

Δ -free process

$G(0)$ = empty graph on $[n]$.

$$G(i+1) = G(i) + e_{i+1}$$

random edge

$E(i)$ = edge set of $G(i)$

$$C(i) = \{ f \in \binom{[n]}{2} \setminus E(i) : \Delta \subseteq G(i) + f \}$$

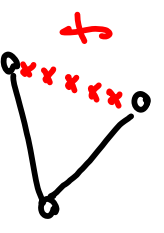
"closed pairs"

$$O(i) = \binom{[n]}{2} \setminus (E(i) \cup C(i))$$

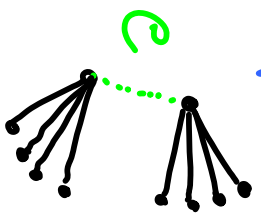
"open pairs"

f is

closed



e is open



e_{i+1} is chosen uniformly at random from $O(i)$.

Set $M = \#$ of steps in the process

note: M is a random variable.

Questions: 1. what is the likely value of M .

2. what is the likely independence # of $G(M)$?

Heuristic:

$G(n)$ resembles $G_{n,i}$

Notation: $G_{n,m}$ has n vertices, m uniform random edges

$$\Pr(xy \in E(G)) \approx \left(1 - \left(\frac{2i}{n}\right)^2\right)^n \approx \exp\left\{-4i^2/n^3\right\}$$

edge probability.

$$E[|O(i)|] \approx \frac{n^2}{2} \exp\left\{-4i^2/n^3\right\}$$

So, the process should end when $i^2/n^3 = \Theta(\log n)$,

which is when $i = \Theta(n^{3/2} \sqrt{\log n})$.

Thm (B, 2009) with high probability

$$M = \Theta(n^{3/2} \sqrt{\log n}) \quad \text{and} \quad \alpha(G_M) = O(n^{1/2} \sqrt{\log n})$$

conjectured by Spencer

note that proving

this gives the Corollary as well as the upper bound on M (since each neighborhood is an independent set).

Corollary $R(B, t) = \Theta(t^2 / \log t)$

note: (i) upper bound on $R(3,t)$ due to Ajtai, Komlos, Szemerédi; (improved by Shearer).

(ii) lower bound previously established by Jeong Han Kim 1995. This was an application of the Rödl nibble.

(iii) See Also Erdős, Suen, Winkler and Bollobás, Riordan and Osthus, Taraz

Random variables

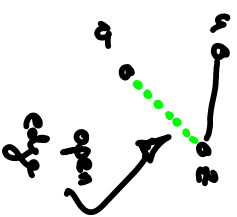
$Q(i) = \#$ of open edges in $G(i)$

This is the variable we care about most as $Q(N) = 0$

$$E[Q(i+1) - Q(i) | \mathcal{F}_i] = \sum_{uv \in Q(i)} \frac{1}{|Q(i)|} [(\# \text{ of edges closed we add } uv = e_{i+1}) + 1]$$

Say z is **partial** w.r.t. uv at step i if

- 1 of uz, vz is in $E(i)$
- 1 of uv, vz is in $Q(i)$



z is **partial** w.r.t. uv .

$Y_{uv}(i) = |\{z : z \text{ is partial w.r.t. } uv \text{ at step } i\}|$

note: $E[Q(i+1) - Q(i) | \mathcal{F}_i] = \sum_{uv \in Q(i)} \frac{Y_{uv}(i)}{Q(i)}$

abuse of notation: Q, Y etc are sets and numbers

$$E[Y_{uv}(i+1) - Y_{uv}(i) | \mathcal{F}_i]$$

$$= - \sum_{z \in Y_{uv}} \frac{1}{Q} (Y_{z z^*}(i) + 1) + ?$$

we can
get new
partial
realizes

notational
convention:

closing
 $z z^*$

choosing $z z^*$

If z is partial w.r.t. uv then let
 $z^* \in \{u, v\}$ s.t. $z z^*$ is open.

Say z is **open** w.r.t. uv in $G(i)$ if uz and
 vz are open

$$X_{uv}(i) = |\{z : z \text{ is open w.r.t. } uv \text{ in } G(i)\}|$$

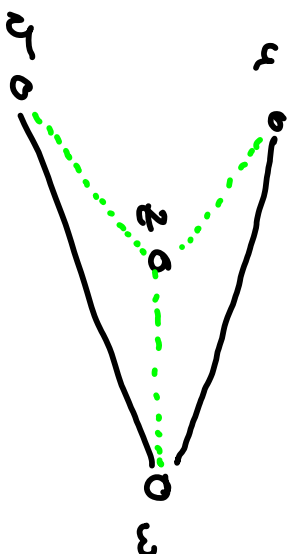
$$\text{then } ? = \frac{2X_{uv}(i)}{Q}$$

$$E[X_{uv}(i+1) - X_{uv}(i) | \mathcal{F}_i] = - \sum_{z \in X_{uv}} \frac{Y_{z z^*}(i) + 1}{Q}$$

is this exactly correct?

what about this situation:

we "count" z twice, but when we choose it only z learns $X_{u,v}$.



observations: (i) u,v is already closed. Do we need to know $X_{u,v}$ for u,v closed? for u,v an edge?

(ii) these vertices w should be rare.

Say z is **closed** w.r.t. u,v in $G(i)$ if uz and vz are edges.

$W_{u,v}(i) = |\{z : z \text{ is closed w.r.t. to } u,v \text{ in } G(i)\}|$

we will not track $W_{u,v}(i)$. We will only bound it.

Trajectories

we set $t = i/n^{3/2}$.

(this should seem natural, given the discussion above)

we would like to show

$$Q(i) \approx p(t)n^2$$

$$Y_{\text{usr}}(i) \approx y(t)n^{1/2}$$

$$X_{\text{usr}}(i) \approx x(t)n$$

} note that these seem sensible. If we have a random graph with about $n^{3/2}$ edges and we add $\epsilon n^{3/2}$ random edges then Q changes by about $\epsilon n^{3/2} \cdot n^{1/2}$ etc.

we use the heuristic to "guess" values of p, y, x .

$$\Pr(\text{usr} \in Q(i)) \approx \left(1 - \left(\frac{2i}{n^2}\right)^2\right)^n = \exp\left\{-4i^2/n^3\right\}$$

$$\text{Set } p(t) = \frac{1}{2} \exp\{-4t^2\}$$

$$Pr(u \in E(i)) \approx \frac{z_i}{n^2} = 2t n^{-1/2}$$

Set

$$y(t) = 4t \exp\{-4t^2\}$$

$$x(t) = \exp\{-8t^2\}$$

We now confirm that these make sense with respect to the expected changes.

Q:

$$E[Q(i+1) - Q(i) | \mathcal{F}_i] = - \sum_{u \in Q(i)} \frac{Y_{u \in Q(i)+1}}{Q(i)} \approx -y(t) n^{1/2}$$

So we should have

$$y' = -y, \quad y(0) = 1/2$$

Note that the 1-step change in $y(t)n^{1/2}$ is about $y'(t)n^{1/2}$

Y:

$$\begin{aligned} E[Y_{uns}(i+1) - Y_{uns}(i) | \mathcal{F}_i] &= - \sum_{z \in Y_{uns}} \frac{Y_{z, z^*}(i) + 1}{Q(i)} + \sum_{z \in X_{uns}} \frac{2}{Q(i)} \\ &\approx - \frac{(y(t) n^{1/2})^2}{g(t) n^2} + \frac{2\chi(t)n}{g(t)n^2} \\ &= - \frac{y^2(t)}{g(t)} \cdot \frac{1}{n} + \frac{2\chi(t)}{g(t)} \cdot \frac{1}{n} \end{aligned}$$

So we
should
have

$$y' = -\frac{y^2}{g} + 2\frac{\chi}{g}, \quad y(0) = 0$$

X:

$$E [X_{n+1}(t_i) - X_{n+1}(t_i) | \mathcal{F}_i] \approx - \sum_{z \in X_{n+1}(t_i)} \frac{Y_{z_n}(t_i) + Y_{z_{n+1}}(t_i) + 2}{Q(t_i)} \approx - \frac{2x(t)y(t)}{g(t)n^2} n^{3/2}$$

So, we are interested in

$$x' = -\frac{2xy}{g}, \quad x(0) = 1$$

The
system

$$\begin{aligned} g(t) &= e^{-4t^2/2} \\ y(t) &= 4t e^{-4t^2} \\ x(t) &= e^{-8t^2} \end{aligned}$$

satisfies all
of the
desired
conditions.

Claim 1: \exists a constant $\mu > 0$ s.t. w.h.p. we have

$$Q(i) = g(t)n^2 \pm f(t)n^{\alpha}$$

$$Y_{us}(i) = y(t)n^{\alpha/2} \pm g(t)n^{\beta} \quad \forall us \notin E(i)$$

$$X_{us}(i) = x(t)n \pm h(t)n^{\rho} \quad \forall us \notin E(i)$$

$$W_{us}(i) < \log^2 n$$

$$\text{for } i=1, \dots, n^{3/2} \sqrt{\log n}$$

as usual, $\alpha, \beta, \rho, f, g, h$
TBA

Proof:

Let the stopping time T be first step at which any of these are violated.

For each $us \in \binom{[n]}{2}$ let $T_{us} = T \wedge \{i : e_i = us\}$
We stop the us -variables at T_{us} .

note: this claim says $G(i)$ indeed resembles $G_{n,i}$.

where $t = 1/n^{3/2}$

We treat the four statements in turn.

Q(i)

This has no appeal to Azuma-Hoeffding or probability theory.

If $i < \mu n^{3/2} \sqrt{\log n}$, then

$$Q(i) = \binom{n}{2} - i - \sum_{j=1}^i Y_{e_j} (j-1)$$

$$= \frac{n^2}{2} - \sum_{j=0}^{i-1} g(j/n^{3/2}) n^{3/2} \pm g(t) n^\beta \pm O(n^{3/2} \sqrt{\log n})$$

$$= \frac{n^2}{2} - n^2 \int_0^t g(s) ds \pm n^{\beta+3/2} \int_0^t g(s) ds \pm O(n^{3/2} \sqrt{\log n})$$

$$= n^2 \left[\frac{1}{2} - \int_0^t 4t e^{-4t^2} dt \right] \pm n^{\beta+3/2} \int_0^t g(s) ds \pm O(n^{3/2} \sqrt{\log n})$$

See

Ex below

This requires

that $f(t) \ll n^{1-\beta}$

$$= n^2 g(t) \pm n^{\beta+3/2} \int_0^t g(s) ds \pm O(n)$$

we need

$$\alpha \geq \beta + 3/2$$

if $\alpha = \beta + 3/2$ \leadsto $f(t) > \int_0^t g(s) ds$

EX

$$\left| \int_{j/n^{3/2}}^{j/n^{3/2} + 1/n^{3/2}} g(s) ds - \frac{g(j/n^{3/2})}{n^{3/2}} \right|$$

why?

$$\leq \int_{j/n^{3/2}}^{j/n^{3/2} + 1/n^{3/2}} \frac{|g'(s)|}{n^{3/2}} ds$$

so $\sum g(j/n^{3/2})$

$$= n^{3/2} \int g + \int_0^t |g'(s)| ds$$

$$= \left| \int_{j/n^{3/2}}^{j/n^{3/2} + 1/n^{3/2}} \underbrace{g(s) - g(j/n^{3/2})}_{\int_s^{j/n^{3/2}} g'(r) dr} ds \right|$$

$X_{uv}(i)$

We begin with the upper bound. Define

$$Z(i) = X_{uv}(i) - x(t)n - h(t)n^p$$

The supermartingale condition

$$E[Z(i+1) - Z(i) | \mathcal{F}_i]$$

converges: if $e_i = uN$
 Then $Z(i) = Z(i-1)$

assumes $h'' n^{p-1} = O(1)$ and
 uses the bound on $W(i)$

$$\leq - \sum_{z \in X_{uv}} \frac{Y_{uz} + Y_{vz}}{Q} - x'(t)n^{-1/2} - h'(t)n^{p-3/2} + O(n^{-1} \log^2 n)$$

$$\leq -2 \frac{(x(t)n + h(t)n^p)(y(t)n^{1/2} + g(t)n^p)}{f(t)n^2 - f(t)n^a} - x'(t)n^{-1/2} - h'(t)n^{p-3/2} + O(n^{-1} \log^2 n)$$

$$\leq - \frac{2xy}{f} n^{-1/2} + \frac{2hy}{f} n^{p-3/2} + \frac{2gx}{f} n^{p-1} + \frac{2hg}{f} n^{p+\beta-2} + \frac{4xyf}{f^2} n^{a-5/2}$$

$$\left| \frac{A}{B \pm \varepsilon} - \frac{A}{B} \right| = \left| \frac{\mp \varepsilon A}{(B \pm \varepsilon) B} \right| < \frac{3/2 \varepsilon A}{B^2}$$

assuming $\varepsilon < B/3$

(we assure that no error term exceeds 5/6 of the main term)

The orange terms cancel exactly.

$$-x'(t) n^{-1/2} - h'(t) n^{\rho-3/2} + O(n^{-1} \log^2 n)$$

we need $h' > \frac{2hy}{f} + \frac{2gx}{f} n^{\frac{1}{2}-\rho+\beta} + \frac{4xyf}{f^2} n^{-1-\rho+\alpha}$

we set $\rho = \alpha - 1$ $\beta = \alpha - 3/2$

with foresight

then we need $h' > 16th + 8fg + 6ytf$ →

this will require $h > e^{8t^2}$

Since we stop at $t = \mu n \sqrt{\log n}$ (you small)
 such an h would be fine.

We apply the variation on Azuma-Hoeffding from

the homework. Note

$$Z_{(i+1)} \geq Z_{(i)} - O(\sqrt{n})$$

$$Z_{(i+1)} \leq Z_{(i)} + O(\sqrt{n})$$

Setting $N = O(\sqrt{n})$

$$\eta = O(1/\sqrt{n})$$

we see that

Each edge we close effects at most 1 vertex in $X_{u,v}(i)$.
 $\therefore X_{u,v}(i+1) \geq X_{u,v}(i) - 1$
 $- Y_{e_{i+1}}(i)$

we achieve this when $X_{u,v}(i+1) = X_{u,v}(i)$

note: we have the condition $n^p < \eta \cdot (n^{3/2} \sqrt{\log n})$ that is needed for the Lemma.

$$\Pr \left(T < \mu n^{3/2} \sqrt{\log n} \text{ and } Z_{u,v}(T) \geq 0 \right)$$

$$\leq \exp \left\{ - \Omega \left(\frac{n^{2p}}{n^{3/2} \sqrt{\log n}} \right) \right\}$$

We need $p > 3/4$ 

note: we can multiply this by $\binom{n}{2}$ and it is still good.

$Y_{u,v}(i)$

Lower bound.

We cannot do the usual thing here. If we consider

$$Z_{u,v}(i) = Y_{u,v}(i) - y^{(t)} n^{1/2} - g^{(t)} n^\beta$$

Then we could have $Z^{(i+1)} > Z^{(i)} + 1/2$ or $Z^{(i+1)} < Z^{(i)} - 1/2$.

So we would take $\eta, N \gg \Omega(1)$. The smallest variation we could keep would be about $n^{3/4}$, this is too big.

Idea: Both additions to and removals from $Y_{u,v}$ are rare. we should treat them separately.

Note that we can write $y^{(t)} = -a(t) + b(t)$ where

$$a(t) = \int_0^t \frac{y^2(s)}{g(s)} ds \quad b(t) = \int_0^t \frac{2x(s)}{g(s)} ds$$

We track removals from $Y_{u,v}$ relative to $a(t)$ and additions to $Y_{u,v}$ relative to $b(t)$.

upper bound - removals

Set $A_{u,v}(i)$ be the # of vertices removed from $Y_{u,v}$ through i steps.

Define

$$Z(i) = A_{u,v}(i) - a(t)n^{\frac{1}{2}} - \frac{g(t)n^\beta}{2}$$

we split the above error evenly between additions and removals.

Supermartingale

$$E[Z(i+1) | \mathcal{F}_i]$$

$$\leq \sum_{z \in Y_{u,v}(i)} \frac{1 + Y_{z,z^*}(i) - 1}{Q(i)} - a'(t)n^{-1} - \frac{g'(t)n^{\beta-3/2}}{2} + O(n^{-2})$$

convention: if $e_i = u,v$
Then $Z(i) = Z(i-1)$

The 1 is for choosing $z z^*$

The -1 is because we don't change Z if $e_{i+1} = u,v$ (note that $u,v \in Y_{z,z^*}$)

Cancels
 The main
 term
 exactly

$$\leq \frac{(y(t)n^{1/2} + g(t)n^\beta)^2}{g(t)n^2 - f(t)n^\alpha} - \frac{y^2(t)}{g(t)} n^{-1} - \frac{g'(t)}{2} n^{\beta-3/2} + O(n^{-3/2})$$

$$\leq \frac{2gy}{f} n^{\beta-3/2} + \frac{2fy^2}{f^2} n^{\alpha-3} + \frac{g^2}{f} n^{2\beta-2} - \frac{g'(t)}{2} n^{\beta-3/2} + O(n^{-3/2})$$

These are the main terms.

Since we already set $\alpha = \beta + 3/2$ we need

again we see that we are getting

$$g' > 16t^2 g + 256t^2 f$$

error = $e^{-\Theta(t^2)}$



Now we apply the modified Azuma tho. along with

$$N = O(\log^2 n)$$

$$\eta = O(1/n)$$

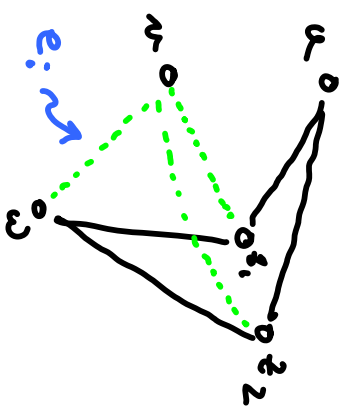
we see that

$$\Pr(T < \mu_n^{3/2} \sqrt{\log n} \wedge A_{\text{unr}}(T) > 0)$$

note:

$$n^\beta < n^{3/2} \sqrt{\log n} \cdot \eta \leq \exp\left\{-\Omega\left(\frac{n^{2\beta}}{n^{1/2} \log^{5/2} n}\right)\right\}$$

So, we need $\beta > 1/4$ 



All the z_i 's are in $W_{u,w}$

note: if we "allowed" us to be added in this calculation then we could get $n^{1/2}$ things removed from Y_{unr} in 1 step.

lower bound - additions

Set $B_{u,v}(i)$ be the # of vertices added to $Y_{u,v}$ through i steps.

Define

$$Z(i) = B_{u,v}(i) - b(t)n^{1/2} + \frac{g(t)n^\beta}{2}$$

submartingale

$$E[Z(i+1) - Z(i) | \mathcal{F}_i]$$

$$\geq \sum_{z \in X_{u,v}} \frac{2}{Q(i)} - b'(t)n^{-1} + \frac{g'(t)n^\beta}{2} + O(n^{-2})$$

Cancel

The main term exactly

$$\geq \frac{2(\alpha(t)n - h(t)n^\beta)}{g(t)n^2 + f(t)n^\alpha} - \frac{2\alpha(t)}{g(t)}n^{-1} + \frac{g'(t)n^\beta}{2} + O(n^{-2})$$

$$\geq -\frac{2\kappa}{f} n^{\rho-2} - \frac{4f\kappa}{f^2} n^{\alpha-3} + \frac{g'}{2} n^{\beta-3/2} + O(n^{-2})$$

Since we are setting $\rho = \beta + 1/2$, $\alpha = \beta + 3/2$ we need.

$$g' > \frac{8\kappa}{f} + 32f \rightarrow$$

Now we apply the modified Azuma - Hoffding with

$$N = O(1) \text{ and } \eta = O(1/n)$$

we see that

$$P(T < \mu n^{3/2} \sqrt{\log n} \text{ and } B_{\mu n}(T) < 0)$$

$$\leq \exp\left\{-\Omega\left(\frac{n^{2\beta}}{n^{1/2} \sqrt{\log n}}\right)\right\}$$

So, we need

$$\beta > 1/4 \rightarrow$$

again,
 $\eta \cdot n^{3/2} \sqrt{\log n} > n^\beta$
 and we
 can multiply
 by (2) with
 no problem

reconciliation

Variable condition →

$$Q: f' > g$$

$$X: h' > 16th + 8fg + 64tg - f$$

$$Y \left\{ \begin{array}{l} g' > 16tg + 256t^2f \\ g' > \frac{8h}{g} + 32f \end{array} \right.$$

first guess

$$\left\{ \begin{array}{l} f = e At^2 \\ g = e Bt^2 \\ h = e Ct^2 \end{array} \right.$$

looking at the exponents:

$$\left. \begin{array}{l} A \geq B \\ C \geq B-4, A-4 \\ B \geq A, C+4 \end{array} \right\}$$

we should take

$$B = A, C = A - 4$$

second guess

$$\begin{cases} f = u(t)e^{At^2} \\ g = v(t)e^{At^2} \\ h = w(t)e^{(A-u)t^2} \end{cases}$$

we will choose

A large

then we need

$$2Atu + u' > v$$

$$2(A-u)tw + w' > \underline{16}tw + 4v + 32tu$$

$$2Atv + v' > \underline{16}tv + 256t^2u$$

$$2Atw + w' > 16w + 16u$$

note: (i) we can ignore the underlined terms by taking A large.

(ii) when t is large we will likely want $v \geq tu$ and $w \geq u$

third guess

$$\begin{cases} f = u(t)e^{At^2} \\ g = u(t) \cdot (t+1)e^{At^2} \\ h = u(t)e^{(A-u)t^2} \end{cases}$$

$$2Atu + u' > (t+1)u$$

$$2(A-4)t u + u' > 16tu + 4(t+1)u + 32tu$$

$$2At(t+1)u + u + u + (t+1)u' > 16t(t+1)u + 256t^2u$$

$$2At(t+1)u + u + (t+1)u' > 16u + 16u$$

note: we can ignore these terms
by taking A large

It seems that we are left with

$$u' > u$$

$$u' > 4u$$

So we take $u = e^{32t}$

$$u' > 31u$$

The following error functions suffice.

$$f = e^{At^2 + 32t}$$

$$g = (t+1)e^{At^2 + 32t}$$

$$h = e^{(A-4)t^2 + 32t}$$

$W_{u,r}(i)$

We apply the first moment method. If $i < T$

$$Pr(W_{u,r}(i+1) = W_{u,r}(i) + 1) = \frac{Y_{u,r}(i)}{Q(i)} < 2 \frac{4t \cdot e^{-4t^2} \cdot n^{1/2}}{(e^{-4t^2/2}) \cdot n^2} = 16t n^{-3/2}$$

So long as t is not very close to 0.

$Pr(T < \mu n^{3/2} \sqrt{\log n})$ and \exists u.r. s.t. $W_{u,r}(T) > \log^2 n$

EX

$$\begin{aligned}
 &< \binom{n}{2} \left(\mu n^{3/2} \sqrt{\log n} \right) (16 \mu \sqrt{\log n} n^{-3/2})^{\log^2 n} \\
 &\leq \frac{n^2}{2} \left[\frac{\mu n^{3/2} \sqrt{\log n}}{\log^2 n} \cdot e \cdot \frac{16 \mu \sqrt{\log n}}{n^{3/2}} \right]^{\log^2 n}
 \end{aligned}$$

this estimate holds even for t near 0

□ Claim 1 is proved.

Note:

Claim 1 establishes the lower bound of the main theorem, as it shows that $Q(i) > 0$ for $i = \mu n^{3/2} \sqrt{\log n}$.