

8/31/09

Knapsack Problem

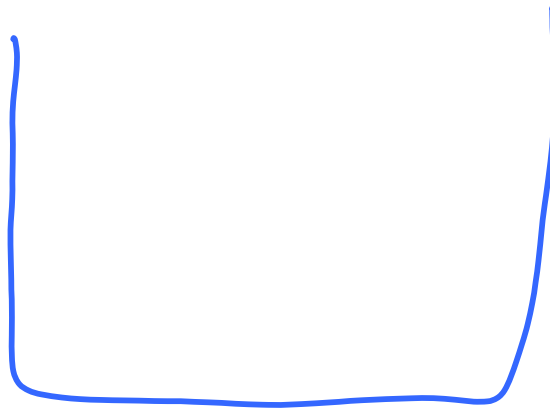
Maximize $C_1x_1 + C_2x_2 + \dots + C_nx_n$

Subject to $w_1x_1 + w_2x_2 + \dots + w_nx_n \leq W$

$x_1, x_2, \dots, x_n \geq 0$ & Integer

$C_1, \dots, C_n, w_1, \dots, w_n, W \geq 0$

Scout



Knapsack

n types of item

Each item have

value C_i

weight w_i

Pack x_1 of item 1

x_2 of item 2

$\vdots$

$$\text{Value} = C_1 x_1 + \dots + C_n x_n$$

$$\text{weight} = w_1 x_1 + \dots + w_n x_n$$

Scout can only carry W on trip

plane

We can solve this in $O(nW)$ time.

One way to solve this is "branch & bound"

We use Dynamic Programming here.

Solution = sequence of decisions.

$x_n = \# \text{ of type } n$

⋮

$$\text{Max Value} = \max_{x_n} \left[c_n x_n + \text{Value}_{n-1}(W - w_n x_n) \right]$$

$$f_r(w) = \text{maximum}_{w=0,1,\dots,W} \begin{array}{l} c_1 x_1 + \dots + c_r x_r \\ w_1 x_1 + \dots + w_r x_r \leq w \end{array}$$

$$= \max_{\substack{\alpha_r \geq 0 \\ \alpha_r \leq \frac{w}{w_r}}} [c_r x_r + f_{r-1}(w - w_r x_r)]$$

$$f_0(w) = 0 \quad \text{or} \quad f_1(w) = c_1 \left\lfloor \frac{w}{w_1} \right\rfloor$$

Run is $O(n W^2)$ time.

Maximize $3x_1 + 2x_2 + 4x_3$.

s.t.

$$2x_1 + 3x_2 + 7x_3 \leq 10$$

$$x_1, x_2, x_3 \geq 0 \text{ \& integer.}$$

ω	f_1	x_1	f_2	x_2	f_3	x_3
0	0	0	0	0	0	0
1	0	0	0	0	0	0
2	1	1	1	0	1	0
3	1	1	2	1	2	0
4	2	2	2	0 or 1	2	0
5	2	2	3	1	3	0
6	3	3	4	2	4	0
7	3	3	4	1 or 2	5	1
8	4	4	5	2	5	0 or 1
9	4	4	6	3	6	0 or 1
10	5	5	6	2 or 3	7	1

