

The topology of competitively constructed graphs

Alan Frieze*, Wesley Pegden

December 5, 2013

Abstract

We consider a simple game, the k -regular graph game, in which players take turns adding edges to an initially empty graph subject to the constraint that the degrees of vertices cannot exceed k . We show a sharp topological threshold for this game: for the case $k = 3$ a player can ensure the resulting graph is planar, while for the case $k = 4$, a player can force the appearance of arbitrarily large clique minors.

1 Introduction

In some sense, restricting one's attention to 3-regular graphs is not a topological constraint at all, in the sense that connected 3-regular graphs can require arbitrarily complex surfaces to be embedded in, or, say, contain arbitrarily large clique minors. In particular, from a topological point of view, vertices of degree 3 are essentially different from vertices of degree 2. One might then expect the presence of degrees greater than 3 to lead to a similar increase in topological trouble. For example, out of the list of 106 forbidden topological minors for embeddability in the projective plane, only 6 are required to ensure embeddability of cubic graphs in the projective plane [1, 3].

We show another kind of topological threshold between degree-3 and degree-4 graph vertices. Consider a game (the k -regular graph game) in which two players take turns adding edges to an initially empty graph. Players are allowed to add edges only between pairs of vertices which were previously nonadjacent and of degree $\leq k - 1$. The game ends when this is no longer possible. In particular, the degree of every vertex in the resulting graph will be exactly k , with at most k exceptions.

Theorem 1.1. *Either player has a strategy to ensure that the 3-regular graph game produces a planar graph.*

On the other hand, for the analogous 4-regular graph game, we have:

Theorem 1.2. *For any ℓ and sufficiently large n , either player has a strategy to ensure that the 4-regular graph game on n vertices produces a graph with a K_ℓ minor.*

*Research supported in part by NSF grant ccf1013110

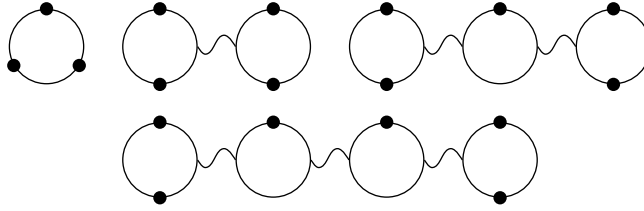


Figure 1: Components of Types 1, 2, 3a and 3b, respectively

Thus there is no surface S for which a player of the 4-regular graph game can ensure that the connected components of the result has a drawing on S .

Note that the the moves of the two players in this game are equivalent, unlike Maker-Breaker games (see [2]). In particular, with a symmetric “normal” win condition—say, the first player in the 6-regular game to break planarity loses—this would be an impartial game, subject to the the Sprague-Grundy theorem [4, 5].

2 Proofs

We begin by proving Theorem 1.1. Call the player with the goal of making the graph planar the *planar player*; his opponent is the *nonplanar player*. At any stage of the game, the *deficit* $\text{def}(v)$ of a vertex v refers to the difference between the current degree and the maximum allowable degree. Thus, in the 3-regular graph game, every vertex begins with deficit 3. The deficit of a set of vertices is the sum of their deficits.

We inductively claim that the planar player can maintain that at any stage, at the end of his move, each connected component C of G can be drawn in the plane such that its positive deficit vertices all lie on its unbounded face, and also that each C is one of the following *Types*:

1. C has deficit ≤ 3 ;
2. C has a bridge e , such that the vertex-sets of the connected components C_1, C_2 of $C \setminus e$ (the *sides* of C) each have deficit exactly 2; or
3. (a) C has bridges e_1, e_2 , such that the component C_0 of $C \setminus e_1$ which does not contain e_2 has deficit 2, the component C_2 of $C \setminus e_2$ which does not contain e_1 has deficit 2, and the component C_1 of $C \setminus \{e_1, e_2\}$ which is not C_0 or C_2 has deficit 1.
- (b) C has bridges e_1, e_2, e_3 , where e_1 and e_3 are in distinct components of $C \setminus e_2$, such that the component C_0 of $C \setminus e_1$ which does not contain e_2, e_3 has deficit 2, the component C_3 of $C \setminus e_3$ which does not contain e_1, e_2 has deficit 2, and the components C_1, C_2 of $C \setminus \{e_1, e_2, e_3\}$ which are not C_0 or C_3 have deficit 1.

Moreover, the planar player will ensure that there is at most one Type 3 component after each of his turns. At any stage, if all components are of one of these types and there is at most one component of Type 3, we say the graph satisfies condition \mathcal{T} .

Figure 1 shows schematics of these component types. Note that if a component C_i has deficit 2, it has either two vertices of deficit 1, or a single vertex of deficit 2; similarly, when a component has deficit 3, it may be distributed among vertices in 3 different ways. Our argument will not be sensitive to these distinct cases. In particular, the drawings in Figure 1 and Figure 2 use circles to denote components after the deletion of the relevant bridges, and dots to denote units of deficit within these components, but we do not assume that distinct dots drawn on individual circles correspond to distinct vertices, and it is thus not valid to assume, for example, that two components with deficit ≥ 2 can be joined by two distinct edges. By the same token, we may not assume that it is legal to draw an edge between the two sides of a Type 2 component (it may be a single edge, for example).

Our inductive argument hinges on the fact that when a planar graph has a bridge, any drawing of it can be “flipped” along the bridge to adjust the order of vertices appearing on its outer face.

Observation 2.1. *If G is drawn in the plane such that the vertices appear in the cyclic order $v_1, \dots, v_k, v_{k+1}, \dots, v_n, \dots, v_{k+1}, v_k, \dots, v_1$ along the outer face and $v_k v_{k+1}$ is a bridge of G , then G can also be drawn in the plane such that the vertex order is $v_1, \dots, v_{k+1}, v_n, v_{n-1}, \dots, v_k, \dots, v_1$ along the outer face. \square*

Since all components are of Type 1 at the beginning of the game (and after any first turn of the game), we assume by induction that condition \mathcal{T} holds and show that the planar player can respond to any move by the nonplanar player to preserve these conditions.

We show that the planar player can maintain his invariant via the following remaining cases:

1. If the nonplanar player has added an edge to a component C of Type 1 or Type 2, the result is already of Type 1; if he has joined two Type 1 components, the result is already of Type 2 or Type 1. (The planar player has a free move.)
2. If the nonplanar player has joined a Type 2 component C to a Type 1 component C' , the planar player makes the result a Type 1 component.
3. If the nonplanar player has joined two Type 2 components C, C' , the planar player can make the result a Type 2 component.
4. If the nonplanar player has joined the Type 3 component C to a Type 2 component C' , the planar player can make the result the Type 3 component.

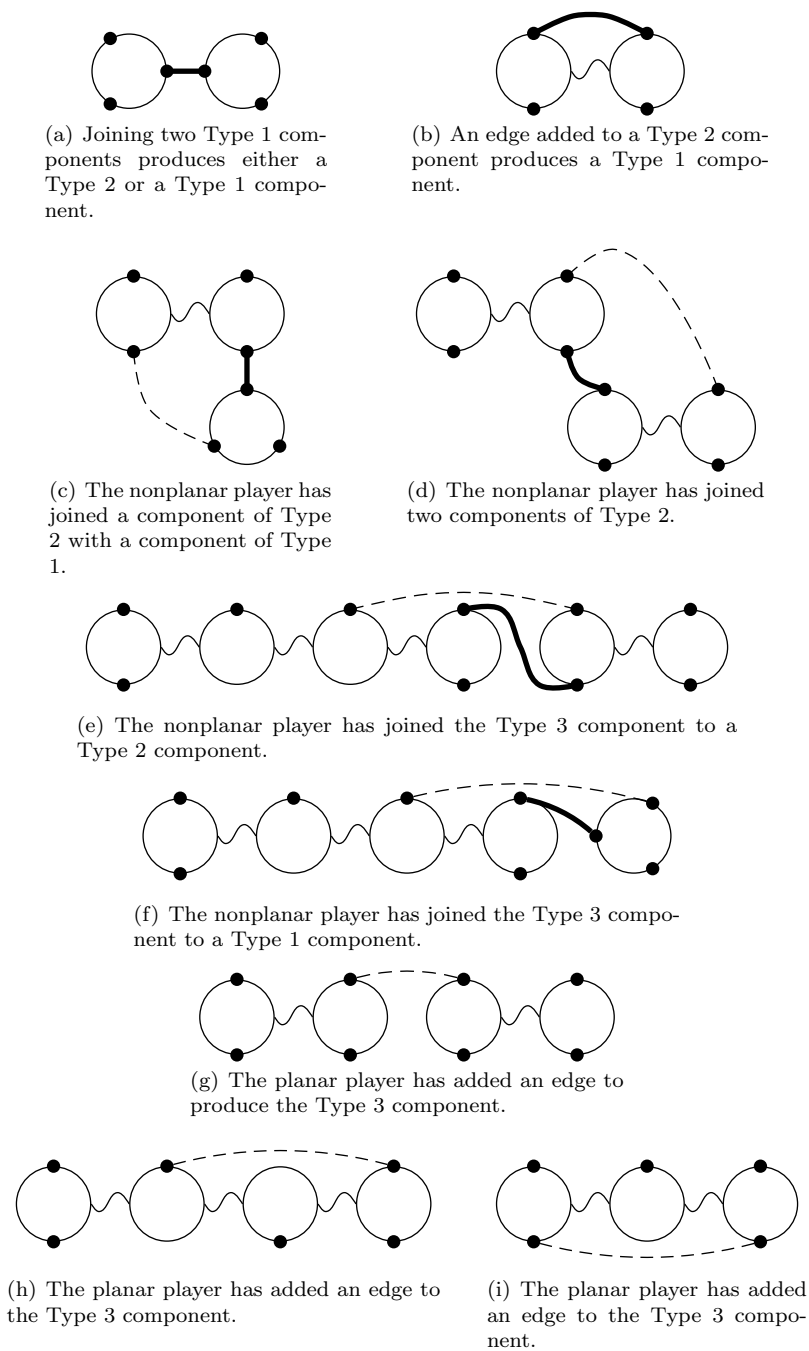


Figure 2: In each case, the nonplanar player has joined two components with the dark edge, and the planar player replies with the dashed edge.

5. If the nonplanar player has joined the Type 3 component C to a Type 1 component C' , the planar player can make the result a Type 3 or Type 2 component.
6. If the nonplanar player has added an edge within the Type 3 component C , the planar player can make the result a Type 1 component.
7. In any other case (a “free move” for the planar player) he either turns a Type 3 component into a Type 2 or Type 1 component, or creates a Type 3 component.

Case proofs:

1. If he adds an edge to a Type 1 or Type 2 component, the deficit of the component after the nonplanar player’s move is $\leq 4 - 2 = 2$. If he joins two Type 1 components, the result is a Type 2 or Type 1 component, as in Figure 2(a).
2. We have that C decomposes as C_1, C_2 joined by a bridge, each of whose vertex sets have deficit equal to 2, and that C can be drawn in the plane with all positive deficit vertices on the unbounded face. Assume that the nonplanar player’s move is an edge from C_2 to C' . Then the planar player chooses an edge between C_1 and C' , as in Figure 2(c) (unless C' now has deficit 0, in which case the planar player has a free move). The result is a Type 1 component.
3. Let the sides of C and C' , respectively be C_1, C_2 and C'_1, C'_2 . Suppose without loss of generality that the nonplanar player’s move is to take an edge from C_2 to C'_1 . The planar player responds with a move from C_2 to C'_2 (Figure 2(d)).
4. Letting C_i ($i = 0, 1, 2$ or $i = 0, 1, 2, 3$) be as in the definition of a Type 3 component, we may assume without loss of generality that the nonplanar player’s edge is from either C_0 or C_1 to C' . In the first case, the planar player responds with an edge from C_1 to C' , as in Figure 2(e). In the second case, the planar player takes an edge from C_0 to C' (the reverse of the case shown Figure 2(e)). Either way, the result is still a single component of Type 3 (of the same subtype 3a or 3b as before the nonplanar player’s move).
5. This situation is analogous to the previous one. Again, the first case is shown, in Figure 2(f). The result is either a component of Type 3a or Type 2.
6. If C was of Type 3a, then it is already a component of Type 1 after the nonplanar player’s move, so the planar player has a free move. If C was of Type 3b, then it is always possible for the planar player to add a second edge to the component, since a Type 3b component admits edges between both the pairs C_0, C_2 and C_1, C_3 of its components under removal of the its

bridges e_1, e_2, e_3 . Adding the second edge brings the deficit to 2, making the result again a Type 1 component.

7. If there remain any two Type 1 components of positive deficit, the planar player can join them to produce a Type 2 (or Type 1) component. Otherwise, if there is already a component C of Type 3, he can add an edge to C to produce a component of Type 1 or Type 2, as in Figures 2(h) and 2(i). Or if there is currently no Type 3 component, the planar player can join a Type 1 or Type 2 component C to a Type 2 component C' to produce a single component of Type 1, 2, or 3 (depending on the type of C); the case where C has Type 2 is shown in Figure 2(g). Finally, if no move described so far is possible because there is at most one component remaining in the graph, and this component is of Type 1 or Type 2, then he can make arbitrary legal moves until the end of the game without endangering planarity. \square

We turn now to the proof of Theorem 1.2. We call the player with the goal of forcing a K_ℓ minor the *minor player*, and the player with the goal of avoiding this the *structure player*. Our proof has the following two ingredients:

Lemma 2.2. *In the course of playing the 4-regular graph game, a player can force the appearance of components of arbitrarily large deficit.*

Lemma 2.3. *Consider labelings of connected bounded degree graphs with vertices labeled by bounded nonnegative labels. When the sum of the labels is sufficiently large, then we can find k disjoint connected subgraphs each with label sums $\geq s$.*

Lemma 2.2 implies the minor player can build arbitrarily large deficit components in the course of play. Applying Lemma 2.3, we see that he can find ℓ disjoint connected subgraphs each of total deficit $\geq \binom{\ell}{2}$. Over the next at most $\binom{\ell}{2}$ moves, the minor player joins previously unconnected pairs of these ℓ subgraphs (the deficit of each subgraph will remain positive while he is not yet finished), creating a K_ℓ minor. \square

All that remains is to prove Lemmas 2.2 and 2.3.

Proof of Lemma 2.2. We divide the minor player's strategy into two rounds. In the first round, he chooses m edges of a matching for some large m . Note that if we ignored the role of the structure player, the result would be a large number of components, each of size 2, of deficit 6. We let C_1, C_2, \dots, C_m denote these pairs of vertices (as sets) which the minor player has joined.

We claim that once the minor player has completed this round, the sum

$$\delta = \sum_{\text{def}(C) \geq 5} \text{def}(C)$$

(taken over all connected components of the graph which at this round have deficit ≥ 5) is large (tends to ∞ with m). To see this, let us allow even that

the structure player is given the m edges of the minor player's matching ahead of time. In absence of the structure player's moves, δ would be $6m$. We classify the structure player's moves in this round into two types:

1. Edges in components which, at the end of this round have deficit ≤ 4 , and
2. Edges in components which, at the end of this round have deficit ≥ 5 .

Let m_i denote the number of moves he makes of Type i (so $m = m_1 + m_2$). We have that

$$\delta \geq 6m - 6m_1/\beta - 2m_2,$$

where the constant β is the minimum number of edges per component required to decrease the deficit of components below 5. We need only show that $\beta > 1$. To bound β , consider any component C of the graph after this first round of play. If it contains m_C edges of the minor player's matching and m'_C edges of the structure player, then its deficit satisfies

$$\text{def}(C) \geq 6m_C - 2m'_C.$$

In particular, $\text{def}(C) \leq 4$ implies that

$$m'_C \geq 3m_C - 2.$$

Moreover, in the case where $m_C = 1$, we see that $m'_C \geq 5$ since the structure player cannot duplicate the minor player's edge (his best case is to complete a K_4). In particular, since $m'_C \geq 4$ in all cases, $\beta \geq 4$, completing the proof that δ becomes arbitrarily large.

We now show that he can force the appearance of a single component of large deficit. He simply chooses one component of deficit ≥ 5 arbitrarily, and, on each turn, grows this component by joining it to a new component of deficit ≥ 5 arbitrarily. Taking into account also the structure player's move, the deficit of this component is increasing by $\geq 5 - 2 - 2 > 0$ on each turn; in particular, it will become arbitrarily large. \square

Proof of Lemma 2.3. Let G be a connected graph with maximum degree Δ with vertices with labels $\ell(v)$ from $0, 1, \dots, b$, and let $\ell(X)$ denote the label sum of a subset $X \subset G$. Consider a spanning tree T of G . We begin by showing that when $\ell(G)$ is sufficiently large, we can find an edge e of T such that the ratio $\rho_e = \ell(C_2)/\ell(C_1)$ for the components C_1, C_2 of $T \setminus e$ satisfies

$$\frac{1}{\Delta} < \rho_e < \Delta. \tag{1}$$

To see this, fix an edge $e \in T$, and let us consider the case where $\rho_e < \frac{1}{\Delta}$. If x is the endpoint of e in the larger label-sum component C_1 of $T \setminus e$, and $e' = \{x, y\}$ is the edge for which y is in the highest label-sum component of $C_1 \setminus x$, then we have that

$$\rho_{e'} \leq \frac{\ell(C_2) + b + \frac{\Delta-2}{\Delta-1}\ell(C_1)}{\frac{1}{\Delta-1}(\ell(C_1) - b)}$$

Letting $\ell(G)$ be sufficiently large that, say, $b < (\frac{1}{\Delta-1} - \frac{1}{\Delta})\ell(C_1)$ (so that also $b < \frac{\ell(C_1)}{\Delta}$) we see that

$$\rho_{e'} \leq \frac{\Delta}{\Delta-1} \frac{\ell(C_2) + \frac{\Delta-1}{\Delta}\ell(C_1)}{\frac{1}{(\Delta-1)}\ell(C_1)} = (\Delta-1) + \Delta\rho_e.$$

In particular, if $\rho_e < \frac{1}{\Delta}$, then $\rho_e < \rho_{e'} < \Delta$; thus, we can walk along the tree to find an edge e satisfying (1).

We now simply apply our ability to find such edges recursively, t times for some t , to divide T into 2^t trees, each with label sums $\geq \frac{\ell(G)}{(\Delta+1)^t}$. \square

3 Discussion

Although we have focused on topological questions regarding the game we have introduced, many other questions seem natural as well. For example, what if we consider subgraphs instead of minors? For example:

Question 1. Which graphs H have the property that a player in the 3-regular graph game on sufficiently many vertices can ensure that the resulting graph contains a copy of H ?

For example, K_4 does not have this property; indeed no cubic graph has this property. This can be seen by showing that either player in the 3-regular graph game has a strategy to ensure that the graph is connected. (The player can maintain the invariant that while disconnected, the graph consists of isolated vertices, plus a single other component with at least one vertex of deficit 2.)

Turning back to topological issues, one can probe the relationship between degree-4 and degree-3 vertices a bit more. Let us define the *cubic+k graph game*, in which degrees of vertices must remain at most 3, except for k special vertices whose degrees may rise to 4. Let now $g^3(k)$ be the minimum genus g such that for any number of vertices n , Player 1 can ensure that any connected component of the result of the cubic+k game can be drawn in some surface of genus g . Then Theorem 1.1 implies that $g^3(0) = 0$, while the proof of Theorem 1.2, which works when some vertices have a degree threshold of 3 so long as sufficiently many have a degree threshold of 4, implies that $g^3(k) \rightarrow \infty$. Thus the asymptotic behavior of $g^3(k)$ can be studied. Rather than the particular rate of growth of $g^3(k)$, however, it may be more interesting to compare with the function $g_4(n)$, which we define as the minimum genus such that Player 1 can ensure that the any connected component of the result of the 4-regular graph game on n vertices can be drawn on some surface of genus g :

Question 2. Is $g^3(n) \sim g_4(n)$?

A first step would be finding a single value of n for which $g^3(n) \neq g_4(n)$.

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