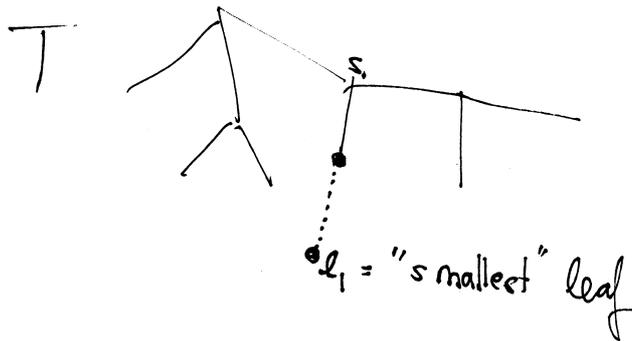
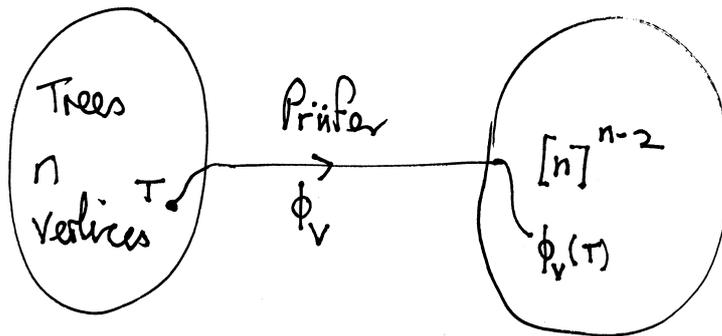
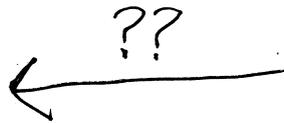


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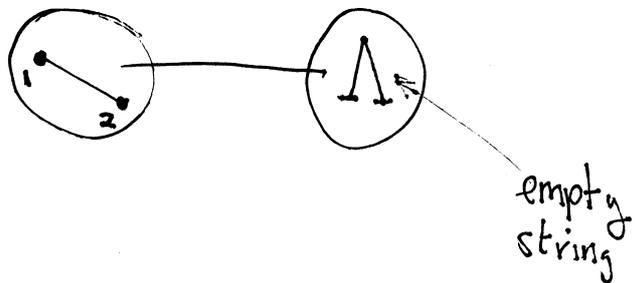
$s_1 \dots$



We show ϕ_v has an inverse.

Prove existence of Φ_V^{-1} by induction on $|V|$.

$|V|=2$.

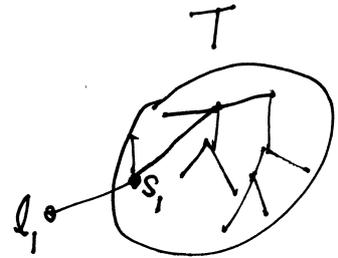


Sequence s_1, s_2, \dots, s_{n-2}

\downarrow

Tree

(1) Find $l_1 =$ smallest value not mentioned in s_1, s_2, \dots, s_{n-2}



(ii) $V \rightarrow V / \{l_1, s_1\}$

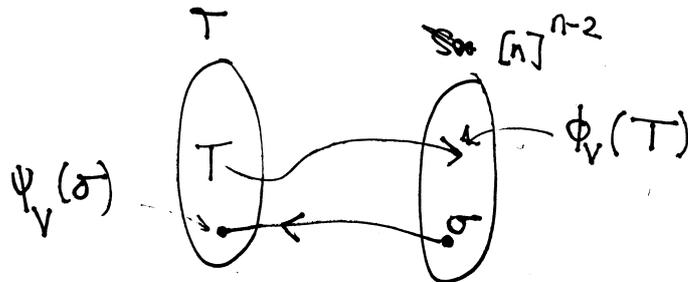
Build $T_1 = T - (l_1, s_1)$

from s_2, s_3, \dots, s_{n-2} [Use induction] (Recursion)

$$\psi_V \left(\underbrace{\phi_V^{-1}(s_1 s_2 \dots s_{n-2})}_{\text{tree } T} \right) = \text{edge}(u, s_1) + \text{tree } T_1,$$

$$T_1 = \phi_{V_1}^{-1}(s_2 s_3 \dots s_{n-2})$$

$$\text{where } V_1 = V \setminus \{u\}.$$



Claim $\psi_V = \phi_V^{-1}$

$$\begin{aligned} & \phi_V(\psi_V(s_1 s_2 \dots s_{n-2})) \\ &= \phi_V(\text{edge}(u, s_1) + T_1) \\ &= s_1 \phi_{V_1}(T_1) \\ &= s_1 s_2 \dots s_{n-2}. \end{aligned}$$

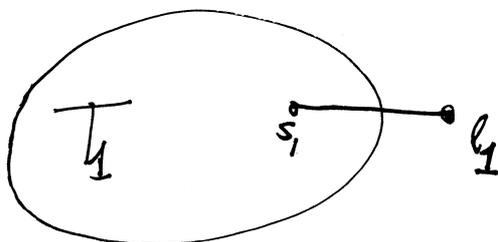
$$T \rightarrow \phi_v(T)$$

$v \in T$ appears exactly $d_T(v) - 1$ times.

[Leaves do not appear in the sequence]

Proof by induction on $|V|$

$$|V| = 2 \quad \checkmark$$



r appears 0 times

$$s_1 \text{ appears } 1 + \underbrace{\text{\#times in } \phi_{s_1}(T_1)}_{\deg_{T_1}(s_1) - 1} + 1$$

$$= \deg_{T_1}(s_1) - 1 + 1$$

$$= \deg_T(s_1) - 1$$

Every other vertex x has

$$\deg_{T_1}(x) = \deg_T(x).$$

trees with degree sequence d_1, d_2, \dots, d_n

$$d_1 + d_2 + \dots + d_n = 2n - 2$$

= # Sequences in $[n]^{n-2}$ where i appears $d_i - 1$ times, $\forall i$

$$= \binom{n-2}{d_1-1, d_2-1, \dots, d_n-1}$$

Cayley's Formula = # spanning trees of K_n

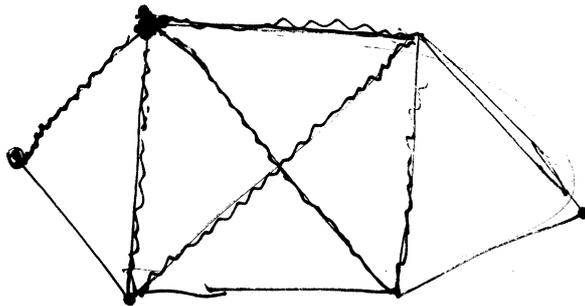
Arbitrary connected graph G etc.

spanning trees of G = determinant of $\begin{pmatrix} \text{Some} \\ (n-1) \times (n-1) \\ \text{matrix} \end{pmatrix}$

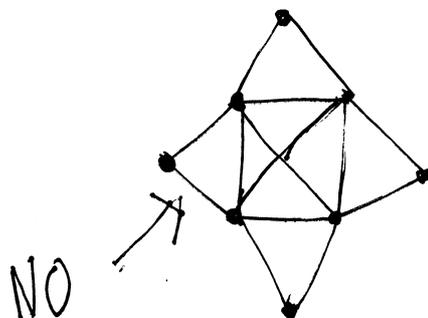
Matrix-Tree Theorem

Eulerian Graphs

An Eulerian cycle is a closed walk through G which goes through each edge exactly once.



Q? Which graphs have Eulerian cycles?



First result in Graph Theory

Bridge of Königsberg Problem.



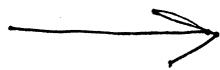
Thm

G has an Euler cycle iff

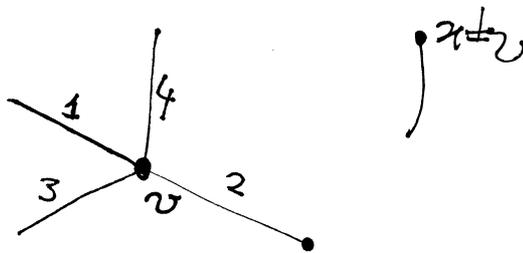
- (i) G is connected
- (ii) Every degree is even.

Proof

Euler cycle



even degrees.



Assume \exists
euler cycle C .

Traverse it.

We won't start at v

We start at x

vertices of odd degree



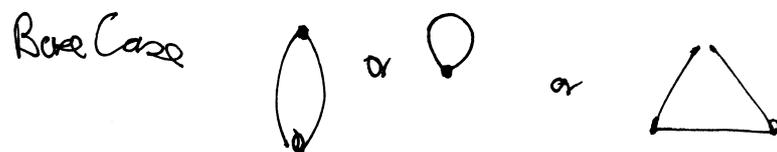
Now eulerian



Euler trail

eulerian \longleftarrow even & connected

By induction on $|E|$.



Inductive Step.

G has at least one cycle C .

Remove C .

