

The giant component via
dynamic concentration

Def: $C_1 = C_1(G)$ is the largest component in graph G .

Theorem (Erdős, Rényi 1960)

Let $p = c/n$ where c is a constant

- (i) If $c < 1$ then $|C_1(G_{n,p})| = O(\log n)$
- (ii) If $c > 1$ then $|C_1(G_{n,p})| = \Omega(n)$

with high probability

Random Graph Process

Let $u_1, v_1, u_2, v_2, u_3, v_3, \dots$

be chosen independently and u.a.r. from $[n]$.

Set $e_i = \{u_i, v_i\}$ for $i=1, 2, \dots$

$G(i) =$ graph on $[n]$ with edge set $\{e_1, \dots, e_i\}$.

Suppose $G(i)$ has connected components C_1, C_2, C_3, \dots

note:

1. Loops and multiple edges can occur.

2. $G(i)$ is very similar to $G_{n,p}$ with $i = np/2$

$X(i) = \sum |c_k|^2$ is the susceptibility of $G(i)$

$$E[X(i+1) - X(i) | G(i)] = \sum_{k \neq j} \frac{|c_k|^2 |c_j|^2}{n} = 2 \sum_k |c_k|^4$$

$$= \frac{2}{n^2} \left(\sum_k |c_k|^2 \right)^2 - \frac{2}{n^2} \sum_k |c_k|^4$$

$$= \frac{2}{n^2} X^2$$

should be
small
call this *

So, we should have

$$X(i) \approx n \alpha(i/n)$$

where

$$x(0) = 1$$

$$\frac{dx}{dt} = 2x^2$$

$$\Rightarrow x(t) = \frac{1}{1-2t}$$

Theorem (B-Kravitz Spencer-Warmald).

Let $\varepsilon > 0$.

With high probability

$$X(i) = \frac{n}{1-2t} (1 + o(1)),$$

uniformly, for $i = 0, 1, \dots, (\frac{1}{2} - \varepsilon)n$.
(where $t = i/n$).

First try at a proof (upper bound)

$$\text{Set } Z(i) = Z^+(i) = X(i) - nX(t) - e(t)$$

a slowly growing function, will also depend on n

note: $Z(0) = -e(0)$

$t = i/n$ this is our time-scaling

- Outline:
1. Show $Z(i)$ is a supermartingale.
 2. Apply Azuma-Hoeffding to conclude.

$$P(\exists i: \text{s.t. } Z(i) > 0) = o(1)$$

Convention: Let the stopping time T be the smallest index i s.t.

$$X(i) \in nX(t) \pm e(t) \quad (1)$$

fails to hold. Set $i_{NT} = \min \{i, T\}$. If $Z(i)$ is a supermartingale then $Z(i_{NT})$ is also a supermartingale.

We work with $Z(i:NT)$, and therefore we can assume that (1) holds.

Convention: Any condition that must be satisfied by a choice of a parameter made later is **stated in red**.

Step 1.

Notation: $\Delta Z = Z(i:1) - Z(i)$

$$E[\Delta Z] \leq \underbrace{\frac{2\chi^2(i)}{n^2}}_{\text{contribution from } \chi(i)} - 2\chi^2(t) - \frac{e'(t)}{n}$$

here we use $\chi'' > 0$ and

$$e'' > 0.$$

$$\leq \frac{2}{n^2} (nx(t) + e(t))^2 - 2x^2(t) - \frac{e'(t)}{n} \quad \text{by (1)}$$

$$\leq \frac{4x(t)e(t)}{n} + \frac{e^2(t)}{n^2} - \frac{e'(t)}{n}$$

to get the supermartingale condition
it suffices to choose $e(t)$ so that

$$e' > \frac{4e}{1-2t} \quad \text{and} \quad e = o(nx)$$

$$\text{We set } e = \frac{S_e}{(1-2t)^3}$$

S_e , the scaling
of the error term,
will be a function
of n

Step 2

To apply Hoeffding-Azuma we need a bound on $|\Delta z|$

e.g. if we assume $|C_k| < B$ (i.e. $B = \log^2 n$ seems sensible)

then $\Delta z = O(B^2)$ and

$$\Pr(z_i - z_0 > A) \leq \exp\left\{-\frac{A^2}{3 \cdot B^4}\right\}.$$

Thus if we take

$$O(\cdot) = S_e = \omega(B^2 \sqrt{n} \log n)$$

this $\log n$ can be any $\omega(1)$ function.

then

$$\Pr(z_i \geq 0) = \Pr(z_i - z_0 > \log n \cdot B \sqrt{n}) < \exp\{-\Omega(\log^2 n)\}$$

So (i) we need a bound $|C_B|$ that holds throughout the process.

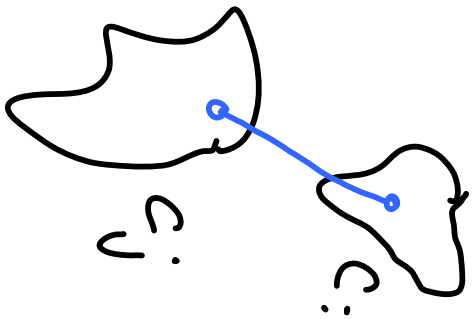
(ii) we also need to set-up a submartingale for the lower bound. And here we need the bound on $|C_B|$ to deal with (ii).

Idea: ignore edges that touch components on n^α or more vertices.

If we do this then

(i) $\Delta X \leq O(n^{2\alpha})$ and so long as $\alpha < 1/4$ the application of Azuma-Hoeffding works.

(ii) There will be an impact on $E[\Delta X]$.



$$\sum_{C_i \text{ bad}, C_j} \frac{|C_i|}{n} \frac{|C_j|}{n} 2|C_i||C_j|$$

recall:
 $X = \sum |C_i|^2$

$$= \sum_{C_i \text{ bad}} \frac{|C_i|^2}{n} \cdot \frac{2 \sum |C_j|^2}{n}$$

if $|C_i|$ or $|C_j|$
 too large add

This as a
 blue edge.

$$\leq \frac{1}{n} \cdot (\# \text{ large})^2 \cdot \frac{2X}{n}$$

constant

In order to maintain the supermartingale condition on $Z(i)$ we need

large components = $O(S_e n^{-2\alpha})$

Next idea: introduce another random variable to keep control on the # of large components.

$$Y(i) = \sum_k |C_k|^3$$

note: ① $Y = O(n) \Rightarrow$ # large components $\leq n^{1-3\alpha}$

Setting $S_\epsilon = n^{1-\alpha/2}$ for a sufficiently small ϵ will satisfy the conditions above.

② $Y = O(n) \Rightarrow$ * $\leq n^{\alpha-1}$. We need this for the lower bound on $X(i)$.

It remains to show that $Y(i)$ is at most linear for $i \leq (\frac{1}{2} - \epsilon)n$.

$$\begin{aligned}
 E[\Delta Y] &= \sum_{i,j} \frac{|c_i|}{n} \frac{|c_j|}{n} (3|c_i|^2 |c_j| + 3|c_i| |c_j|^2) \\
 &\leq \frac{6}{n^2} \left[\sum_i |c_i|^2 \right] \left[\sum_j |c_j|^3 \right] \\
 &= \frac{6XY}{n^2}
 \end{aligned}$$

If we want to show

$$Y(i) \leq n y(i/n)$$

we take the solution of

$$y' = 6xy \quad y(0) = 1$$

so

$$y = \frac{1}{(1-2t)^3}$$

note: as we need only an upper bound we can ignore

① the diagonal term

② the edges we fail to add.

The application of Azure-Health is very similar to the one above.



