

9/18/09

$$\alpha = \frac{1}{2}$$

$$\begin{bmatrix} 3 & 7 & 4 & 2 \\ 1 & 5 & 6 & 3 \\ 2 & 8 & 4 & 7 \\ 5 & 9 & 3 & 1 \end{bmatrix}$$

$$\begin{array}{cccccc} i & 1 & 2 & 3 & 4 \\ \text{min} & 4 & 1 & 1 & 4 \end{array}$$

$$y_1 = 2 + \frac{1}{2} y_4 = 3$$

$$y_2 = 1 + \frac{1}{2} y_1 = 5/2$$

$$y_3 = 2 + \frac{1}{2} y_1 = 7/2$$

$$y_4 = 1 + \frac{1}{2} y_4 = 2$$

$$\alpha = \frac{1}{2}$$

$$\begin{bmatrix} 3 & 7 & 4 & 2 \\ 1 & 5 & 6 & 3 \\ 2 & 8 & 4 & 7 \\ 5 & 9 & 3 & 1 \end{bmatrix}$$

$y,$

y_2

y

y_4

3

$$5/2$$

$\frac{7}{2}$

2

$$i = 1 :$$

$$z + \frac{1}{2}y = 9/2$$

$$7 + \frac{1}{2}y_2 = 33/4$$

$$4 + \frac{1}{2} y_3 = 23/4$$

$$2^{\frac{1}{2}} g_4 = 3^*$$

$$i = 2 :$$

$$1 + \frac{1}{2} y_1 = 5/2 \quad *$$

$$5 + \frac{1}{2} y_2 =$$

6 + $\frac{1}{2} \times 5 = 6.5$

$$3 \cdot \frac{1}{2} y_4 =$$

$$\alpha = \frac{1}{2} \quad \begin{bmatrix} 3 & 7 & 4 & 2 \\ 1 & 5 & 6 & 3 \\ 2 & 8 & 4 & 7 \\ 5 & 9 & 3 & 1 \end{bmatrix} \quad \begin{matrix} y_1 & y_2 & y_3 & y_4 \\ 3 & 5/2 & 7/2 & 2 \end{matrix}$$

OPTIMAL

$i = 3:$

$$\begin{aligned} 2 + \frac{1}{2} y_1 &= 7/2 \quad * \\ 8 + \frac{1}{2} y_2 &= \\ 4 + \frac{1}{2} y_3 &= \\ 7 + \frac{1}{2} y_4 &= \end{aligned}$$

$i = 4:$

$$\begin{aligned} 5 + \frac{1}{2} y_1 &= \\ 9 + \frac{1}{2} y_2 &= \\ 3 + \frac{1}{2} y_3 &= \\ 1 + \frac{1}{2} y_4 &= 2 \quad * \end{aligned}$$

Theorem

π is optimal for all i
iff, $\forall i$

$$\textcircled{\times} y_i = c_{i, \pi(i)} + \alpha y_{\pi(i)} = \min_j [c_{ij} + \alpha y_j]$$

stopping condition for algorithm

Proof

Suppose that $(*)$ is true.

Let $\hat{\pi}$ be any other policy
and let $\hat{y}$ be its values.

$$\text{Let } \xi_i = \hat{y}_i - y_i$$

Want to show that $\xi_i \geq 0, \forall i$

$$\hat{y}_i = C_{i, \hat{\pi}(i)} + \alpha \hat{y}_{\hat{\pi}(i)}$$

$$y_i \leq C_{i, \hat{\pi}(i)} + \alpha y_{\hat{\pi}(i)}$$

optimality

⊗

$$x_i \geq \alpha \sum \hat{\pi}(i)$$

$$\geq \alpha^2 \sum \hat{\pi}^2(i)$$

$$\geq \alpha^k \sum \hat{\pi}^k(i) \rightarrow 0.$$

Now suppose that $\textcircled{\forall}$

fails:

Let $m(i) = \text{index of min in } \textcircled{\forall}$ i.e.

$$C_{i, m(i)} + \alpha y_{m(i)} = \min_j [C_{i, j} + \alpha y_j]$$

$m(i) = \pi(i)$ unless there is a strictly
better j .

$$U = \{j : m(i) \neq \pi(i)\} \neq \emptyset$$

Now define $\tilde{\pi}(i) = \pi(i)$
and let $\tilde{y}$ be its values.

$$y_i \geq c_{i, \tilde{\pi}(i)} + \alpha y_{\tilde{\pi}(i)}$$

$$\tilde{y}_i = c_{i, \tilde{\pi}(i)} + \alpha \tilde{y}_{\tilde{\pi}(i)}$$

$$D_i = y_i - \tilde{y}_i.$$

Show $D_i \geq 0, \forall i$ and $D_i > 0, i \in U$

As before

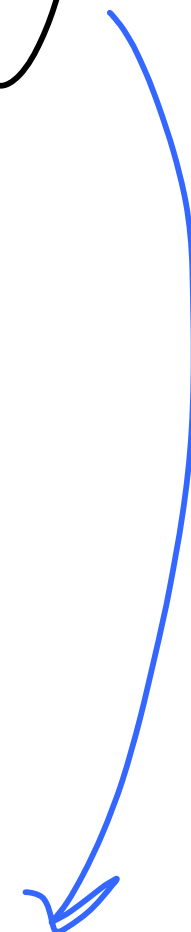
$$\eta_i \geq \alpha \eta_{\frac{\alpha}{1-\alpha}(i)}$$

and $\eta_i > \alpha \eta_{\frac{\alpha}{1-\alpha}(i)}, i \in U$

We know

$$\begin{aligned} \eta_i &\geq \alpha^k \eta_{\frac{\alpha^k}{1-\alpha^k}(i)} \quad \forall k \\ &\vdots \\ &\geq 0 \quad \text{in limit.} \end{aligned}$$

& $\eta_i > 0, \text{ for } i \in U$



Introduce randomness

Decisions:

D_i

$d \in D_i$

