

9/9/15

Probabilistic Method

Coloring problem: A_1, A_2, \dots, A_n
are subsets of A , $|A_i| = k$,
 $i = 1, 2, \dots, n$.

If $n < 2^{k-1}$ then there exists
a partition of $A = \underset{\text{red}}{R} \cup \underset{\text{blue}}{B}$

such that

$$A_i \cap R \neq \emptyset \text{ \& } A_i \cap B \neq \emptyset$$

for $i = 1, 2, \dots, n$

Note: - if n is too large then it may not be possible to find such a coloring.

Suppose $n = \binom{2k}{k}$ and $\{A_1, A_2, \dots, A_n\} = \left(\begin{matrix} [2k] \\ k \end{matrix} \right)$.

In any "coloring" there is a set R or B of size $\geq k$. Suppose R is a k -subset of R - is a red.

The largest possible value of n (in terms of k) is not

How can we find a "good" coloring?

We choose R, B uniformly at random and show that with positive probability we get a "good" coloring.

$$\text{BAD}(i) = \{A_i \subseteq R \text{ or } A_i \subseteq B\}.$$

$$\text{BAD} = \bigcup_{i=1}^n \text{BAD}(i)$$

Want to show that

$$P(\text{BAD}) < 1$$

$\Rightarrow \exists$ a coloring with desired property.

If there were no coloring then we would have $P(\text{BAD}) = 1$



$$P(\text{BAD}) = P\left[\bigcup_{i=1}^n \text{BAD}(i)\right]$$

$$\leq \sum_{i=1}^n P[\text{BAD}(i)]$$

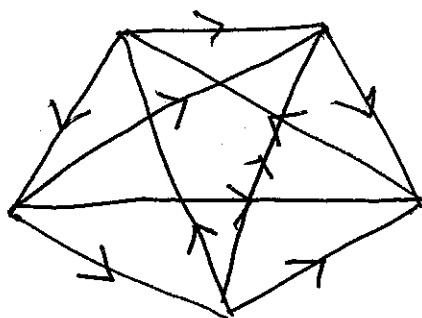
$$= \sum_{L=1}^n \left(\frac{1}{2}\right)^{k-1}$$

$$= \frac{n}{2^{k-1}} < 1$$



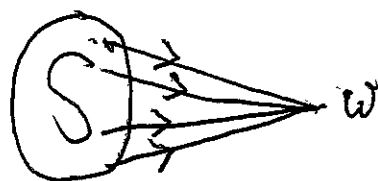
Tournaments

A tournament is an orientation of a complete graph K_n .

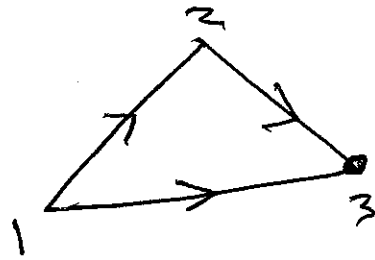


For $1 \leq i < j \leq n$
there is an edge
 (i, j) or (j, i) but
not both

Let \mathcal{P}_k be the property that
for every set $S \subseteq [n]$, $|S| = k$
there exists $w \notin S$ such that
 w "beats" S i.e. all edges are ^{oriented} from
 S to w .

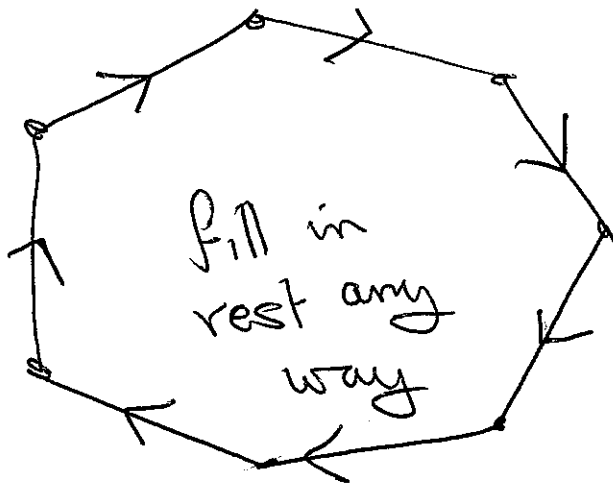


Not every tournament will have property \mathcal{A}_1 , say.



$S = \{3\}$ breaks \mathcal{A}_1 .

But there exist tournaments with \mathcal{A}_1



Theorems

For every k there exists a tournament with property \mathcal{A}_k .

To prove there exist a tournament with property \mathcal{P}_k , we choose n large and choose our tournament T at random i.e. we independently orient the edges in either direction.

Bad events: ~~for~~ for $|S| = k$ we let

$$\mathcal{B}_S = \left\{ \nexists w \text{ that beats } S \right\}.$$

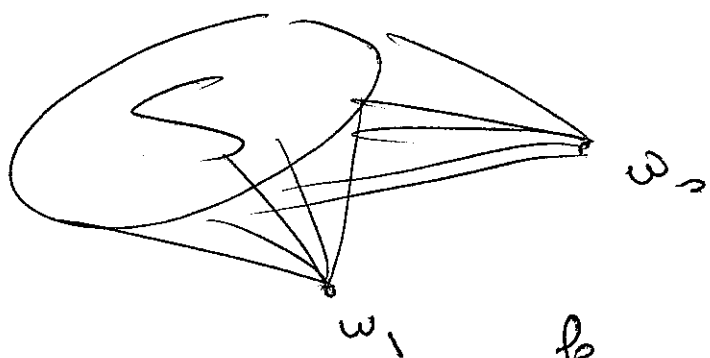
$$\overline{\mathcal{P}_k} = \bigcup_{|S|=k} \mathcal{B}_S$$

$$P(\overline{\mathcal{P}_k}) \leq \sum_{|S|=k} P(\mathcal{B}_S)$$

$$\mathcal{B}_S = \bigcap_{w \notin S} \mathcal{B}_{S,w}$$

where $\mathcal{B}_{S,w} = \{ w \text{ does not beat } S \}$

The events $\mathcal{B}_{S,w}$ are independent for a fixed S .



\mathcal{B}_{S,w_1} & \mathcal{B}_{S,w_2}
depend on
disjoint "coin flips".

$$P_i(\neg \mathcal{B}_{S,w}) = \left(\frac{1}{2}\right)^k$$

$$P_i(\mathcal{B}_{S,w}) = 1 - \left(\frac{1}{2}\right)^k$$

$$P_i(\mathcal{B}_S) = \left(1 - \left(\frac{1}{2}\right)^k\right)^{n-k}$$

$$P(\text{TO}_{k,r}) \approx \sum_{S} P(\mathcal{B}_S)$$

$$\approx \sum_{S} \left(1 - \frac{1}{2^k}\right)^{n-k}$$

$$\approx \binom{n}{k} \left(1 - \frac{1}{2^k}\right)^{n-k}$$

$$< 1 \quad \text{if} \quad n \geq k.$$