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Partially Ordered Sets [POSETS]

A poset is a set P plus a binary relation $\leq$ such that

- ① $a \leq a, \forall a$ (reflexive)
- ② $a \leq b$ and $b \leq c \Rightarrow a \leq c$ (transitive)
- ③ $a \leq b$ and $b \leq a \Rightarrow a = b$ (anti-symmetry)

Examples

① $P = \{1, 2, \dots\} \quad \leq = \leq \quad \left[\begin{smallmatrix} \text{total} \\ \text{order} \\ \text{here} \end{smallmatrix} \right]$

② $P = \{1, 2, \dots\} \quad a \leq b \text{ iff } a \text{ divides } b$

③ $P = \{A_1, A_2, \dots, A_m\}$ where A_i are sets and
 $\leq = \subseteq$

a & b are comparable if $a \leq b$ or $a \geq b$

otherwise a, b are incomparable. In a linear order, all pairs are comparable.

$a < b$ if $a \leq b$ and $a \neq b$

A **chain** is a sequence

$$a_1 < a_2 < a_3 < \dots < a_s$$

[length = s]

An **anti-chain** is a set A where

$a, b \in A \Rightarrow a \not\leq b$ are incomparable.

Sperner family $\equiv$ anti-chain in example 3

$$\emptyset \subseteq \{1, 2\} \subseteq \{1, 2, 5, 8\} \subseteq \{1, 2, 5, 8, 12\}$$

Chain in example 3

$$\{\{1, 2, 3\}, \{1, 2, 4\}, \{2, 3, 4\}\}$$

is an anti-chain

If C is a chain and A is an anti-chain

then

$$|A \cap C| \leq 1$$

Suppose $x, y \in A \cap C$

$x, y \in C \Rightarrow x < y$ or $y < x$ — contradicts $x, y \in A$

Two theorems relate maximum size of chain (anti-chain) to minimum number of anti-chains (chains) needed to cover P .

A chain cover of P is a set of chains $C_1, C_2, \dots, C_k$ such that $C_1 \cup C_2 \cup \dots \cup C_k = P$.
Its size is k .

An anti-chain cover is a set of anti-chains $A_1, A_2, \dots, A_k$ such that $A_1 \cup A_2 \cup \dots \cup A_k = P$ [size = k]

If P can be covered with k chains [anti-chains]
then largest anti-chain [Chain] has size $\leq k$

Pigeon hole principle.

Anti-chain of size $k+1$ would have two
[chain] elements in same chain
[anti-chain]

Theorem P is a finite poset.

$\nu = \min \{m : P \text{ can be covered with } m \text{ anti-chains}\}$

$\mu = \max \text{ size of a chain.}$

Then $\mu = \nu$.

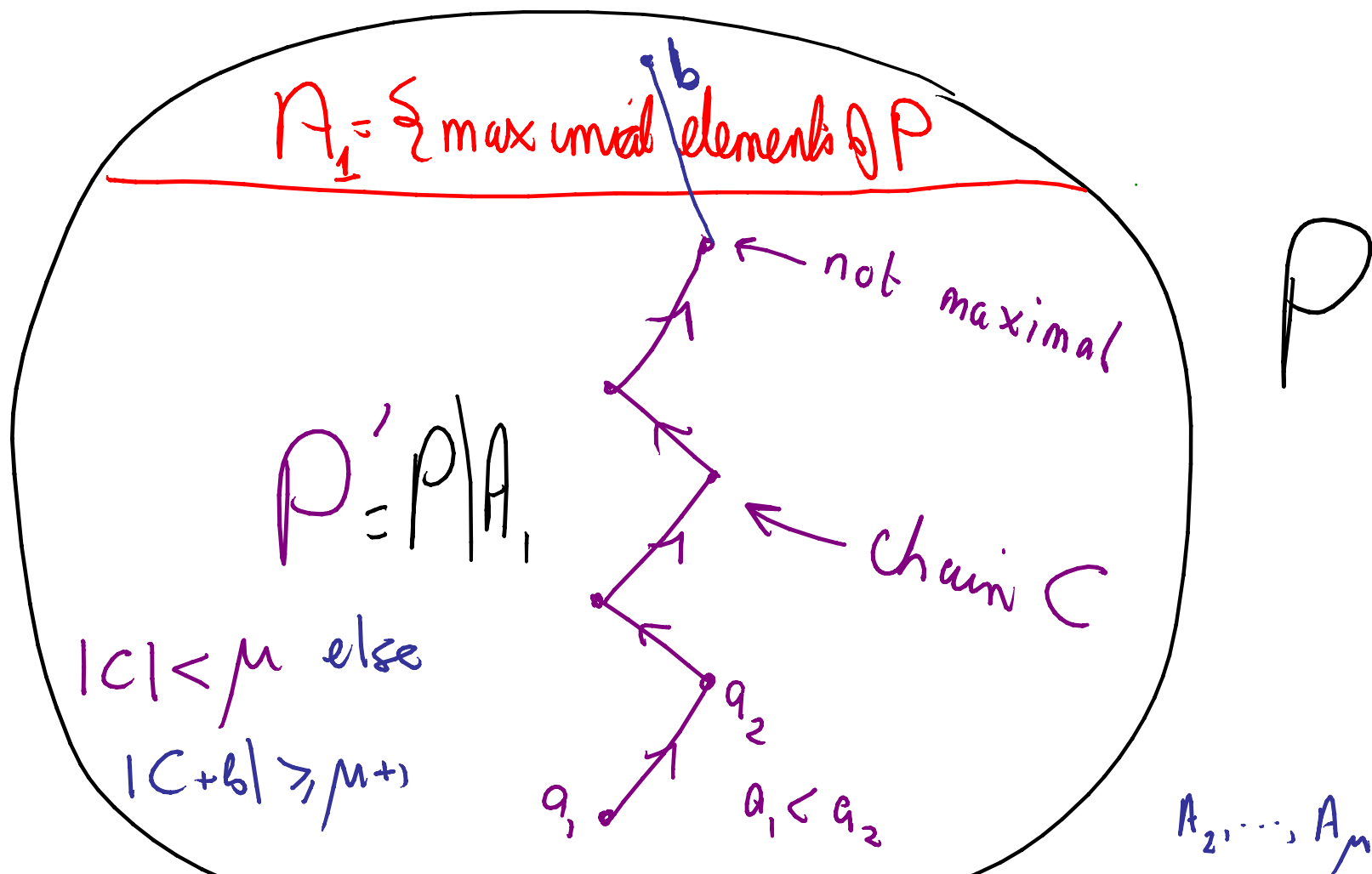
Proof

By induction on μ .

$\mu=1 \Rightarrow P$ is an anti chain

a is maximal, if $\nexists b$ such that $a < b$

μ = maximum chain length



By induction, P' can be partitioned into $\mu - 1$ anti-chains.

$$P = \{1, 2, 3, \dots, 10\}$$

$\leq \equiv$ divisibility

$$A_1 = \{6, 7, 8, 9, 10\}$$

Dilworth's Theorem

P is a finite poset.

$$\alpha = \min \{ m : P = C_1 \vee C_2 \vee \dots \vee C_m \}$$

chains

$\beta = \max$ size of anti-chain

Then

$$\alpha = \beta$$

Applications

① Erdős, Szekeres : $a_1, a_2, \dots, a_{n^2+1}$ contains a monotone subsequence of length $n+1$.

$$P = \{ (i, a_i) : 1 \leq i \leq n^2+1 \}$$

$$(i, a_i) < (j, a_j) \text{ iff } i < j \text{ and } a_i \leq a_j$$

A chain corr. to a monotone increasing subsequence.

$$(i_1, a_{i_1}) < (i_2, a_{i_2}) < \dots$$

$$i_1 < i_2 \text{ and } a_{i_1} \leq a_{i_2} \text{ etc}$$

Suppose ~~$\exists$~~ a monotone increasing subsequence of length $n+1$

$\Rightarrow$ minimum cover by chains
need $\geq \lceil \frac{n^2+1}{n} \rceil = n+1$ chains

$\Rightarrow \exists$ an ant-chain of size $n+1$
 $(i_1, a_{i_1}), (i_2, a_{i_2}), \dots, (i_{n+1}, a_{i_{n+1}})$
where $i_1 < i_2 < \dots < i_{n+1}$. So $a_{i_1} > a_{i_2} > \dots > a_{i_{n+1}}$

N intervals $I_1 = [a_1, b_1], I_2 = [a_2, b_2], \dots, I_N = [a_N, b_N]$ on the
real line $a_i \leq b_i, 1 \leq i \leq N = n^2 + 1$.

Either $\exists$ $n+1$ disjoint intervals $I_{i_1}, I_{i_2}, \dots, I_{i_{n+1}}$

or $\exists$ $n+1$ intervals $I_{i_j}, i_j \in J$ s.t. $\bigcap_{j \in J} I_{i_j} \neq \emptyset$

Extrema:

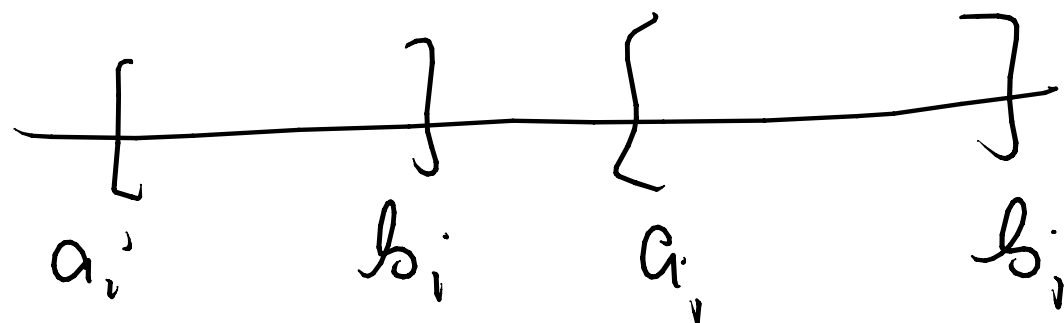


a

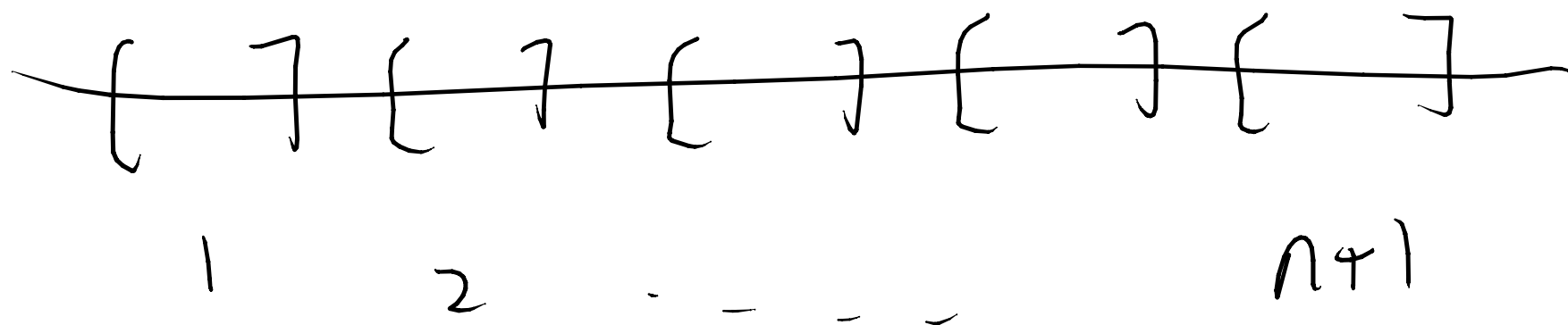


$$P = \{1, 2, \dots, N\}$$

$$i < j \Rightarrow$$



A chain of length $n+1$ looks like

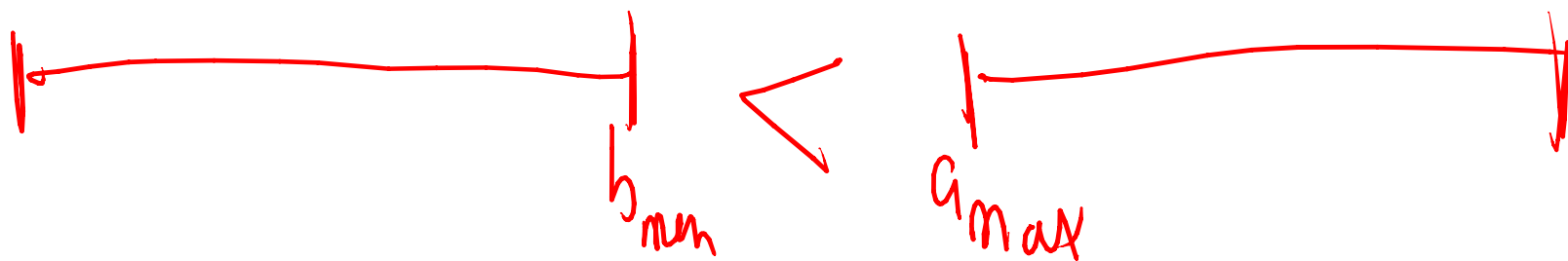


If there is no chain of length $n+1$ then
there must be an anti-chain of size

$$\left\lceil \frac{n+1}{2} \right\rceil = n+1$$

Anti-chain = $\{I_j, j \in J\}$

Claim: $\bigcap_{j \in J} I_j \neq \emptyset \equiv \begin{matrix} b_{\min} = \text{smallest } b_j \\ a_{\max} = \text{largest } a_j \end{matrix}$



$b_{\min} = \text{smallest } b_i$

$a_{\max} = \text{largest } a_i$

$$[a_{\max}, b_{\min}] \subseteq I_j, \quad j \in J$$