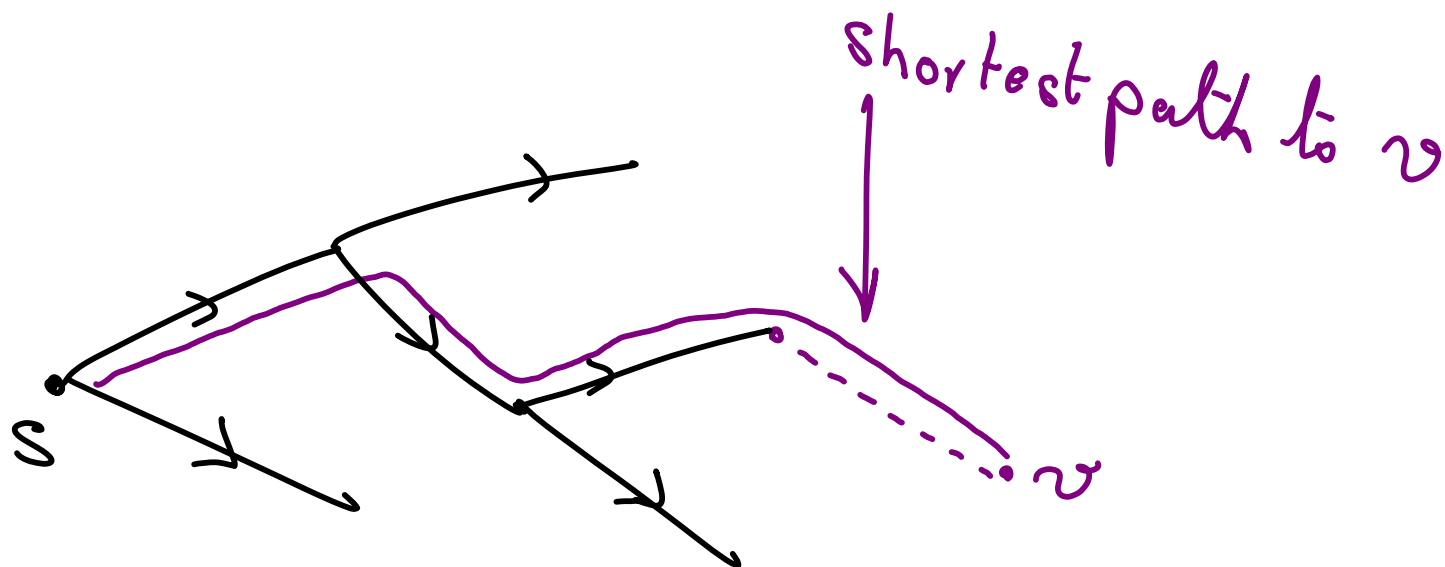
$d(v)$ = min
length
path
the from



S

claim: $v \in S \Rightarrow d(v)$ = length
shortest path from
 S to v

True initially



Any other path to v

