

Class 04

Theorem 1. *There are n^{n-2} different trees with vertex set $[n]$.*

Proof Let \mathcal{T} be the set of trees with vertex set $[n]$. We define a map $\mathcal{P} : \mathcal{T} \rightarrow [n]^{n-2}$ such that \mathcal{P} is a bijection. Then $|\mathcal{T}| = |[n]^{n-2}| = n^{n-2}$.

We define a sequence $(x_i, y_i, T_i), i = 1, 2, \dots, n-1$ where $T_1 = T$ and x_i is the least valued vertex of degree 1 in T_i and y_i is the **unique** neighbour of x_i in T_i .

$$\mathcal{P}(T) = y_1 y_2 \cdots y_{n-2}.$$

We need to show that we can recover T from $\mathcal{P}(T)$ i.e. \mathcal{P} is 1-1 and onto. We can prove this by induction on n . Observe first that vertex v occurs $\deg_T(v) - 1$ times among y_1, y_2, \dots, y_{n-2} . This is because it occurs $\deg_T(v)$ times among $x_1, y_1, x_2, y_2, \dots, x_{n-1}, y_{n-1} = n$ and exactly once among $x_1, x_2, \dots, x_{n-1}, y_{n-1}$.

Thus x_1 is the least element of V which does not appear in y_1, y_2, \dots, y_{n-1} and then T is T_1 plus vertex x_1 and edge (x_1, y_1) . The rest of the code defines T_1 , by induction. (Note that T_1 is a tree on $[n] \setminus \{x_1\}$ and so we really need to define the code for trees with vertices in any ordered set). \square