

09/27/06

Boole's Inequality (Union Bound)

$$P(A \cup B) = P(A) + P(B) - P(A \cap B) \geq 0$$
$$\leq P(A) + P(B)$$

If $A \cap B = \emptyset$ [Disjoint events]

$$P(A \cup B) = P(A) + P(B)$$

$$P\left(\bigcup_{i=1}^n A_i\right) \leq P\left(\bigcup_{i=1}^{n-1} A_i\right) + P(A_n)$$

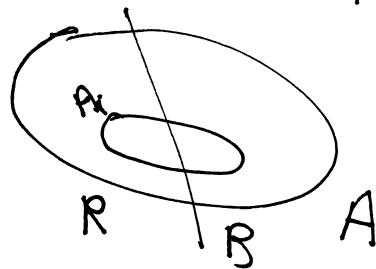
/ induction

$$\leq P(A_1) + \dots + P(A_{n-1}) + P(A_n)$$

Coloring Problem

$A_1, A_2, \dots, A_n$ are subsets of A

Question: Does $\exists$ partition of $A = R \cup B$
such that $A_i \cap R \neq \emptyset, A_i \cap B \neq \emptyset, 1 \leq i \leq n$



$$A = \{1, 2, 3, 4\}$$

$$A_1 = \{1, 2, 3\}$$

R R B

$$B_1 = \{2, 3, 4\}$$

R B B

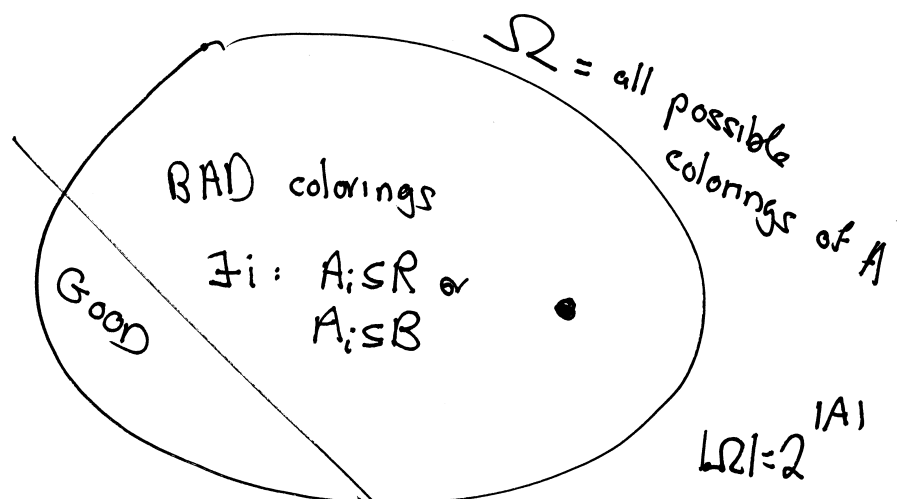
Thm

IF $|A_i| = k \forall i$ and $n < 2^{k-1}$

then a colouring exists.

PROBABILISTIC METHOD

PAUL
ERDŐS



Theorem claims that $\text{Good} \neq \emptyset$

If you make Ω into a probability space
and make each colouring equally likely
then

$$P(\text{Good}) > 0$$

or

$$P(\text{BAD}) < 1$$

$$P(\text{Good}) + P(\text{BAD}) = 1$$

$$P(\text{BAD}) < 1$$

$$\text{BAD} = \bigcup_{i=1}^n \text{BAD}_i$$

$$\{A_i \subseteq R \text{ or } A_i \subseteq B\}$$

$$P(\text{BAD}) \leq \sum_{i=1}^n P(\text{BAD}_i)$$

To choose an element from Ω , I randomly color each element of A .

$$P(\text{BAD}) \leq \sum_{i=1}^n \left(\frac{1}{2^k} + \frac{1}{2^k} \right) = \frac{n}{2^{k-1}} < 1$$

□

Can you always 2-colour $A_1, \dots, A_n$?

$$A = [2k-1] \quad n = \binom{2k-1}{k} > 2^{k-1}$$

$$A_1, A_2, \dots, A_n = \binom{A}{k}$$

In any 2-coloring of A , there is mono-chromatic Set.

$$A = R \cup B$$

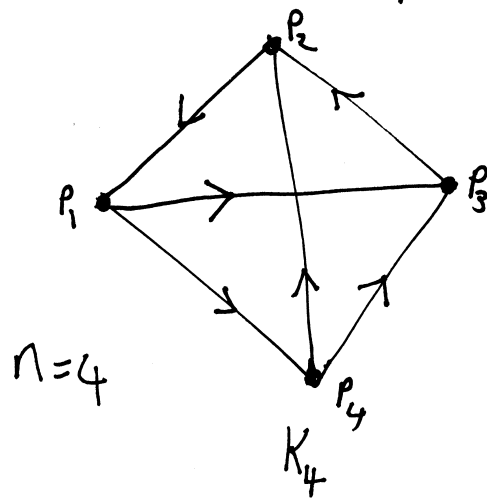
$$|R| \geq k \text{ or } |B| \geq k$$

Assume $|R| \geq k$ then

any k -subset of R is all Red.

Tournaments

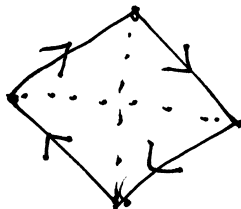
n players. Each play each other. No ties.



"Orientation of K_n "

For k . Does $\exists$ a tournament such $|S| = k$
 then $\exists$ player w_s who beats all of them.

$k=1$: easy



$k=2$: think about it.

We prove the existence of a tournament
by showing that a random tournament
has the property with positive probability.

BAD: $\exists S$ s.t. nobody beat S

$$\text{BAD} = \bigcup_{\substack{S \subseteq [n] \\ |S|=k}} \text{BAD}_S$$

$$P(\text{BAD}) \leq \sum_{\substack{S \subseteq [n] \\ |S|=k}} P(\text{BAD}_S)$$

$$\left(1 - \frac{1}{2^k}\right)^{n-k}$$



$$\begin{aligned} P(x \text{ beats } S) &= \frac{1}{2^k} \\ P(x \text{ does not beat } S) &= 1 - \frac{1}{2^k} \end{aligned}$$

$$P(\text{BAD}) \leq \binom{n}{k} \left(1 - \frac{1}{2^k}\right)^{n-k}$$

k is fixed (say 100), n can be arbitrarily large

$$\leq n^k e^{-\frac{n-k}{2^k}}$$

$$\underline{1-x \leq e^{-x}}$$

$$\rightarrow 0$$

$$k \log n \ll \frac{n-k}{2^k}$$