The **girth** of a graph is the length of the shortest cycle in G. At first sight it seems that a graph having large girth would necessarily be colourable with few colours i.e. there might exist a function f such that if G has girth at least a then its chromatic number would be at most f(a). This is ruled out by the following theorem of Erdős:

**Theorem 1.** Let a, b > 0 be positive integers. Then there exists a graph with girth  $g \ge a$  and chromatic number  $\chi \ge b$ .

**Proof** This follows from the following lemma:

**Lemma 1.** Let d be large enough so that  $\frac{d}{3 \log d} \geq b$ . Then there exists a graph G with  $n \geq 100d^a$  such that

- (i) There are at most  $2d^a$  cycles of length  $\leq a$ .
- (ii) There is no independent set of vertices of size  $\geq \frac{2 \log d}{d} n$ .

From the lemma, let G' be obtained from G by deleting one vertex from each cycle of length  $\leq a$ . Then G' has girth  $\geq a$  and

$$\chi \ge \frac{n - 2d^a}{\frac{2\log d}{d}n} \ge \frac{d}{3\log d} \ge b$$

**Proof of Lemma:** Let p = d/n and consider  $G_{n,p}$ .

Markov Inequality: If X is a non-negative random variable then

$$\mathbf{Pr}(X \ge t) \le \frac{\mathbf{E}(X)}{t}$$
 for any  $t > 0$ .

**Proof:** 

$$\begin{split} \mathbf{E}(X) &= \mathbf{E}(X \mid X \geq t) \mathbf{Pr}(X \geq t) + \mathbf{E}(X \mid X < t) \mathbf{Pr}(X < t) \\ &\geq \mathbf{E}(X \mid X \geq t) \mathbf{Pr}(X \geq t) \\ &\geq t \mathbf{Pr}(X \geq t). \end{split}$$

(i) Let X be the number of cycles of length at most a. Then

$$\begin{split} \mathbf{E}(X) &\leq \sum_{k=3}^{a} \binom{n}{k} \frac{(k-1)!}{2} p^{k} \\ &\leq \sum_{k=3}^{a} \frac{n^{k}}{k!} \frac{(k-1)!}{2} \frac{d^{k}}{n^{k}} \\ &= \sum_{k=3}^{a} \frac{d^{k}}{2k} \\ &< d^{a}. \end{split}$$

Applying the Markov inequality we get

$$\mathbf{Pr}(X \ge 2d^a) \le \frac{E(X)}{2d^a} \le \frac{1}{2}.$$

(ii) Let  $k = \left\lceil \frac{2 \log d}{d} n \right\rceil$  and Y be the number of independent sets of size k in  $G_{n,p}$ . Then

$$\mathbf{E}(X) = \binom{n}{k} \left(1 - \frac{d}{n}\right)^{\binom{k}{2}}$$

$$\leq \left(\frac{ne}{k} \cdot \left(1 - \frac{d}{n}\right)^{\frac{k-1}{2}}\right)^{k}$$

$$\leq \left(\frac{de}{2\log d} \cdot e^{d/2n} \cdot e^{-dk/2n}\right)^{k}$$

$$\leq \left(\frac{de}{2\log d} \cdot (1 + o(1)) \cdot \frac{1}{d}\right)^{k}$$

$$\to 0.$$

 $\operatorname{So}$ 

$$\mathbf{Pr}(G_{n,p} \text{ satisfies (i) and (ii)}) \ge \frac{1}{2} - o(1) > 0.$$