

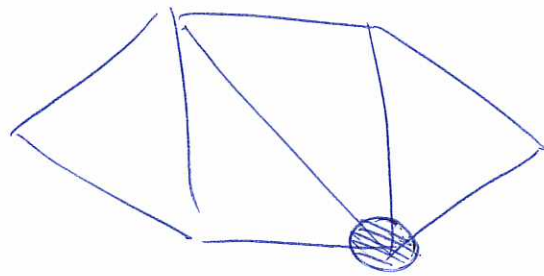
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## Geography

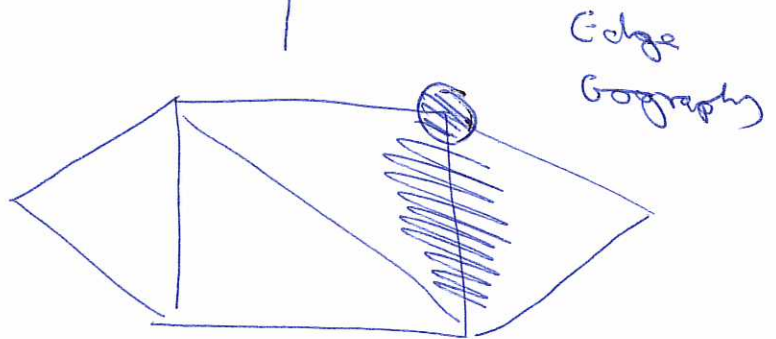
Start with a chip on a graph vertex  $x$  and move the chip to a neighbor  $y$ .

Vertex Geography: delete  $x$

Edge Geography: delete  $y$



Move



Edge Geography

Graph can be directed or undirected.

Winner = person who makes the last move.

Complexity: (Can you tell "quickly"  $N$  from  $P$ )

Digraph: Edge & vertex geography are PSPACE-hard.

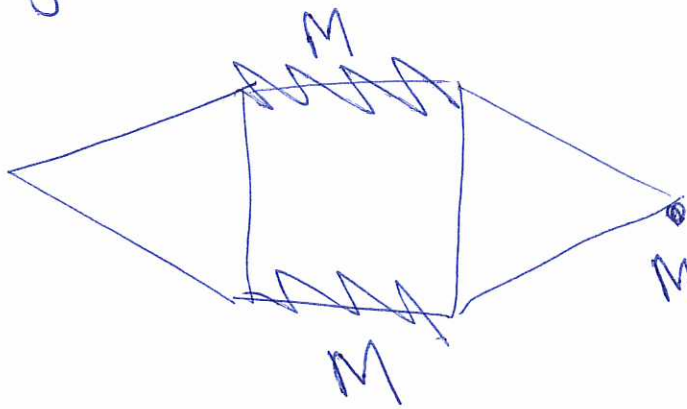
Graphs: Edge geography is PSPACE-hard (

Vertex geography is in  $P$

Edge geography on a hypertite graph is in  $P$ ]

# Vertex Geograph in (undirected) graphs

Need some simple observations about matchings.



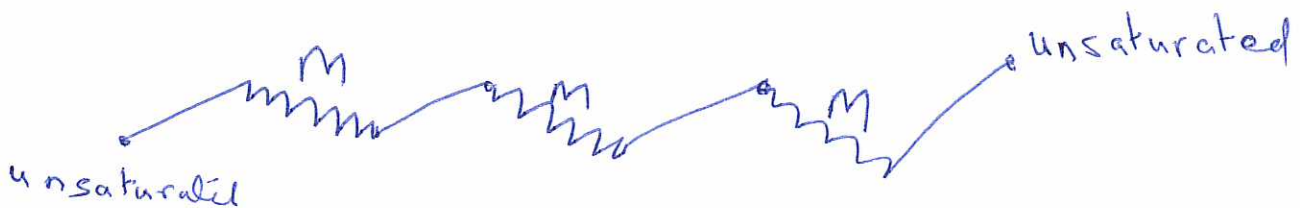
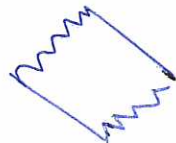
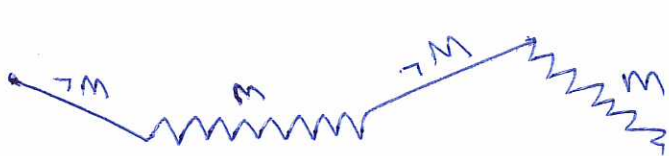
Matching  $M =$   
 $\{ \text{disjoint edges} \}$

M-unsaturated

$M$  is perfect - if no vertex is  $M$ -unsaturated

## Alternating Paths

A path is  $M$ -alternating if the edges are  
 $M, \neg M, M, \neg M, \dots$  etc



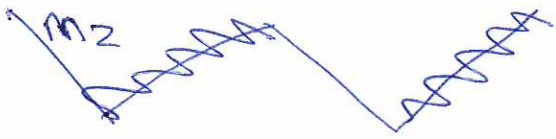
# Lemma

$M$  is a maximum size matching iff there are no  $M$ -augmenting paths.

## Proof

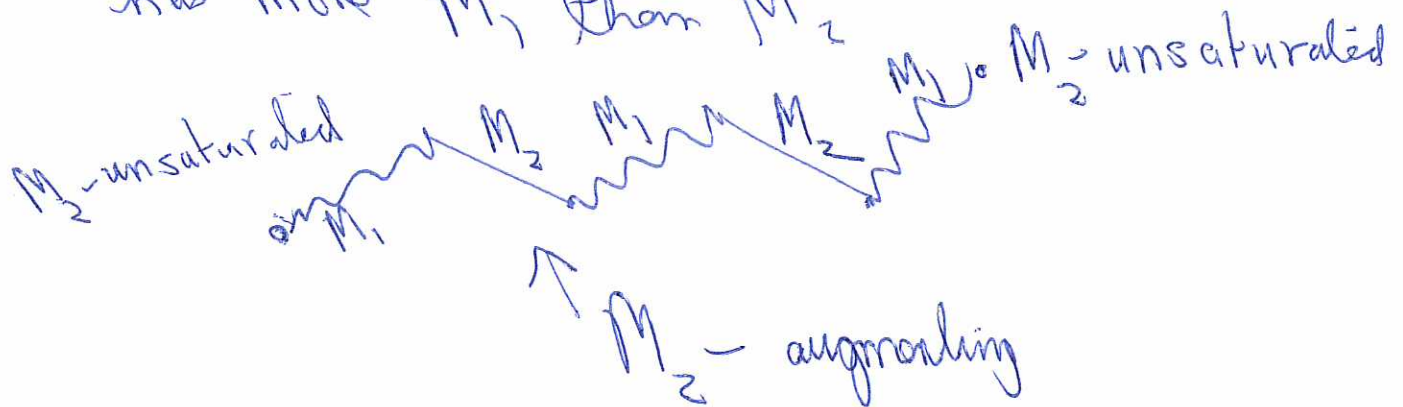
If there is an augmenting path we remove its  $M$ -edge from  $M$  and add the  $\neg M$  edges to get a matching one size bigger.

Suppose  $M_1$  &  $M_2$  are matchings &  $|M_1| > |M_2|$

$M_1 \oplus M_2$ :  = alternating paths & cycle.

$M_1 \setminus M_2 \cup M_2 \setminus M_1$

If  $|M_1| > |M_2|$  then at least one path has more  $M_1$  than  $M_2$



# Geography

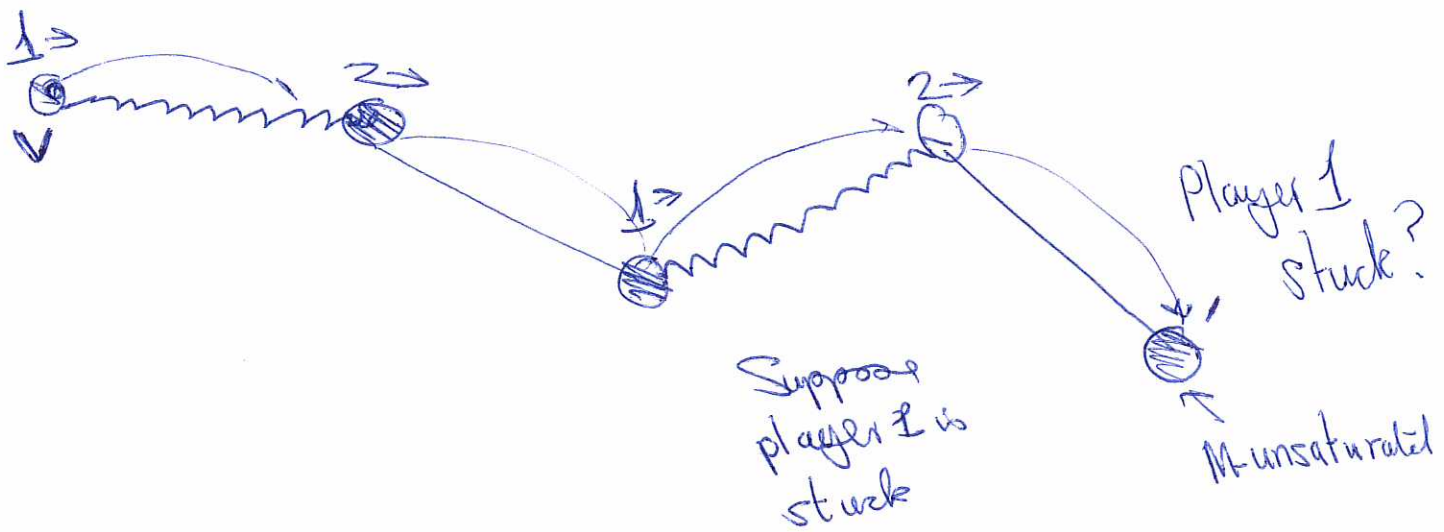
$(G, v)$  is an  $N$ -position iff every maximum matching in  $G$ , covers  $v$ .

## Proof

Suppose  $v$  is in every maximum matching.

Player 1 chooses some maximum matching  $M$ .

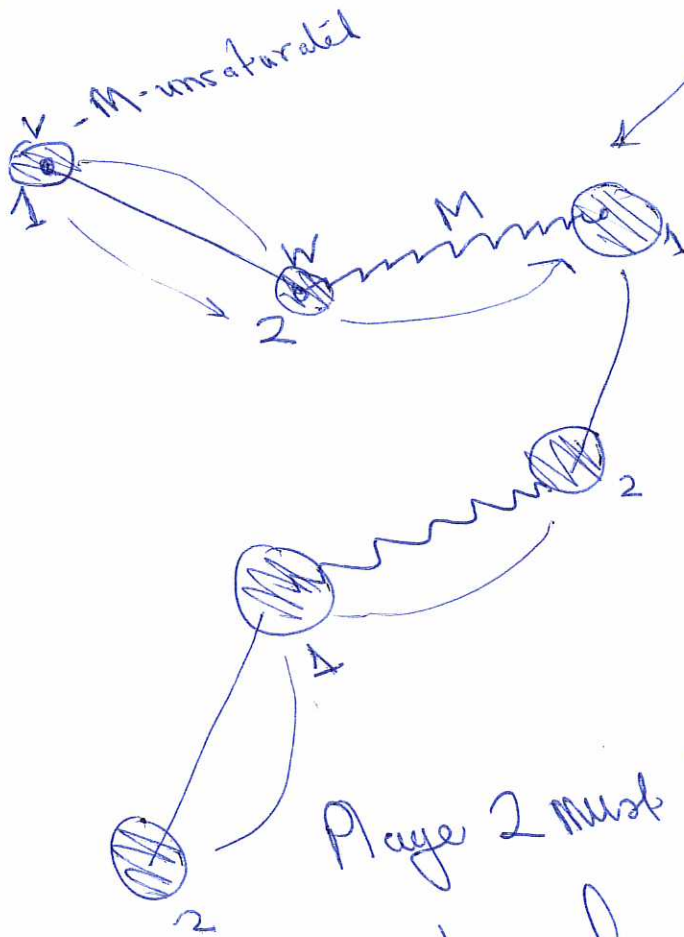
Moves along  $M$ -edge



Alternating path give a matching  $M'$  s.t.

$|M'| = |M| + 1$   $M'$  doesn't cover  $v$ . - contradiction

Suppose next that  $\exists$  maximum matching  $M$  which does not cover  $v$ .



exists edge  $w$  is  $M$ -unsaturated and  $M + \{v, w\}$  is a bigger matching than  $M$ .

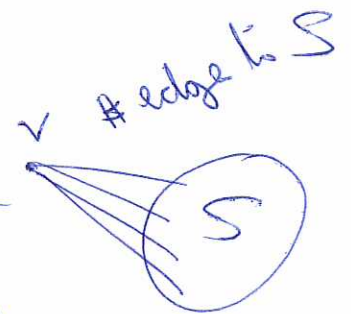
Player 2 must have a ~~more~~ more edge we have found an augmenting path i.e.  $M$  is not maximum.

# Undirected Edge Graph - Bipartite Graph.

~~Let~~  $G$  is a graph,  $S$  is an even kernel if

①  $S$  is an independent set i.e. no edge is contained in  $S$

②  $v \notin S$  then  $\deg_S(v)$  is even



## Lemma

If  $S$  is an even kernel and  $v \notin S$  then  $(G, v)$  is a P-position.

## Proof

