

Ch5 Coupling

def: A coupling of μ and ν is a pair of random variables (X, Y) s.t. the marginal distribution of X is μ and of Y is ν .

Note: recall $\|\mu - \nu\|_{TV} = \min_{(X,Y) \text{ a coupling of } \mu, \nu} \Pr(X \neq Y)$ (Prop 4.7)

Ex: Simple random walk on $\{0, 1, \dots, n\}$ (we go ± 1 unless we are at 0 or n)

We suspect that $x \leq y \Rightarrow P^t(x, n) \leq P^t(y, n)$

How to prove? Coupling!

Define two walks X_t, Y_t , $X_0 = x, Y_0 = y$.

Such that they always both go same direction at each step (if possible).

Notes: • X_t and Y_t almost surely meet up and stay together

• (X_t, Y_t) is a coupling of the distributions $P^t(x, \cdot)$ and $P^t(y, \cdot)$

• $x \leq y \Rightarrow X_t \leq Y_t \quad \forall t$

Therefore $P^t(x, n) = \Pr(X_t = n) \leq \Pr(Y_t = n) = P^t(y, n)$ $\square$

def: A coupling of Markov chains w/ trans matrix P is a pair (X_t, Y_t) s.t. both (X_t) and (Y_t) are MC's with matrix P .

Notation: if (X_t, Y_t) is a coupling of MC's with P then we write $\text{and } X_0 \geq x, Y_0 \geq y$

$\Pr_{x,y}$ for probabilities on the common space.

Thm: let (X_t, Y_t) be a coupling such that

$$\star X_s = Y_s \Rightarrow X_{s+1} = Y_{s+1}$$

and $X_0 = x, Y_0 = y$. define $\tau_{\text{couple}} := \min \{t : X_t = Y_t\}$

$$\text{then } \|P^t(x, \cdot) - P^t(y, \cdot)\|_{TV} \leq \Pr_{x,y}(\tau_{\text{couple}} > t).$$

Pf: Note $P^t(x, z) = \Pr_{x,y}(X_t = z)$, $P^t(y, z) = \Pr_{x,y}(Y_t = z)$

so (X_t, Y_t) is a coupling of $P^t(x, \cdot)$ and $P^t(y, \cdot)$

$$\text{Now } \|P^t(x, \cdot) - P^t(y, \cdot)\|_{TV} \leq \Pr_{x,y}(X_t \neq Y_t) = \Pr_{x,y}(\tau_{\text{couple}} > t)$$



def: Markovian coupling, p. 65

$$\text{Cor: } d(t) \leq \max_{x,y \in \Omega} \Pr_{x,y}(\tau_{\text{couple}} > t)$$

$$\text{Pf: } d(t) \leq \bar{d}(t) = \max_{x,y} \|P^t(x, \cdot) - P^t(y, \cdot)\|_{TV} \leq \text{RHS}$$

Ex: Mixing time for lazy random walk on cycle:

Use the coupling (X_t, Y_t) with $X_0 = x, Y_0 = y$

where they never move simultaneously (i.e. either X moves or Y moves)

And let D_t be the clockwise distance from X_t to Y_t .

X_t meets Y_t when $D_t = 0$ or n and then ~~run~~ run X_t and Y_t together

Note that D_t is a random walk (nonlazy) on $\{0, \dots, n\}$

and
$$d(t) \leq \max_{x,y} \Pr_{x,y}(\tau_{\text{couple}} > t) \leq \frac{\max_{x,y} E_{x,y}(\tau_{\text{couple}})}{t}$$

but $E_{x,y}(\tau_{\text{couple}}) = k(n-k)$ where $D_0 = k$

so $d(t) \leq \frac{n^2}{4t}$ and $t_{\text{mix}} \leq n^2$

Ex: random ^{lazy} walk on torus (describe torus, p 65)

Thm: for the lru on $\mathbb{Z}_n^d$, $t_{\text{mix}}(\epsilon) \leq C(d) n^2 \log_2(\epsilon^{-1})$
 where $C(d)$ is a constant depending on d .

PF (by coupling): Lets couple. $\vec{X}_0 = \vec{x}$, $\vec{Y}_0 = \vec{y}$.

To move: randomly pick a coordinate (of the d dimensions)

If $\vec{X}_t$ and $\vec{Y}_t$ agree in this coordinate, move ^(or stay) together.

If not, flip a coin to see whether $\vec{X}$ moves or $\vec{Y}$ moves
 and flip again to see which direction.

define $\tau_i := \min \{ t \geq 0 : X_t^i = Y_t^i \}$

Coordinate i is selected with probability $\frac{1}{d}$, so the expected wait time between picking i again is d .

Ex 5.3: if $X_1, X_2, \dots$ are ind. with mean μ and τ is $\mathbb{Z}^+$ -valued independent of the X 's, then $E\left(\sum_{i=1}^{\tau} X_i\right) = \mu E(\tau)$

$\Rightarrow E_{x,y}(\tau_i) \leq \frac{dn^2}{4}$

Now $E_{x,y}(\tau_{\text{couple}}) = E_{x,y}\left(\max_i \tau_i\right) \leq \frac{d^2 n^2}{4}$

Now $\Pr_{X,Y}(\tau_{\text{couple}} > t) \leq \frac{\mathbb{E}_{X,Y}(\tau_{\text{couple}})}{t} \leq \frac{d^2 n^2}{4t} \Rightarrow t_{\text{mix}} \leq d^2 n^2$

so $t_{\text{mix}}(\epsilon) \leq \lceil \log_2 \epsilon^{-1} \rceil t_{\text{mix}} \quad (4.36) \quad \square$

LRW on hypercube $\mathbb{Z}_2^n$

Let $\vec{X}_0 = \vec{X}$, $\vec{Y}_0 = \vec{Y}$ and couple as follows:

Pick a random coordinate and "refresh" both X and Y in that coordinate by replacing both with a 0 or 1, chosen by a coin toss.

$$\tau_{\text{couple}} \leq \min \{ t \geq 0 : \text{all coordinates have been picked} \}$$

Now recall Section 2.2, coupon collector problem.

$$d(n \log n + cn) \leq \Pr(\text{there are still any unpicked coordinates at step } n \log n + cn) \leq e^{-c}$$

this immediately gives $t_{\text{mix}}(\epsilon) \leq n \log n + \log\left(\frac{1}{\epsilon}\right)n$

distance between $P^t(x, \cdot)$ and $P^{t+1}(x, \cdot)$ for lazy P .

Prop: let Q be irreducible trans. matrix and $P = \frac{1}{2}(I+Q)$.

$$\text{then } \|P^t(x, \cdot) - P^{t+1}(x, \cdot)\|_{TV} \leq \frac{12}{\sqrt{t}}$$

PF: we construct $(X_t), (Y_t)$ both started at x
with trans matrix P s.t. $P_t(X_t \neq Y_{t+1}) \leq \frac{12}{\sqrt{t}}$
and this suffices since (X_t, Y_{t+1}) is a coupling of
the distributions $P^t(x, \cdot)$ and $P^{t+1}(x, \cdot)$.

Let (Z_t) be a MC with matrix Q started at x ,
and (W_t) an i.i.d sequence of 0-1 valued random variables.
Let $N_t := \sum_{s=1}^t W_s$ and $Y_t = Z_{N_t}$.

Note that Y_t is an MC with matrix P .

$$\text{Let } X_t := \begin{cases} Z_{t-(N_{t+1}-W_1)} & \text{if } X_{t-1} \neq Y_t \\ Y_{t+1} & \text{otherwise} \end{cases}$$

X_t is also a MC with matrix P .

$$\text{Let } \tau := \min \{t \geq 0 : X_t = Y_{t+1}\}.$$

Case 1: $W_1 = 0$: then $Y_1 = X_0 = x$ so $\tau = 0$

Case 2: $W_1 = 1$ then $\tau \leq \min \{t \geq 0 : N_{t+1} = t - (N_{t+1} - W_1)\}$

$$\text{and } N_{t+1} = t - (N_{t+1} - W_1) \Leftrightarrow 2N_{t+1} - t = 1$$

but $(2N_{t+1} - t)_{t=0}^{\infty}$ is a random walk starting at 2

but th 2.17 $\Rightarrow \Pr(\tau > t) \leq 12/\sqrt{t}$

Grand Coupling

def: a collection $\{X_t^x : x \in \Omega, t=0,1,\dots\}$ of random variables

Such that for each $x \in \Omega$, (X_t^x) is a MC with matrix P started at x is called a Grand Coupling.

Note: can construct using random mapping model:

let $f: \Omega \times \Lambda \rightarrow \Omega$ and Z a Λ -valued random var

such that $P(x,y) = \Pr(f(x,Z)=y)$. (f,Z) exists by Prop 15

Now let $Z_1, Z_2, \dots$ be i.i.d copies of Z and

$$X_0^x := x, \quad X_t^x = f(X_{t-1}^x, Z_t) \quad \text{for } t \geq 1$$

~~Thm~~ Recall the Metropolis chain on colorings.

$\tilde{\Omega}$ = space of colorings, improper allowed

Thm: ^{Let} $|G| = n$, max degree Δ . If $q > 3\Delta$ and

$C_{\text{met}}(\Delta, q) = 1 - \frac{3\Delta}{q}$ then for the Metropolis chain on G

$$t_{\text{mix}}(\epsilon) \leq C_{\text{met}}(\Delta, q)^{-1} n \left[\log n + \log \frac{1}{\epsilon} \right]$$

Pf: use a grand coupling. Keep track of the MC for every possible start coloring in $\tilde{\Omega}$.

~~Def~~ def: for $x, y \in \tilde{\Omega}$, $\rho(x, y) := \sum_{v \in V} \mathbb{1}_{\{x(v) \neq y(v)\}}$

Note: ~~ρ~~ ρ is a metric.

Suppose $\rho(x, y) = 1$ so x and y disagree only at v_0 .

Let $N = x(N(v_0))$, the set of colors of neighbors of v_0

$$\text{Then } \Pr[\rho(x_i^x, x_i^y) = 0] = \frac{1}{n} \cdot \frac{q - |N|}{q} \geq \frac{q - \Delta}{nq}$$

There is a case where $\rho(x_i^x, x_i^y) = 2$ (see p. 72)

$$\text{but } \Pr[\rho(x_i^x, x_i^y) = 2] \leq \frac{\Delta}{n} \cdot \frac{2}{q}$$

$$\text{Thus } \mathbb{E}[\rho(x_i^x, x_i^y) - 1] \leq \frac{3\Delta - q}{nq}$$

$$\text{So } \mathbb{E}[\rho(x_i^x, x_i^y)] \leq 1 - \frac{C_{\text{met}}(\Delta, q)}{n}$$

Now if $\rho(x, y) = r$ then $\exists$ a sequence

$$X = x_0, x_1, \dots, x_r = y \quad \text{s.t.} \quad \rho(x_i, x_{i+1}) = 1$$

$$\text{SO } \Delta\text{-ineq} \Rightarrow \mathbb{E}[\rho(x^x, x^y)] \leq \mathbb{E}\left[\sum_{k=1}^r \rho(x_i^{x_k}, x_i^{x_{k+1}})\right] \\ \leq \rho(x, y) \cdot \left(1 - \frac{c_{\text{net}}}{n}\right)$$

$$\text{Now } \mathbb{E}[\rho(x^x, x^y) \mid x_{t+1}^x = x_{t+1}, x_{t+1}^y = y_{t+1}] = \mathbb{E}[\rho(x_i^{x_{t+1}}, x_i^{y_{t+1}})] \\ \leq \rho(x_{t+1}, y_{t+1}) \left(1 - \frac{c_{\text{net}}}{n}\right)$$

$$\text{SO } \mathbb{E}[\rho(x^x, x^y)] \leq \mathbb{E}[\rho(x_{t+1}^x, x_{t+1}^y)] \left(1 - \frac{c_{\text{net}}}{n}\right) \\ \leq \dots \leq \rho(x, y) \left(1 - \frac{c_{\text{net}}}{n}\right)^t$$

$$\text{Now } \Pr[X_r^x \neq X_r^y] = \Pr[\rho(x^x, x^y) \geq 1] \leq \rho(x, y) \left(1 - \frac{c_{\text{net}}}{n}\right)^t \\ \leq n e^{-\frac{c_{\text{net}}}{n} t}$$

thus $d(t) \leq n e^{-\frac{c_{\text{net}}}{n} t}$ and the result follows $\square$