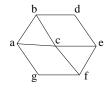
Paths and Walks

 $W=(v_1,v_2,\ldots,v_k)$ is a walk in G if $(v_i,v_{i+1})\in$ E for $1 \leq i < k$.

A path is a walk in which the vertices are distinct.

 W_1 is a path, but W_2, W_3 are not.

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 $W_1 = a,b,c,e,d$ $W_2 = a,b,a,c,e$

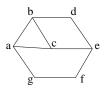
 $W_3 = g,f,c,e,f$

for v_1, v_k

A walk is *closed* if $v_1=v_k$. A *cycle* is a closed walk in which the vertices are distinct except

b, c, e, d, b is a cycle.

b, c, a, b, d, e, c, b is not a cycle.



Connected components

We define a relation \sim on V. $a \sim b$ iff there is a walk from a to b.





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 $a \sim b$ but $a \not\sim d$.

Claim: \sim is an equivalence relation.

reflexivity $v \sim v$ as v is a (trivial) walk from v to v.

Symmetry $u \sim v$ implies $v \sim u$.

 $(u = u_1, u_2 \dots, u_k = v)$ is a walk from u to v implies $(u_k, u_{k-1}, \dots, u_1)$ is a walk from v to u.

Transitivity $u \sim v$ and $v \sim w$ implies $u \sim w$. $W_1 = (u = u_1, u_2 \dots, u_k = v)$ is a walk from u to v and $W_2 = (v_1 = v, v_2, v_3, \dots, v_{\ell} =$ w) is a walk from v to w imples that $(W_1, W_2) = (u_1, u_2, \dots, u_k, v_2, v_3, \dots, v_\ell)$ is a walk from u to w.

The equivalence classes of \sim are called connected components.

In general $V = C_1 \cup V_2 \cup \cdots \cup C_r$ where C_1, C_2, \ldots , C_r are the connected comonents.

We let $\omega(G)(=r)$ be the number of components of G.

G is connected iff $\omega(G)=1$ i.e. there is a walk between every pair of vertices.

Thus C_1, C_2, \ldots, C_r induce connected subgraphs $G[C_1], \ldots, G[C_r]$ of G

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For a walk W we let $\ell(W)=$ no. of edges in W.



Lemma 1 Suppose W is a walk from vertex a to vertex b and that W minimises ℓ over all walks from a to b. Then W is a path.

Proof Suppose $W=(a=a_0,a_1,\ldots,a_k=b)$ and $a_i=a_j$ where $0\leq i< j\leq k$. Then $W'=(a_0,a_1,\ldots,a_i,a_{j+1},\ldots,a_k)$ is also a walk from a to b and $\ell(W')=\ell(W)-(j-i)<\ell(W)$ – contradiction.

Corollary 1 If $a \sim b$ then there is a path from a to b.

So G is connected $\leftrightarrow \forall a,b \in V$ there is a path from a to b.

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In BFS we construct A_0, A_1, A_2, \ldots , by

 $A_{t+1} = \{ w \notin A_0 \cup A_1 \cup \cdots \cup A_t : \exists \text{ an edge } (u,w) \text{ such that } u \in A_t \}.$

Note: no edges (a,b) between A_k and A_ℓ for $\ell-k\geq 2$, else $w\in A_{k+1}\neq A_\ell$. (1)

In this way we can find all vertices in the same component ${\it C}$ as ${\it v}$.

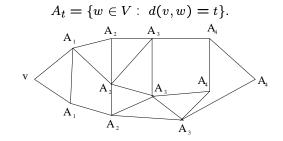
By repeating for $v' \notin C$ we find another component etc.

Breadth First Search - BFS

Fix $v \in V$. For $w \in V$ let

d(v,w) = length of shortest path from v to w.

For t = 0, 1, 2, ..., let



 $A_0 = \{v\}$ and $v \sim w \leftrightarrow d(v, w) < \infty$.

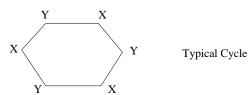
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Characterisation of bipartite graphs

Theorem 1 G is bipartite $\leftrightarrow G$ has no cycles of odd length.

Proof \rightarrow : $G = (X \cup Y, E)$.



Suppose $C=(u_1,u_2,\ldots,u_k,u_1)$ is a cycle. Suppose $u_1\in X$. Then $u_2\in Y,u_3\in X,\ldots,u_k\in Y$ implies k is even.

 \leftarrow Assume G is connected, else apply following argument to each component.

Choose $v \in V$ and construct A_0, A_1, A_2, \ldots , by BFS.

 $X = A_0 \cup A_2 \cup A_4 \cup \cdots$ and $Y = A_1 \cup A_3 \cup A_5 \cup \cdots$

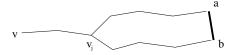
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We need only show that X and Y contain no edges and then all edges must join X and Y. Suppose X contains edge (a,b) where $a\in A_k$ and $b\in A_\ell.$

(i) If $k \neq \ell$ then $|k - \ell| \geq 2$ which contradicts (1)





There exist paths $(v=v_0,v_1,v_2,\ldots,v_k=a)$ and $(v=w_0,w_1,w_2,\ldots,w_k=b).$

Let
$$j = \max\{t: v_t = w_t\}$$
.

$$(v_j, v_{j+1}, \ldots, v_k, w_k, w_{k-1}, \ldots, w_j)$$

is an odd cycle — length 2(k-j)+1 — contradiction. $\hfill\Box$

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