

11/9/11

Partially Ordered Sets

A Poset is a set P plus an order relation \leq

(I) $a \leq a$ (Reflexive)

(II) $a \leq b \& b \leq c \Rightarrow a \leq c$ (transitive)

(III) $a \leq b \& b \leq a \Rightarrow a = b$ (anti-symmetry)

Examples

(I) $P = \{1, 2, 3, \dots\}$ and $a \leq b$ has usual meaning
(total order)

(II) $P = \{1, 2, 3, \dots\}$ and $a \leq b$ iff a divides b

(III) $P = \{A_1, A_2, \dots, A_m\}$

m elements

$A_i \subseteq A$
 $i=1, 2, \dots, m$

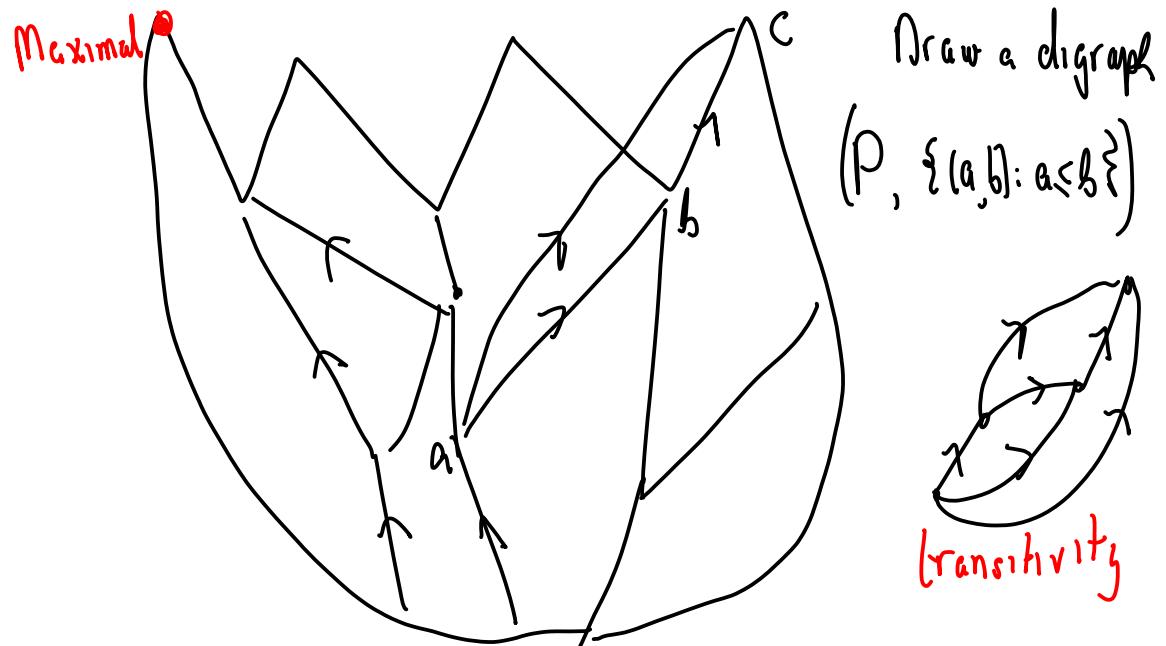
$A_i \leq A_j$ iff $A_i \subseteq A_j$

(III) $2 \nleq 3 \wedge 3 \nleq 2$

2 and 3 are
incomparable

A pair of elements a, b are comparable if

$$a \leq b \text{ or } b \leq a.$$



A poset without incomparable pairs is called a total order.

$a < b$ if $a \leq b$ and $a \neq b$

A chain is a sequence

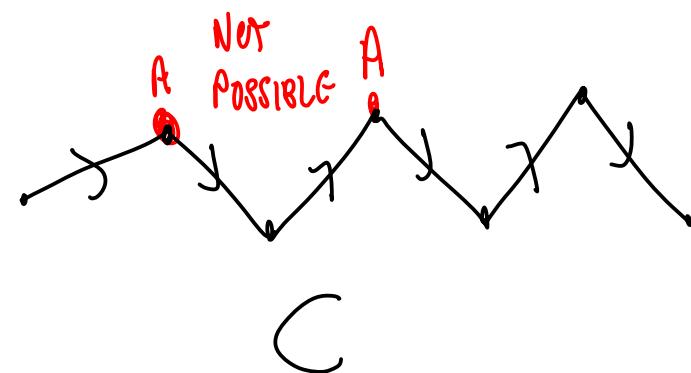
$$a_1 < a_2 < a_3 < \dots < a_k$$

A set A is an anti-chain if every $a \neq b \in A$ are incomparable

A Sperner family is an anti-chain
for example $\{\{1\}, \{2\}, \{3\}\}$.

If C is a chain and A is an
anti-chain then

$$|A \cap C| \leq 1$$



Questions

Cover = every element in one of the sets

- ① Size of largest chain \vee # of anti-chain needed to cover P



anti-chains needed \geq size of largest chain

- ② Size of largest anti-chain \vee # of chain needed to cover P

chain needed \geq size of largest anti-chain

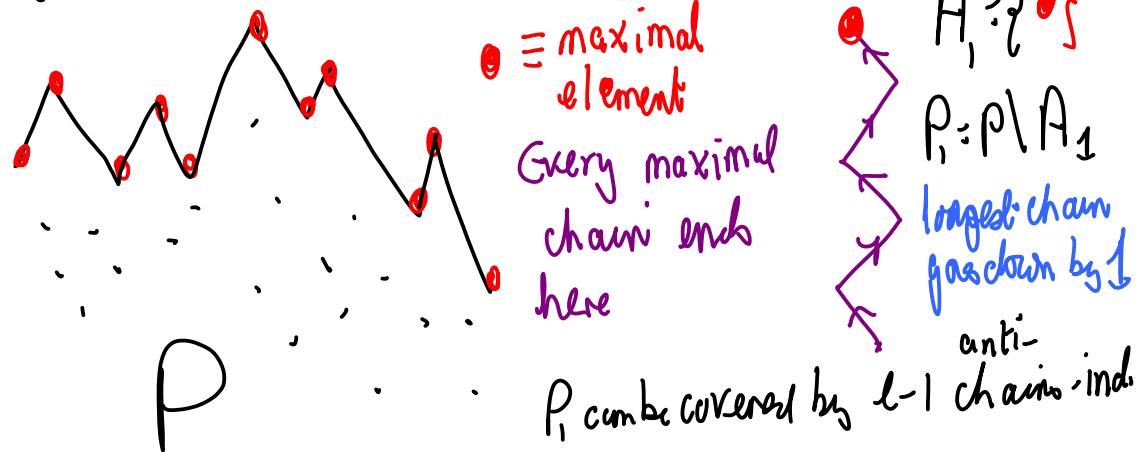
Theorem

min. # anti-chains needed to cover P

: max. size of a chain.

Proof

By induction on length of longest chain $\leq l$



Dilworth's Theorem

min # chains needed to cover P

: max size of largest anti-chain

Proof - later

Application 1

Erdős-Szekeres Theorem

a_1, a_2, \dots, a_n are real.

l^+ = length of longest monotone increasing sequence

l^- = length of longest monotone decreasing sequence

$$l^+ \times l^- \geq n.$$

$\max\{l^+, l^-\} > \sqrt{n}$
 $n = k^2 + 1 \quad > \sqrt{k^2 + 1} = k + 1$

What is the poset?

$$P = \{ (a_i, i) : i=1, 2, \dots, n \}$$

$$(a_i, i) \leq (a_j, j) \text{ if } a_i \leq a_j \text{ & } i < j$$

Chains give rise to monotone increasing sequences.

Anti-chains give rise to monotone decreasing sequences

$$A = \{ (a_{l_1}, l_1), (a_{l_2}, l_2), \dots \}$$

$$l_1 < l_2 < l_3 < \dots \quad a_{l_1} > a_{l_2} \text{ else not anti-chain}$$

Dilworth's Theorem \Rightarrow

$$\underbrace{|\text{longest chain}|}_{c} \times \underbrace{|\text{largest anti-chain}|}_{a} \geq |P|$$

$$P = C_1 \cup C_2 \cup \dots \cup C_{a=}$$

$|C_i| \leq c$