Solving semidefinite-quadratic-linear programs using SDPT3

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Abstract

This paper discusses computational experiments with linear optimization problems involving semidefinite, quadratic, and linear cone constraints (SQLPs). Many test problems of this type are solved using a new release of SDPT3, a MATLAB implementation of infeasible primal-dual path-following algorithms. The software developed by the authors uses Mehrotra-type predictor-corrector variants of interior-point methods and two types of search directions: the HKM and NT directions. A discussion of implementation details is provided and computational results on problems from the SDPLIB and DIMACS Challenge collections are reported.

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1 Introduction

Conic linear optimization problems can be expressed in the following standard form:

$$\begin{array}{lll} \min & \langle c,\, x\rangle & \\ \text{s.t.} & \langle a_k,\, x\rangle \;=\; b_k, & k=1,\ldots,m, \\ & x \in K & \end{array} \tag{1}$$

where K is a closed, convex pointed cone in a finite dimensional inner product space endowed with an inner product $\langle\cdot,\,\cdot\rangle$. By choosing K to be the semidefinite, quadratic (second-order), and linear cones respectively, one obtains the well-known special cases of semidefinite, second-order cone, and linear programming problems. Recent years have seen a dramatic increase in the number of subclasses of conic optimization problems that can be solved efficiently by interior-point methods. In addition to the ongoing theoretical work that derived convergence guarantees and convergence rates for such algorithms, many groups of researchers have also implemented these algorithms and developed public domain software packages that are capable of solving conic optimization problems of ever increasing size and diversity. This paper discusses the authors' contribution to this effort through the development of the software SDPT3. Our earlier work on SDPT3 is presented in [20, 23].

The current version of SDPT3, version 3.0, can solve conic linear optimization problems with inclusion constraints for the cone of positive semidefinite matrices, the second-order cone, and/or the polyhedral cone of nonnegative vectors. In other words, we allow K in (1) to be a Cartesian product of cones of positive semidefinite matrices, second-order cones, and the nonnegative orthant. We use the following standard form of such problems, henceforth called SQLP problems:

$$\begin{array}{llccccc} (P) & \min & \sum_{j=1}^{n_s}\langle c_j^s,\, x_j^s\rangle & + & \sum_{i=1}^{n_q}\langle c_i^q,\, x_i^q\rangle & + & \langle c^l,\, x^l\rangle \\ & \text{s.t.} & \sum_{j=1}^{n_s}(A_j^s)^T\mathbf{svec}(x_j^s) & + & \sum_{i=1}^{n_q}(A_i^q)^T x_i^q & + & (A^l)^T x^l \;=\; b, \\ & & x_j^s \in K_s^{s_j}\ \forall j, & & x_i^q \in K_q^{q_i}\ \forall i, & & x^l \in K_l^{n_l}. \end{array}$$

Here, c_j^s, x_j^s are symmetric matrices of dimension s_j and $K_s^{s_j}$ is the cone of positive semidefinite symmetric matrices of the same dimension. Similarly, c_i^q, x_i^q are vectors in $\mathbb{R}^{q_i}$ and $K_q^{q_i}$ is the second-order cone defined by $K_q^{q_i} := \{x \in \mathbb{R}^{q_i} : x_1 \geq \|x_{2:q_i}\|\}$. Finally, c^l, x^l are vectors of dimension n_l and $K_l^{n_l}$ is the cone $\mathbb{R}_+^{n_l}$. In the notation above, A_j^s denotes the $\bar{s}_j \times m$ matrix with $\bar{s}_j = s_j(s_j+1)/2$ whose columns are obtained using the **svec** operator from m symmetric $s_j \times s_j$ constraint matrices corresponding to the jth semidefinite block x_j^s. For a definition of the vectorization operator **svec** on symmetric matrices, see, e.g., [20]. The matrices A_i^q's are $q_i \times m$ dimensional constraint matrices corresponding to the ith quadratic block x_i^q, and A^l is the $l \times m$ dimensional constraint matrix corresponding to the linear block x^l. The notation $\langle p,\, q\rangle$ denotes the standard inner product in the appropriate space.

The software also solves the dual problem associated with the problem above:

$$
\begin{array}{lllcll}
(D) & \max & b^T y & & & \\
& \text{s.t.} & A_j^s y & + & z_j^s & = & c_j^s, \quad j = 1 \ldots, n_s \\
& & A_i^q y & + & z_i^q & = & c_i^q, \quad i = 1 \ldots, n_q \\
& & A^l y & + & z^l & = & c^l, \\
& z_j^s \in K_s^{s_j} \ \forall j, & & z_i^q \in K_q^{q_i} \ \forall i, & & z^l \in K_l^{n_l}. &
\end{array}
$$

This package is written in MATLAB version 5.3 and is compatible with MATLAB version 6.0. It is available from the internet sites:

`http://www.math.nus.edu.sg/~mattohkc/index.html`

`http://www.math.cmu.edu/~reha/sdpt3.html`

This software package was originally developed to provide researchers in semidefinite programming with a collection of reasonably efficient and robust algorithms that can solve general SDPs with matrices of dimensions of the order of a hundred. The current release, version 3.0, expands the family of problems solvable by the software in two dimensions. First, this version is much faster than the previous release [23], especially on large sparse problems, and consequently can solve much larger problems. Second, the current release can also directly solve problems that have second-order cone constraints — with the previous version it was necessary to convert such constraints to semidefinite cone constraints.

In this paper, the vector 2-norm and Frobenius norm are denoted by $\|\cdot\|$ and $\|\cdot\|_F$, respectively. In the next section, we discuss the algorithm used in the software and several implementation details. Section 3 describes the initial iterates generated by our software while Section 4 briefly explains how to use the software and its data storage scheme. In Section 5, we present and comment on the results of our computational experiments with our software on problems from the SDPLIB and DIMACS libraries.

2 A primal-dual infeasible-interior-point algorithm

The algorithm implemented in SDPT3 is a primal-dual interior-point algorithm that uses the path-following paradigm. In each iteration, we first compute a *predictor* search direction aimed at decreasing the duality gap as much as possible. After that, the algorithm generates a Mehrotra-type corrector step [14] with the intention of keeping the iterates close to the central path. However, we do not impose any neighborhood restrictions on our iterates.[1] Initial iterates need not be feasible — the

[1]This strategy works well on most problems we tested. However, it should be noted that the occasional failure of the software on problems with poorly chosen initial iterates is likely due to the lack of a neighborhood enforcement in the algorithm.

algorithm tries to achieve feasibility and optimality of its iterates simultaneously.

It should be noted that our implementation allows the user to switch to a primal-dual path-following algorithm that does not use corrector steps and sets a centering parameter to be used in such a framework. The choices we make on the parameters used by the algorithm are based on minimizing either the number of iterations or the CPU time of the linear algebra involved in computing the Schur complement matrix and its Cholesky factorization. What follows is a pseudo-code for the algorithm we implemented. Note that this description makes references to later parts of this section where many details related to the algorithm are explained.

Algorithm IPC. *Suppose we are given an initial iterate (x^0, y^0, z^0) with x^0, z^0 strictly satisfying all the conic constraints. Decide on the type of search direction to use. Set $\gamma^0 = 0.9$. Choose a value for the parameter* expon *used in e.*

For $k = 0, 1, \ldots$

(Let the current and the next iterate be (x, y, z) and (x^+, y^+, z^+) respectively. Also, let the current and the next step-length parameter be denoted by γ and γ^+ respectively.)

- *Set $\mu = \langle x,\, z\rangle/n$, and*

$$\texttt{relgap} \;=\; \frac{\langle x,\, z\rangle}{1 + \max(|\langle c,\, x\rangle|, |b^T y|)}, \quad \phi \;=\; \max\left(\frac{\|r_p\|}{1+\|b\|}\,,\, \frac{\|R_d\|}{1+\|c\|}\right). \tag{2}$$

  *Stop the iteration if the infeasibility measure ϕ and the relative duality gap (*relgap*) are sufficiently small.*

- *(Predictor step)*
  Solve the linear system (10), with $\sigma = 0$ in the right-side vector (12). Denote the solution of (4) by $(\delta x, \delta y, \delta z)$. Let α_p and β_p be the step-lengths defined as in (33) and (34) with $\Delta x, \Delta z$ replaced by $\delta x, \delta z$, respectively.

- *Take σ to be*

$$\sigma \;=\; \min\left(1,\, \left[\frac{\langle x + \alpha_p\,\delta x,\, z + \beta_p\,\delta z\rangle}{\langle x,\, z\rangle}\right]^e\right),$$

  where the exponent e is chosen as follows:

$$e \;=\; \begin{cases} \max[expon, 3\min(\alpha_p, \beta_p)^2] & \text{if } \mu > 10^{-6}, \\ expon & \text{if } \mu \le 10^{-6}. \end{cases}$$

- *(Corrector step)*
  Solve the linear system (10) with R_c in the the right-hand side vector (12) replaced by

$$\begin{aligned} \tilde{R}_c^s &= \mathbf{svec}\left[\sigma\mu I - H_P(\mathbf{smat}(x^s)\mathbf{smat}(z^s)) - H_P(\mathbf{smat}(\delta x^s)\mathbf{smat}(\delta z^s))\right] \\ \tilde{R}_c^q &= \sigma\mu e^q - T_G(x^q, z^q) - T_G(\delta x^q, \delta z^q) \\ \tilde{R}_c^l &= \sigma\mu e^l - diag(x^l) z^l - diag(\delta x^l)\delta z^l. \end{aligned}$$

  Denote the solution of (4) by $(\Delta x, \Delta y, \Delta z)$.

- *Update* (x, y, z) *to* (x^+, y^+, z^+) *by*

$$x^+ \;=\; x + \alpha\,\Delta x, \quad y^+ \;=\; y + \beta\,\Delta y, \quad z^+ \;=\; z + \beta\,\Delta z,$$

where α *and* β *are computed as in (33) and (34) with* γ *chosen to be* $\gamma = 0.9 + 0.09\min(\alpha_p, \beta_p)$.

- *Update the step-length parameter by*

$$\gamma^+ \quad = \quad 0.9 + 0.09\min(\alpha, \beta).$$

2.1 The search direction

To simplify discussion, we introduce the following notation, which is also consistent with the internal data representation in SDPT3:

$$A^s = \begin{bmatrix} A_1^s \\ \vdots \\ A_{n_s}^s \end{bmatrix}, \quad A^q = \begin{bmatrix} A_1^q \\ \vdots \\ A_{n_q}^q \end{bmatrix}.$$

Similarly, we define

$$x^s = \begin{bmatrix} \mathbf{svec}(x_1^s) \\ \vdots \\ \mathbf{svec}(x_{n_s}^s) \end{bmatrix}, \quad x^q = \begin{bmatrix} x_1^q \\ \vdots \\ x_{n_q}^q \end{bmatrix}. \tag{3}$$

The vectors c^s, z^s, c^q, and z^q are defined analogously. We will use corresponding notation for the search directions as well. Finally, let

$$A^T = \begin{bmatrix} A^s \\ A^q \\ A^l \end{bmatrix}, \; x = \begin{bmatrix} x^s \\ x^q \\ x^l \end{bmatrix}, \; c = \begin{bmatrix} c^s \\ c^q \\ c^l \end{bmatrix}, \; z = \begin{bmatrix} z^s \\ z^q \\ z^l \end{bmatrix},$$

and

$$n = \sum_{j=1}^{n_s} s_j \;+\; n_q \;+\; n_l.$$

With the notations introduced above, the primal and dual equality constraints can be represented respectively as

$$Ax = b, \quad A^T y + z = c.$$

In this paper, we assume that A has full row rank. In our software, dependent constraints are automatically removed, if there are any.

The main step at each iteration of our algorithms is the computation of the search direction $(\Delta x, \Delta y, \Delta z)$ from the *symmetrized Newton equation* with respect to an

invertible block diagonal scaling matrix P for the semidefinite block and a block scaling matrix G for the quadratic block. The matrices P and G are usually chosen as a function of the current iterate x, z and we will elaborate on specific choices below. The search direction $(\Delta x, \Delta y, \Delta z)$ is obtained from the following system of equations:

$$
\begin{array}{rcrcrcl}
A^T\Delta y & + & \Delta z & = & R_d & := & c - z - A^T y \\
A\Delta x & & & = & r_p & := & b - Ax \\
\mathcal{E}^s\Delta x^s & + & \mathcal{F}^s\Delta z^s & = & R_c^s & := & \mathbf{svec}\left(\sigma\mu I - H_P\left(\mathbf{smat}(x^s)\mathbf{smat}(z^s)\right)\right) \\
\mathcal{E}^q\Delta x^q & + & \mathcal{F}^q\Delta z^q & = & R_c^q & := & \sigma\mu e^q - T_G(x^q, z^q) \\
\mathcal{E}^l\Delta x^l & + & \mathcal{F}^l\Delta z^l & = & R_c^l & := & \sigma\mu e^l - \mathcal{E}^l\mathcal{F}^l e^l,
\end{array}
\tag{4}
$$

where $\mu = \langle x,\, z\rangle/n$ and σ is the centering parameter. The notation **smat** denotes the inverse map of **svec** and both are to be interpreted as blockwise operators if the argument consists of blocks. Here H_P is the symmetrization operator whose action on the jth semidefinite block is defined by

$$
\begin{aligned}
H_{P_j} : I\!R^{s_j\times s_j} &\longrightarrow I\!R^{s_j\times s_j} \\
H_{P_j}(U) &= \tfrac{1}{2}\left[P_j U P_j^{-1} + P_j^{-T} U^T P_j^T\right],
\end{aligned}
\tag{5}
$$

with P_j the jth block of the block diagonal matrix P and $\mathcal{E}^s$ and $\mathcal{F}^s$ are symmetric block diagonal matrices whose jth blocks are given by

$$
\mathcal{E}_j^s = P_j \circledast P_j^{-T} z_j^s, \qquad \mathcal{F}_j^s = P_j x_j^s \circledast P_j^{-T}, \tag{6}
$$

where $R \circledast T$ is the symmetrized Kronecker product operation described in [20].

In the quadratic block, e^q denotes the blockwise identity vector, i.e.,

$$
e^q = \begin{bmatrix} e_1^q \\ \vdots \\ e_{n_q}^q \end{bmatrix},
$$

where e_j^q is the first unit vector in $I\!R^{q_j}$. Let the arrow operator defined in [3] be denoted by $\mathbf{Arw}(\cdot)$. Then the operator $T_G(x^q, z^q)$ is defined as follows:

$$
T_G(x^q, z^q) = \begin{bmatrix} \mathbf{Arw}\left(G_1 x_1^q\right)\left(G_1^{-1} z_1^q\right) \\ \vdots \\ \mathbf{Arw}\left(G_{n_q} x_{n_q}^q\right)\left(G_{n_q}^{-1} z_{n_q}^q\right) \end{bmatrix}, \tag{7}
$$

where G is a symmetric block diagonal matrix that depends on x, z and G_i is the ith block of G. The matrices $\mathcal{E}^q$ and $\mathcal{F}^q$ are block diagonal matrices whose the ith blocks are given by

$$
\mathcal{E}_i^q = \mathbf{Arw}\left(G_i^{-1} z_i^q\right) G_i, \quad \mathcal{F}_i^q = \mathbf{Arw}\left(G_i x_i^q\right) G_i^{-1}. \tag{8}
$$

In the linear block, e^l denotes the n_l-dimensional vector of ones, and $\mathcal{E}^l = \mathrm{diag}(x^l)$, $\mathcal{F}^l = \mathrm{diag}(z^l)$.

For future reference, we partition the vectors R_d, Δx, and Δz in a manner analogous to c, x, and z as follows:

$$R_d = \begin{bmatrix} R_d^s \\ R_d^q \\ R_d^l \end{bmatrix}, \quad \Delta x = \begin{bmatrix} \Delta x^s \\ \Delta x^q \\ \Delta x^l \end{bmatrix}, \quad \Delta z = \begin{bmatrix} \Delta z^s \\ \Delta z^q \\ \Delta z^l \end{bmatrix}. \tag{9}$$

Assuming that $m = \mathcal{O}(n)$, we compute the search direction via a Schur complement equation as follows (the reader is referred to [2] and [20] for details). First compute Δy from the Schur complement equation

$$M\Delta y \;=\; h, \tag{10}$$

where

$$\begin{aligned} M &= (A^s)^T(\mathcal{E}^s)^{-1}\mathcal{F}^s A^s + (A^q)^T(\mathcal{E}^q)^{-1}\mathcal{F}^q A^q + (A^l)^T(\mathcal{E}^l)^{-1}\mathcal{F}^l A^l & (11)\\ h &= r_p \;-\; (A^s)^T(\mathcal{E}^s)^{-1}(R_c^s - \mathcal{F}^s R_d^s) \\ &\quad - (A^q)^T(\mathcal{E}^q)^{-1}(R_c^q - \mathcal{F}^q R_d^q) \;-\; (A^l)^T(\mathcal{E}^l)^{-1}(R_c^l - \mathcal{F}^l R_d^l). & (12) \end{aligned}$$

Then compute Δx and Δz from the equations

$$\begin{aligned} \Delta z &= R_d - A^T\Delta y & (13)\\ \Delta x^s &= (\mathcal{E}^s)^{-1}R_c^s - (\mathcal{E}^s)^{-1}\mathcal{F}^s\Delta z^s & (14)\\ \Delta x^q &= (\mathcal{E}^q)^{-1}R_c^q - (\mathcal{E}^q)^{-1}\mathcal{F}^q\Delta z^q & (15)\\ \Delta x^l &= (\mathcal{E}^l)^{-1}R_c^l - (\mathcal{E}^l)^{-1}\mathcal{F}^l\Delta z^l. & (16) \end{aligned}$$

2.2 Two choices of search directions

We start by introducing some notation that we will use in the remainder of this paper. For a given q_i-dimensional vector x_i^q, we let x_i^0 denote its first component and x_i^1 denote its subvector consisting of the remaining entries, i.e.,

$$\begin{bmatrix} x_i^0 \\ x_i^1 \end{bmatrix} = \begin{bmatrix} (x_i^q)_1 \\ (x_i^q)_{2:q_i} \end{bmatrix}. \tag{17}$$

We will use the same convention for $z_i^q, \Delta x_i^q$, etc. Also, we define the following function from $K_q^{q_i}$ to $\mathbb{R}_+$:

$$\gamma(x_i^q) \;:=\; \sqrt{(x_i^0)^2 - \langle x_i^1,\, x_i^1\rangle}. \tag{18}$$

Finally, we use X and Z for $\mathbf{smat}(x^s)$ and $\mathbf{smat}(z^s)$, where the operation is applied blockwise to form a block diagonal symmetric matrix of order $\sum_{j=1}^{n_s} s_j$.

In the current release of this package, the user has two choices of scaling operators parametrized by P and G, resulting in two different search directions: the HKM direction [10, 12, 15], and the NT direction [18]. See also Tsuchiya [24, 25] for the second-order case.

(1) **The HKM direction.** This choice uses the scaling matrix $P = Z^{1/2}$ for the semidefinite blocks and a symmetric block diagonal scaling matrix G for the quadratic blocks where the ith block G_i is given by the following equation:

$$G_i \;=\; \begin{bmatrix} z_i^0 & (z_i^1)^T \\ z_i^1 & \gamma(z_i^q)I + \dfrac{z_i^1 (z_i^1)^T}{\gamma(z_i^q) + z_i^0} \end{bmatrix}. \tag{19}$$

(2) **The NT direction.** This choice uses the scaling matrix $P = N^{-1}$ for the semidefinite blocks, where N is a matrix such that $D := N^T Z N = N^{-1} X N^{-T}$ is a diagonal matrix [20], and G is a symmetric block diagonal matrix whose ith block G_i is defined as follows. Let

$$\omega_i = \sqrt{\frac{\gamma(z_i^q)}{\gamma(x_i^q)}}, \quad \xi_i = \begin{bmatrix} \xi_i^0 \\ \xi_i^1 \end{bmatrix} = \begin{bmatrix} \frac{1}{\omega_i} z_i^0 + \omega_i x_i^0 \\ \frac{1}{\omega_i} z_i^1 - \omega_i x_i^1 \end{bmatrix}. \tag{20}$$

Then

$$G_i = \omega_i \begin{bmatrix} t_i^0 & (t_i^1)^T \\ t_i^1 & I + \dfrac{t_i^1 (t_i^1)^T}{1 + t_i^0} \end{bmatrix}, \quad \text{where } \begin{bmatrix} t_i^0 \\ t_i^1 \end{bmatrix} = \frac{1}{\gamma(\xi_i)} \begin{bmatrix} \xi_i^0 \\ \xi_i^1 \end{bmatrix}. \tag{21}$$

2.3 Computation of the search directions

The size and the density of the Schur complement matrix M defined in (10) is the main determinant of the cost of each iteration in our algorithm. The density of this matrix depends on two factors: (i) The density of the constraint coefficient matrices A^s, A^q, and A^l, and (ii) any additional fill-in introduced because of the terms $(\mathcal{E}^s)^{-1}\mathcal{F}^s$, $(\mathcal{E}^q)^{-1}\mathcal{F}^q$, and $(\mathcal{E}^l)^{-1}\mathcal{F}^l$ in (10).

2.3.1 Semidefinite blocks

For problems with semidefinite blocks, it appears that there is not much one can do about additional fill-in, since $(\mathcal{E}^s)^{-1}\mathcal{F}^s$ is dense and structure-less for most problems. One can take advantage of sparsity in A^s in related computations, however, and we discussed some of these issues, such as blockwise computations, in our earlier papers [20, 23].

The way we exploit sparsity of A_j^s in the computation of $M_j^s := (A_j^s)^T (\mathcal{E}_j^s)^{-1} \mathcal{F}_j^s A_j^s$ basically follows the approach in [7]. We will not go into the details here but just briefly highlight one issue that is often critical in cutting down the computation

time in forming M_j^s. Let $A_j^s(:,k)$ be the kth column of A_j^s. In computing the kth column of M_j^s, typically a matrix product of the form $x_j^s\,\mathbf{smat}(A_j^s(:,k))\,(z_j^s)^{-1}$ or $w_j^s\,\mathbf{smat}(A_j^s(:,k))\,w_j^s$ is required for the HKM direction or NT direction, respectively. In many large SDP problems, the matrix $\mathbf{smat}(A_j^s(:,k))$ is usually very sparse, and it is important to store this matrix as a sparse matrix in MATLAB and perform sparse-dense matrix-matrix multiplication in the matrix products just mentioned whenever possible. Also, entries of this product only need to be computed if they contribute to an entry of M, i.e., if they correspond to a nonzero entry of $A_j^s(:,k')$ for some k'.

2.3.2 Quadratic and linear blocks

For linear blocks, $(\mathcal{E}^l)^{-1}\mathcal{F}^l$ is a diagonal matrix and it does not introduce any additional fill-in. This matrix does, however, affect the conditioning of the Schur complement matrix and is a popular subject of research in implementations of interior-point methods for linear programming.

From equation (11), it is easily shown that the contribution of the quadratic blocks to the matrix M is given by

$$M^q \;=\; (A^q)^T(\mathcal{E}^q)^{-1}\mathcal{F}^q A^q \;=\; \sum_{i=1}^{n_q} \underbrace{(A_i^q)^T(\mathcal{E}_i^q)^{-1}\mathcal{F}_i^q A_i^q}_{M_i^q}. \tag{22}$$

For the HKM direction, $(\mathcal{E}^q)^{-1}\mathcal{F}^q$ is a block diagonal matrix whose ith block is given by

$$\begin{aligned}(\mathcal{E}_i^q)^{-1}\mathcal{F}_i^q \;&=\; G_i^{-1}\mathbf{Arw}\,(G_i x_i^q)\,G_i^{-1} \\ &=\; \frac{1}{\gamma^2(z_i^q)}\left(\langle x_i^q,\, z_i^q\rangle \begin{bmatrix} -1 & 0 \\ 0 & I \end{bmatrix} + \begin{bmatrix} x_i^0 \\ x_i^1 \end{bmatrix}\begin{bmatrix} z_i^0 \\ -z_i^1 \end{bmatrix}^T + \begin{bmatrix} z_i^0 \\ -z_i^1 \end{bmatrix}\begin{bmatrix} x_i^0 \\ x_i^1 \end{bmatrix}^T\right)\end{aligned} \tag{23}$$

(Note that $\mathbf{Arw}\left(G_i^{-1}z_i^q\right) = I$.) Thus, we see that matrix $(\mathcal{E}_i^q)^{-1}\mathcal{F}_i^q$ in M_i^q is the sum of a diagonal matrix and a rank-two symmetric matrix. Hence

$$M_i^q \;=\; \frac{\langle x_i^q,\, z_i^q\rangle}{\gamma^2(z_i^q)}(A_i^q)^T J_i A_i^q \;+\; u_i^q(v_i^q)^T + v_i^q(u_i^q)^T, \tag{24}$$

where

$$J_i = \begin{bmatrix} -1 & 0 \\ 0 & I \end{bmatrix}, \quad u_i^q = (A_i^q)^T\begin{bmatrix} x_i^0 \\ x_i^1 \end{bmatrix}, \quad v_i^q = (A_i^q)^T\left(\frac{1}{\gamma^2(z_i^q)}\begin{bmatrix} z_i^0 \\ -z_i^1 \end{bmatrix}\right). \tag{25}$$

The appearance of the outer-product terms in the equation above is potentially alarming. If the vectors u_i^q, v_i^q are dense, then even if A_i^q is sparse, the corresponding matrix M_i^q, and hence the Schur complement matrix M, will be dense. A direct factorization of the resulting dense matrix will be very expensive for even moderately high m.

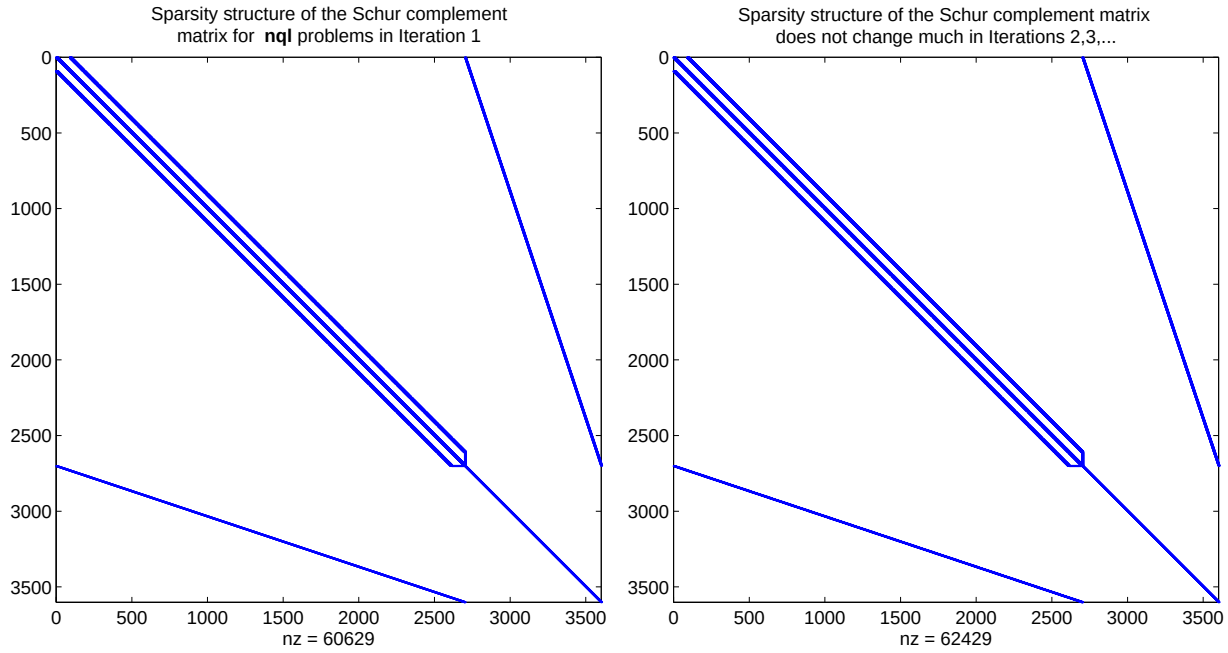

Figure 1: *The output of the* `spy` *function in* MATLAB *on the Schur complement matrix for* `nql30`. *Later iterations introduce less than 3% new nonzero elements.*

The observed behavior of the density of the Schur complement matrix on test problems depends largely on the particular problem structure. When the problem has many small quadratic blocks, it is often the case that each block appears in only a small fraction of the constraints. In this case, all A_i^q matrices are sparse and the vectors u_i^q and v_i^q turn out to be sparse vectors for each i. Consequently, the Schur complement matrix remains relatively sparse for these problems and it can be factorized directly and cheaply. In Figure 1, the density structures of the Schur complement matrices in the first and later iterations of our algorithm applied to the the problem `nql30` depict the situation and are typical for all `nql` and `qssp` problems. Since we initially choose multiples of unit vectors for our variables, all the nonzero elements of the Schur complement matrix in the first iteration come from the nonzero elements of the constraint matrices. Later iterations introduce fewer than 3% new nonzero elements.

The situation is drastically different for problems where one of the quadratic blocks, say the ith block, is large. For such problems the vectors u_i^q, v_i^q are typically dense, and therefore, M_i^q is likely be a dense matrix even if the data A_i^q is sparse. However, observe that M_i^q is a rank-two perturbation of a sparse matrix when A_i^q is sparse. In such a situation, it may be advantageous to use the Sherman-Morrison-Woodbury update formula [9] when solving the Schur complement equation (10). This is a standard strategy used in linear programming when there are dense columns in the constraint matrix and this is the approach we used in our implementation of SDPT3.

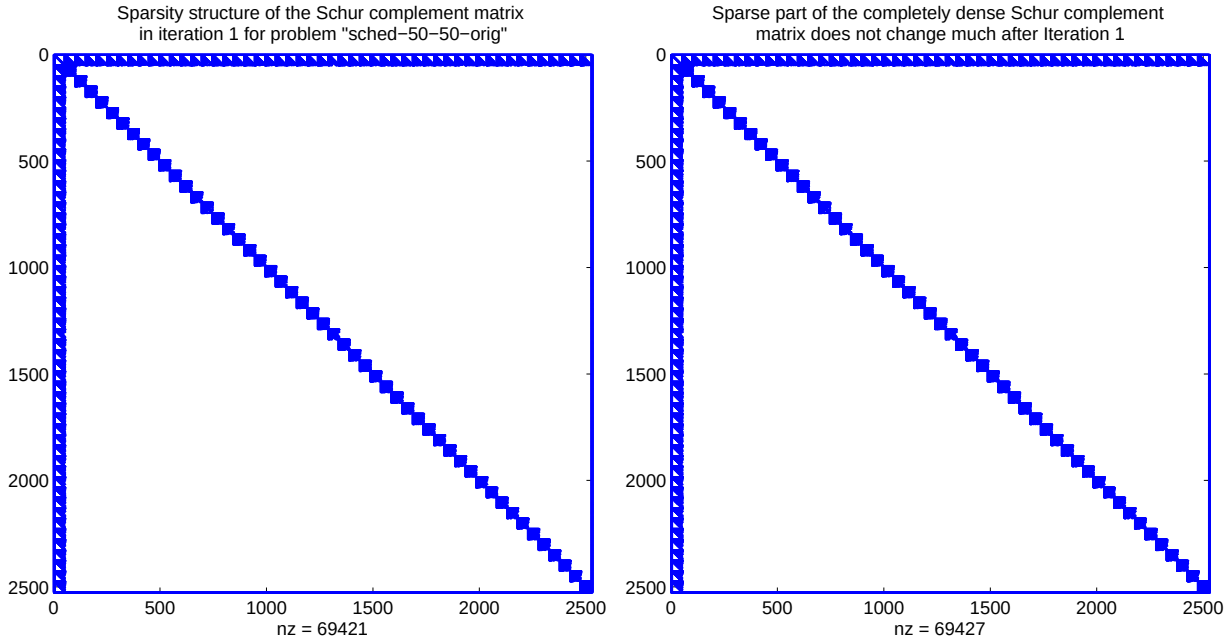

Figure 2: *The output of the* `spy` *function in* MATLAB *for problem* `sched-50-50-orig` *on (i) the complete Schur complement matrix in the first iteration, (ii) the sparse portion of the Schur complement matrix in the following iterations.*

This approach helps tremendously on the scheduling problems from the DIMACS Challenge set. Figure 2 depicts the Schur complement matrix M in the first iteration and its sparse portion in the following iterations. While these two matrices have almost identical sparsity patterns, the complete Schur complement matrix becomes completely dense after the first iteration.

To apply the Sherman-Morrison-Woodbury formula, we need to modify the sparse portion of the matrix M_i^q slightly. Since the diagonal matrix J_i has a negative component, the matrix $(A_i^q)^T J_i A_i^q$ need not be a positive definite matrix, and therefore the Cholesky factorization of the sparse portion of M_i^q need not exist. To overcome this problem, we use the following identity:

$$M_i^q \;=\; \frac{\langle x_i^q,\, z_i^q\rangle}{\gamma^2(z_i^q)}(A_i^q)^T A_i^q \;+\; u_i^q(v_i^q)^T + v_i^q(u_i^q)^T - 2\frac{\langle x_i^q,\, z_i^q\rangle}{\gamma^2(z_i^q)}k_i k_i^T, \tag{26}$$

where u_i^q and v_i^q are as in (25) and

$$k_i \;\;=\;\; (A_i^q)^T\, e_i^q. \tag{27}$$

Note that if A_i^q is a large sparse matrix with a few dense rows, we also use the Sherman-Morrison-Woodbury formula to handle the matrix $(A_i^q)^T A_i^q$ in (26).

We end our discussion on the computation of the HKM direction with the following

formula that is needed in the computation of the right-hand-side vector (12):

$$(\mathcal{E}_i^q)^{-1}(R_c^q)_i \;=\; \frac{\sigma\mu}{\gamma^2(z_i^q)} \left[\begin{array}{c} z_i^0 \\ -z_i^1 \end{array} \right] \;-\; x_i^q. \tag{28}$$

Just as for the HKM direction, we can obtain a very simple formula for $(\mathcal{E}_i^q)^{-1}\mathcal{F}_i^q$ for the NT direction. By noting that $G_i x_i^q = G_i^{-1} z_i^q$, it is easy to see that the ith block $(\mathcal{E}_i^q)^{-1}\mathcal{F}_i^q = G_i^{-2}$, and a rather straightforward algebraic manipulation gives the following identity:

$$(\mathcal{E}_i^q)^{-1}\mathcal{F}_i^q \;=\; G_i^{-2} \;=\; \frac{1}{\omega_i^2}\left(\left[\begin{array}{cc} -1 & 0 \\ 0 & I \end{array} \right] + 2 \left[\begin{array}{c} t_i^0 \\ -t_i^1 \end{array} \right] \left[\begin{array}{c} t_i^0 \\ -t_i^1 \end{array} \right]^T \right). \tag{29}$$

For the NT direction, the formula in (28) also holds and we have:

$$M_i^q \;=\; \frac{1}{\omega_i^2}\left((A_i^q)^T J_i A_i^q \;+\; 2u_i^q (u_i^q)^T \right), \text{ with } u_i^q = (A_i^q)^T \left[\begin{array}{c} t_i^0 \\ -t_i^1 \end{array} \right]. \tag{30}$$

We note that the identity (30) describing the NT direction was observed by other authors — see, e.g., [8]. The identities (23) and (24), however, appear to be new in the literature. It is straightforward, if a bit tedious, to verify these formulas. In addition to simplifying the search direction computation, these identities can be used to provide a simple proof of the scale-invariance of the HKM search direction in second-order cone programming. In [25], Tsuchiya proves this result and the scale-invariance of the NT direction using two-page arguments for each proof. We refer the reader to [20] for a description of scale-invariance and provide the following simple and instructive proof:

Proposition 1 *Consider a pure second-order cone programming problem ($n_s = 0$ and n_l=0). The HKM and NT directions for this problem are scale-invariant.*

Proof. The scaled problem is constructed as follows: Let $F_i \in \mathcal{G}_i$ denote a scaling matrix for block i where $\mathcal{G}_i$ is the automorphism group of the cone $K_q^{q_i}$. For future reference, note that we have

$$F_i^T \bar{J}_i F_i = \bar{J}_i, \quad F_i^T \bar{J}_i = \bar{J}_i F_i^{-1}, \quad \bar{J}_i F_i = F_i^{-T} \bar{J}_i, \quad \bar{J}_i F_i^T = F_i^{-1} \bar{J}_i, \tag{31}$$

where $\bar{J}_i = -J_i$, and J_i is as in (25). Let $F = \mathrm{diag}[F_1, \ldots, F_{n_q}]$ and define the scaled quantities as follows:

$$\hat{A}^q = F^{-1}A^q, \quad \hat{b} = b, \quad \hat{c}^q = F^{-1}c^q, \quad \hat{x}^q = F^T x^q, \quad \hat{y} = y, \quad \hat{z}^q = F^{-1}z^q.$$

Note that $\hat{r}_p = r_p$ and $\hat{R}_d = F^{-1}R_d$. First, we consider the HKM direction. We observe that $\gamma^2(\hat{z}_i^q) = (\hat{z}_i^q)^T \bar{J}_i \hat{z}_i^q = (z_i^q)^T F_i^{-T} \bar{J}_i F_i^{-1} z_i^q = \gamma^2(z_i^q)$. Now, from equation (24) and (31) it follows that each M_i^q, and therefore M^q, is invariant with respect to this automorphic scaling. Using (28), we see that h in (12) is also invariant. Now, if

we denote the HKM search direction for the scaled problem by $(\widehat{\Delta x}^q, \widehat{\Delta y}, \widehat{\Delta z}^q)$ and the corresponding direction for the unscaled problem by $(\Delta x^q, \Delta y, \Delta z^q)$ we immediately obtain $\widehat{\Delta y} = \Delta y$ from equation (10), $\widehat{\Delta z}^q = F^{-1}\Delta z^q$ from (13) and $\widehat{\Delta x}^q = F^T \Delta x^q$ from (15) and (28). Thus, the HKM direction for SOCP is scale-invariant.

To prove the result for the NT direction, we first observe that $\gamma(\hat{x}_i^q) = \gamma(x_i^q)$ and that ω_i defined in (20) remains unchanged after scaling. The scaled equivalent of ξ_i defined in (20) is $\hat{\xi}_i = \frac{1}{\omega_i}\hat{z}_i^q + \omega_i \bar{J}_i \hat{x}_i^q = F_i^{-1}(\frac{1}{\omega_i} z_i^q + \omega_i \bar{J}_i x_i^q) = F_i^{-1}\xi_i$. Thus, with the scaled quantities, we obtain

$$\hat{t}_i = \begin{bmatrix} \hat{t}_i^0 \\ \hat{t}_i^1 \end{bmatrix} = F_i^{-1} \begin{bmatrix} t_i^0 \\ t_i^1 \end{bmatrix}.$$

Now, from equation (30) and (31) it follows that each M_i^q, and therefore M^q, is invariant with respect to this automorphic scaling. Continuing as above, we conclude that the NT direction for SOCP must be scale-invariant as well. □

2.4 Step-length computation

Once a direction Δx is computed, a full step will not be allowed if $x + \Delta x$ violates the conic constraints. Thus, the next iterate must take the form $x + \alpha \Delta x$ for an appropriate choice of the step-length α. In this subsection, we discuss an efficient strategy to compute the step-length α.

For semidefinite blocks, it is straightforward to verify that, for the jth block, the maximum allowed step-length that can be taken without violating the positive semidefiniteness of the matrix $x_j^s + \alpha_j^s \Delta x_j^s$ is given as follows:

$$\alpha_j^s \;=\; \begin{cases} \dfrac{-1}{\lambda_{\min}((x_j^s)^{-1}\Delta x_j^s)}, & \text{if the minimum eigenvalue } \lambda_{\min} \text{ is negative} \\ \infty & \text{otherwise.} \end{cases} \tag{32}$$

If the computation of eigenvalues necessary in α_j^s above becomes expensive, then we resort to finding an approximation of α_j^s by estimating extreme eigenvalues using Lanczos iterations [22]. This approach is quite accurate in general and represents a good trade-off between the effort versus quality of the resulting stepsizes.

For quadratic blocks, the largest step-length α_i^q that keeps the next iterate feasible with respect to the kth quadratic cone can be computed as follows. Let

$$a_i = \gamma^2(\Delta x_i^q), \quad b_i = \langle \Delta x_i^q, -J_i x_i^q \rangle, \quad c_i = \gamma^2(x_i^q),$$

where J_i is the matrix defined in (25) and let

$$d_i \;=\; b_i^2 - a_i c_i.$$

We want the largest α with $a_i\alpha^2 + 2b_i\alpha + c_i > 0$ for all smaller positive values. This

is given by

$$\alpha_i^q \;=\; \begin{cases} \dfrac{-b_i - \sqrt{d_i}}{a_i} & \text{if } a_i < 0 \text{ or } b_i < 0, a_i \le b_i^2/c_i \\ \dfrac{-c_i}{2b_i} & \text{if } a_i = 0, b_i < 0 \\ \infty & \text{otherwise.} \end{cases}$$

For the linear block, the maximum allowed step-length α_i^l for the hth component is given by

$$\alpha_h^l \;=\; \begin{cases} \dfrac{-x_h^l}{\Delta x_h^l}, & \text{if } \Delta x_h^l < 0 \\ \infty & \text{otherwise.} \end{cases}$$

Finally, an appropriate step-length α that can be taken in order for $x+\alpha\Delta x$ to satisfy all the conic constraints takes the form

$$\alpha \;=\; \min\left(1,\; \gamma \min_{1\le j\le n_s} \alpha_j^s,\; \gamma \min_{1\le i\le n_q} \alpha_i^q,\; \gamma \min_{1\le h\le n_l} \alpha_h^l\right), \tag{33}$$

where γ (known as the step-length parameter) is typically chosen to be a number slightly less than 1, say 0.98, to ensure that the next iterate $x + \alpha\Delta x$ stays strictly in the interior of all the cones.

For the dual direction Δz, we let the analog of α_j^s, α_i^q and α_h^l be β_j^s, β_i^q and β_h^l, respectively. Similar to the primal direction, the step-length that can be taken by the dual direction Δz is given by

$$\beta \;=\; \min\left(1,\; \gamma \min_{1\le j\le n_s} \beta_j^s,\; \gamma \min_{1\le i\le n_q} \beta_i^q,\; \gamma \min_{1\le h\le n_l} \beta_h^l\right). \tag{34}$$

2.5 Sherman-Morrison-Woodbury formula and iterative refinement

In this subsection, we discuss how we solve the Schur complement equation when M is a low rank perturbation of a sparse matrix. As discussed in Section 2.3 such situations arise when the SQLP does not have a semidefinite block, but has large quadratic blocks or the constraint matrices A_i^q, A^l have a small number of dense rows. In such a case, the Schur complement matrix M can be written in the form

$$M \;=\; H + UU^T \tag{35}$$

where H is a sparse symmetric matrix and U has only few columns. If H is non-singular, then by the Sherman-Morrison-Woodbury formula, the solution of the Schur complement equation is given by

$$\Delta y \;=\; \hat{h} - H^{-1}U\left(I + U^T H^{-1} U\right)^{-1} U^T \hat{h}, \tag{36}$$

where $\hat{h} = H^{-1}h$.

Computing Δy via the Sherman-Morrison-Woodbury update formula above is not always stable, and the computed solution for Δy can be highly inaccurate when H is ill-conditioned. To overcome such a difficulty, we combine the Sherman-Morrison-Woodbury update with iterative refinement [11]. It is noted in [11] that iterative refinement is beneficial even if the residuals are computed only at the working precision. Our numerical experience with the SQLP problems from the DIMACS Challenge set confirmed that iterative refinement very often does greatly improve the accuracy of the computed solution for Δy via the Sherman-Morrison-Woodbury formula. However, we must mention that iterative refinement can occasionally fail to provide any significant improvement. We have not yet incorporated a stable and efficient method for computing Δy when M has the form (35), but note that Goldfarb, Scheinberg, and Schmieta [8] discuss a stable product-form Cholesky factorization approach to this problem.

3 Initial iterates

Our algorithms can start with an infeasible starting point. However, the performance of these algorithms is quite sensitive to the choice of the initial iterate. As observed in [7], it is desirable to choose an initial iterate that at least has the same order of magnitude as an optimal solution of the SQLP. If a feasible starting point is not known, we recommend that the following initial iterate be used:

$$
\begin{aligned}
y^0 &= 0, \\
(x_j^s)^0 = \xi_j^s\, I_{s_j}, \quad (z_j^s)^0 &= \eta_j^s\, I_{s_j}, \quad j = 1, \ldots, n_s, \\
(x_i^q)^0 = \xi_i^q\, e_i^q, \quad (z_i^q)^0 &= \eta_i^q\, e_i^q, \quad i = 1, \ldots, n_q, \\
(x^l)^0 = \xi^l\, e^l, \quad (z^l)^0 &= \eta^l\, e^l,
\end{aligned}
$$

where I_{s_j} is the identity matrix of order s_j, and

$$
\xi_j^s = s_j \max\nolimits_{1\le k\le m} \frac{1+|b_k|}{1+\|A_j^s(:,k)\|}, \quad \eta_j^s = \frac{1}{\sqrt{s_j}}\left[1+\max(\max_k\{\|A_j^s(:,k)\|\},\, \|c_j^s\|_F)\right],
$$

$$
\xi_i^q = \sqrt{q_i} \max_{1\le k\le m} \frac{1+|b_k|}{1+\|A_i^q(:,k)\|}, \quad \eta_i^q = \sqrt{q_i}\,[1+\max(\max_k\{\|A_i^q(:,k)\|\},\, \|c_i^q\|)],
$$

$$
\xi^l = \max_{1\le k\le m} \frac{1+|b_k|}{1+\|A^l(:,k)\|}, \quad \eta^l = 1+\max(\max_k\{\|A^l(:,k)\|\},\, \|c^l\|),
$$

where $A_j^s(:,k)$ denotes the kth column of A_j^s, and $A_i^q(:,k)$ and $A_i^l(:,k)$ are defined similarly.

By multiplying the identity matrix I_{s_i} by the factors ξ_i^s and η_i^s for the semidefinite blocks, and similarly for the quadratic and linear blocks, the initial iterate has a better chance of having the appropriate order of magnitude.

The initial iterate above is set by calling `infeaspt.m`, with initial line

```
function  [X0,y0,Z0] = infeaspt(blk,A,C,b,options,scalefac),
```

where `options = 1 (default)` corresponds to the initial iterate just described, and `options = 2` corresponds to the choice where the blocks of `X0`, `Z0` are `scalefac` times identity matrices or unit vectors, and `y0` is a zero vector.

4 The main routine

The main routine that corresponds to the infeasible path-following algorithms described in Section 2 is `sqlp.m`:

```
[obj,X,y,Z,gaphist,infeashist,info,Xiter,yiter,Ziter] =
sqlp(blk,A,C,b,X0,y0,Z0,OPTIONS).
```

Input arguments.

`blk`: a cell array describing the block structure of SQLP problem (see below).

`A`, `C`, `b`: SQLP data (see below).

`X0`, `y0`, `Z0`: an initial iterate.

`OPTIONS`: a structure array of parameters (see below).

If the input argument `OPTIONS` is omitted, default values are used.

Output arguments. The names chosen for the output arguments explain their contents. The argument `info` is a 5-dimensional vector containing performance information; see [23] for details. The argument `(Xiter,yiter,Ziter)` is new in this release: it is the last iterate of `sqlp.m`, and if desired, the user can continue the iteration process with this as the initial iterate. Such an option allows the user to iterate for a certain amount of time, stop to analyze the current solution, and continue if necessary.

Note that, while `(X,y,Z)` normally gives approximately optimal solutions, if `info(1)` is 1 the problem is suspected to be primal infeasible and `(y,Z)` is an approximate certificate of infeasibility, with $\mathtt{b}^T\mathtt{y} = 1$, `Z` in the appropriate cone, and $\mathtt{A}^T\mathtt{y} + \mathtt{Z}$ small, while if `info(1)` is 2 the problem is suspected to be dual infeasible and `X` is an approximate certificate of infeasibility, with $\langle \mathtt{C}, \mathtt{X} \rangle = -1$, `X` in the appropriate cone, and `A X` small.

A structure array for parameters.

`sqlp.m` use a number of parameters which are specified in a MATLAB structure array called `OPTIONS` in the m-file `parameters.m`. If desired, the user can change the values of these parameters. The meaning of most of the specified fields in `OPTIONS` are given in [23]. The new fields are self-explanatory.

C Mex files used.

Our software uses a number of Mex routines generated from C programs written to carry out certain operations that MATLAB is not efficient at. In particular, operations such as extracting selected elements of a matrix, and performing arithmetic operations on these selected elements are all done in C. As an example, the vectorization operation **svec** is coded in the C program `mexsvec.c`.

Our software also uses a number of Mex routines generated from Fortran programs written by Esmond Ng and Barry Peyton and Joseph Liu for computing sparse Cholesky factorization. These programs are adapted from the LIPSOL software written by Yin Zhang [28].

Cell array representation for problem data.

Our implementation SDPT3 exploits the block structure of the given SQLP problem. In the internal representation of the problem data, we classify each semidefinite block into one of the following two types:

1. a dense or sparse matrix of dimension greater than or equal to 30;
2. a sparse block-diagonal matrix consisting of numerous sub-blocks each of dimension less than 30.

The reason for using the sparse matrix representation to handle the case when we have numerous small diagonal blocks is that it is less efficient for MATLAB to work with a large number of cell array elements compared to working with a single cell array element consisting of a large sparse block-diagonal matrix. Technically, no problem will arise if one chooses to store the small blocks individually instead of grouping them together as a sparse block-diagonal matrix.

For the quadratic part, we typically group all quadratic blocks (small or large) into a single block, though it is not mandatory to do so. If there are a large number of small blocks, it is advisable to group them all together as a single large block consisting of numerous small sub-blocks for the same reason we mentioned before.

Let $L = n_s + n_q + 1$. For each SQLP problem, the block structure of the problem data is described by an $L \times 2$ cell array named `blk`, The content of each of the elements of the cell arrays is given as follows. If the jth block is a semidefinite block consisting of a single block of size $\mathtt{s_j}$, then

$\mathtt{blk\{j,1\} = 's'}$ $\quad \mathtt{blk}\{j,2\} = [\mathtt{s_j}]$

$\mathtt{A\{j\}} = [\bar{\mathtt{s}}_\mathtt{j}\ \mathtt{x\,m\ sparse}]$

$\mathtt{C\{j\},\ X\{j\},\ Z\{j\}} = [\mathtt{s_j\,x\,s_j\ \ double\ or\ sparse}]$,

where $\bar{\mathtt{s}}_\mathtt{j} = \mathtt{s_j}(\mathtt{s_j}+1)/2$.

If the jth block is a semidefinite block consisting of numerous small sub-blocks, say p of them, of dimensions $\mathtt{s_{j1}}, \mathtt{s_{j2}}, \ldots, \mathtt{s_{jp}}$ such that $\sum_{k=1}^{p} \mathtt{s_{jk}} = \mathtt{s_j}$, then

$\mathtt{blk\{j,1\} = 's'}$ $\quad \mathtt{blk}\{j,2\} = [\mathtt{s_{j1}}\ \mathtt{s_{j2}}\ \cdots\ \mathtt{s_{jp}}]$

$\texttt{A\{j\}} = [\bar{\texttt{s}}_\texttt{j}\ \texttt{x m sparse}]$

$\texttt{C\{j\}, X\{j\}, Z\{j\}} = [\texttt{s}_\texttt{j}\ \texttt{x}\ \texttt{s}_\texttt{j}\ \texttt{sparse}]$,

where $\bar{\texttt{s}}_\texttt{j} = \sum_{\texttt{k}=1}^{\texttt{p}} \texttt{s}_{\texttt{jk}}(\texttt{s}_{\texttt{jk}}+1)/2$.

The above storage scheme for the data matrix A_j^s associated with the semidefinite blocks of the SQLP problem represents a departure from earlier versions of our implementation, such as the one described in [23]. Previously, the semidefinite part of A was represented by an $n_s \times m$ cell array, where A{j,k} corresponds to the kth constraint matrix associated with the jth semidefinite block, and it was stored as an individual matrix in either dense or sparse format. Now, we store all the constraint matrices associated with the jth semidefinite block in vectorized form as a single $\bar{s}_j \times m$ matrix where the kth column of this matrix corresponds to the kth constraint matrix. The data format we used in earlier versions of SDPT3 was more natural but, for the sake of computational efficiency, we adopted our current data representation. The reason for such a change is again due to the fact that it is less efficient for MATLAB to work with a single cell array with many cells. We also avoid explicit loops over the index k. In the next section, we will discuss the consequence of this modification in our storage scheme.

If the ith block is a quadratic block consisting of numerous sub-blocks, say p of them, of dimensions $\texttt{q}_{\texttt{i1}}, \texttt{q}_{\texttt{i2}}, \ldots, \texttt{q}_{\texttt{ip}}$ such that $\sum_{k=1}^{p} \texttt{q}_{\texttt{ik}} = \texttt{q}_\texttt{i}$, then

$\texttt{blk\{i,1\} = 'q'}$ $\qquad \texttt{blk}\{\texttt{i},2\} = [\texttt{q}_{\texttt{i1}}\ \texttt{q}_{\texttt{i2}}\ \cdots\ \texttt{q}_{\texttt{ip}}]$

$\texttt{A\{i\}} = [\texttt{q}_\texttt{i}\ \texttt{x m sparse}]$

$\texttt{C\{i\}, X\{i\}, Z\{i\}} = [\texttt{q}_\texttt{i}\ \texttt{x 1 double or sparse}]$.

If the ith block is the linear block, then

$\texttt{blk\{i,1\} = 'l'}$ $\qquad \texttt{blk}\{\texttt{i},2\} = \texttt{n}_\texttt{l}$

$\texttt{A\{i\}} = [\texttt{n}_\texttt{l}\ \texttt{x m sparse}]$

$\texttt{C\{i\}, X\{i\}, Z\{i\}} = [\texttt{n}_\texttt{l}\ \texttt{x 1 double or sparse}]$.

Caveats.

The user should be aware that SQLP is more complicated than linear programming. For example, it is possible that both primal and dual problems are feasible, but their optimal values are not equal. Also, either problem may be infeasible without there being a certificate of that fact (so-called weak infeasibility). In such cases, our software package is likely to terminate after some iterations with an indication of short step-length or lack of progress. Also, even if there is a certificate of infeasibility, our infeasible-interior-point methods may not find it. In our very limited testing on strongly infeasible problems, most of our algorithms have been quite successful in detecting infeasibility.

5 Computational experiments

Here we describe the results of our computational testing of SDPT3, on problems from the SDPLIB collection of Borchers [4] as well as the DIMACS Challenge test problems [16]. In both, we solve a selection of the problems; in the DIMACS problems, these are selected as the more tractable problems, while our subset of the SDPLIB problems is more representative (but we cannot solve the largest two `maxG` problems). Since our algorithm is a primal-dual method storing the primal iterate `X`, it cannot exploit common sparsity in `C` and the constraint matrix as well as dual methods or nonlinear-programming based methods. We are therefore unable to solve the largest problems.

Most of the results given were obtained on a Sun UltraSPARC 170E (170 MHz) with 512MB of memory running Solaris 2.6, using MATLAB 5.3 with the numerics library based on LAPACK distributed by MathWorks. (We had some difficulty with MATLAB 6 using some of our codes, so use 5.3 throughout.) For some of the larger problems we require more memory, and so used (one processor of) a Sun UltraSPARC 420R (450 MHz) with 2 GB of memory.

The test problems are listed in Tables 1 and 2, along with their dimensions.

5.1 Cholesky factorization

Earlier versions of SDPT3 were intended for problems that always have semidefinite cone constraints. As we indicated above, for such problems, the Schur complement matrix M in (11) is a dense matrix after the first iteration. To solve the associated linear system (10), we first find a Cholesky factorization of M and then solve two triangular systems. When M is dense, a reordering of the rows and columns of M does not alter the efficiency of the Cholesky factorization and specialized sparse Cholesky factorization routines are not useful. Therefore, earlier versions of SDPT3 (up to version 1.3) simply used MATLAB's `chol` routine for Cholesky factorizations. For versions 2.1 and 2.2, we introduced our own Cholesky factorization routine `mexchol` that utilizes loop unrolling and provided 2-fold speed-ups on some architectures compared to MATLAB's `chol` routine. However, in newer versions of MATLAB that use numerics libraries based on LAPACK, MATLAB's `chol` routine is more efficient than our Cholesky factorization routine `mexchol` for dense matrices. Thus, in version 3.0, we use MATLAB's `chol` routine whenever M is dense.

For the solution of most second-order cone programming problems in DIMACS test set, however, MATLAB's `chol` routine is not competitive. This is largely due to the fact that the Schur complement matrix M is often sparse for SOCPs and LPs, and MATLAB cannot sufficiently take advantage of this sparsity. To solve such problems more efficiently we imported the sparse Cholesky solver in Yin Zhang's LIPSOL [28], an interior-point code for linear programming problems. It should be noted that LIPSOL uses Fortran programs developed by Esmond Ng, Barry Peyton, and Joseph Liu for Cholesky factorization. When SDPT3 uses LIPSOL's Cholesky solver, it first generates a symbolic factorization of the Schur complement matrix to

Problem	m	semidefinite blocks	second-order blocks	linear block
bm1	883	882	–	–
copo14	1275	[14 x 14]	–	364
copo23	5820	[23 x 23]	–	1771
filter48-socp	969	48	49	931
filtinf1	983	49	49	945
hamming-7-5-6	1793	128	–	–
hamming-9-8	2305	512	–	–
hinf12	43	24	–	–
hinf13	57	30	–	–
minphase	48	48	–	–
nb	123	–	[793 x 3]	4
nb-L1	915	–	[793 x 3]	797
nb-L2	123	–	[1637, 838 x 3]	4
nb-L2-bessel	123	–	[123, 838 x 3]	4
nql30	3601	–	[900 x 3]	5560
nql60	14401	–	[3600 x 3]	21920
nql180	129601	–	[32400 x 3]	195360
qssp30	5674	–	[1891 x 4]	3600
sched-50-50-orig	2527	–	[2474, 3]	2502
sched-50-50-scaled	2526	–	2475	2502
sched-100-50-orig	4844	–	[4741, 3]	5002
sched-100-50-scaled	4843	–	4742	5002
sched-100-100-orig	8338	–	[8235, 3]	10002
sched-100-100-scaled	8337	–	8236	10002
sched-200-100-orig	18087	–	[17884, 3]	20002
sched-200-100-scaled	18086	–	17885	20002
torusg3-8	512	512	–	–
toruspm3-8-50	512	512	–	–
truss5	208	[33 x 10, 1]	–	–
truss8	496	[33 x 19, 1]	–	–

Table 1: Selected DIMACS Challenge Problems. Notation like [33 x 19] indicates that there were 33 semidefinite blocks, each a symmetric matrix of order 19, etc.

Problem	m	semidefinite blocks	linear block
arch8	174	161	174
control7	666	[70, 35]	–
control10	1326	[100, 50]	–
control11	1596	[110, 55]	–
gpp250-4	251	250	–
gpp500-4	501	500	–
hinf15	91	37	–
mcp250-1	250	250	–
mcp500-1	500	500	–
qap9	748	82	–
qap10	1021	101	–
ss30	132	294	132
theta3	1106	150	–
theta4	1949	200	–
theta5	3028	250	–
theta6	4375	300	–
truss7	86	[150 x 2, 1]	–
truss8	496	[33 x 19, 1]	–
equalG11	801	801	–
equalG51	1001	1001	–
equalG32	2001	2001	–
maxG11	800	800	–
maxG51	1000	1000	–
maxG32	2000	2000	–
qpG11	800	1600	–
qpG112	800	800	800
qpG51	1000	2000	–
qpG512	1000	1000	1000
thetaG11	2401	801	–
thetaG11n	1600	800	–
thetaG51	6910	1001	–
thetaG51n	5910	1000	–

Table 2: Selected SDPLIB Problems. Note that `qpG112` is identical to `qpG11` except that the structure of the semidefinite block is exposed as a sparse symmetric matrix of order 800 and a diagonal block of the same order, which can be viewed as a linear block, and similarly for `qpG512`. Also, `thetaG11n` is a more compact formulation of `thetaG11`, and similarly for `thetaG51n`.

determine the pivot order by examining the sparsity structure of this matrix carefully. Then, this pivot order is re-used in later iterations to compute the Cholesky factors. Contrary to the case of linear programming, however, the sparsity structure of the Schur complement matrix can change during the iterations for SOCP problems. If this happens, the pivot order has to be recomputed. We detect changes in the sparsity structure by monitoring the nonzero elements of the Schur complement matrix. Since the default initial iterates we use for an SOCP problem are unit vectors but subsequent iterates are not, there is always a change in the sparsity pattern of M after the first iteration. After the second iteration, the sparsity pattern remains unchanged for most problems, and only one more change occurs in a small fraction of the test problems.

The effect of including a sparse Cholesky solver option for SOCP problems was dramatic. We observed speed-ups up to two orders of magnitude. Version 3.0 of SDPT3 automatically makes a choice between MATLAB's built-in `chol` routine and the sparse Cholesky solver based on the density of the Schur complement matrix. The cutoff density is provided in the `spdensity` field of the `OPTIONS` structure array.

5.2 Vectorized matrices vs. sparse matrices

The current release, version 3.0, of the code stores the constraint matrix in "vectorized" form as described in Sections 2 and 4. In the previous version 2.2, `A` was a doubly subscripted cell array of symmetric matrices for the semidefinite blocks, as we outlined at the end of the previous section. The result of the change is that much less storage is required for the constraint matrix, and that we save a considerable amount of time in forming the Schur complement matrix M in (11) by avoiding loops over the index k. Operations relating to forming and factorizing the Schur complement and hence computing the predictor search direction comprise much of the computational work for most problem classes, ranging from 25% for `qpG11` up to 99% for the larger `theta` problems, the `control` problems, `copo14`, `hamming-7-5-6`, and the `nb` problems. Other significant parts are computing the corrector search direction (between less than 1% and 75%) and computing step lengths (between less than 1% and 60%).

While we now store the constraint matrix in vectorized form, the parts of the iterates `X` and `Z` corresponding to semidefinite blocks are still stored as matrices, since that is how the user wants to access them.

Results are given in Tables 3 through 6: Tables 3 and 4 give results on the DIMACS problems for both SDPT3-3.0 and SDPT3-2.2, while Tables 5 and 6 give the comparable results for the SDPLIB problems. In all of these, the format is the same. We give the number of iterations required, the time in seconds, and four measures of the precision of the computed answer. The first column gives the logarithm (to base 10) of the total complementary slackness; the second the scaled primal infeasibility $\|Ax - b\|/(1 + \max \|b_k\|)$, and the third $\|A^T y + z - c\|/(1 + \max |c|)$, where the norm is subordinate to the inner product and the maximum taken over all components of c; and the last the maximum of 0 and $\langle c, x\rangle - b^T y$.

In general, the codes solved the problems to reasonable accuracy. On the DIMACS problems, only 1 digit of precision was obtained for `bm1` (NT only), `filtinf1`, `nql180`,

and `torusg3-8`, and even less for some of the `sched-orig` problems. However, the optimal value for the latter problems and for `torusg3-8` is at least of order 10^5, so the relative accuracy is better than it appears. The other measures of accuracy were also reasonable, except for the NT search direction applied to `nql60` and again three of the `sched-orig` problems. For the SDPLIB problems, again the accuracy is reasonable; on the `equalG` problems, the optimal value is of order 10^3, so the duality gap is not too excessive. We note that on three problems, `qpG11`, `qpG51`, and `sched-100-100-orig`, the algorithm terminated with an indication that `X` and `Z` were not both positive definite. However, this is a conservative test designed to stop if numerical difficulties are imminent. Using SeDuMi's `eigK.m` routine to check the iterates, it was found that in all cases both were feasible in the conic constraints.

To compare the two codes in terms of time, we consider only the problems that both codes could solve, and omit the simplest problems with times under 20 seconds (the `hinf` problems, `minphase`, and `truss5` and `truss7`). For the remaining problems, we compute the ratio of the times taken by the two codes, take its logarithm to base 2, and then plot the results in decreasing order of absolute values. The results are shown in Figures 1 and 2 for the HKM and NT search directions. A bar of height 1 indicates that SDPT3-3.0 was 2 times faster than SDPT3-2.2, of -1 the reverse, and of 3 that SDPT3-3.0 was 8 times faster. Note that the new version using vectorized matrices is uniformly faster for the HKM direction, and faster or comparable for the NT direction.

5.3 HKM vs. NT

The new version of the code allows only two search directions, HKM and NT. Version 2.2 also allowed the AHO direction of Alizadeh, Haeberly, and Overton [2] and the GT (Gu-Toh) direction — see [21], but these are uncompetitive when the problems are of large scale. We intend to keep Version 2.2 of the code available for those who wish to experiment with these other search directions, which tend to give more accurate results on smaller problems.

To compare the two remaining search directions, we again use a bar chart to show their relative performance as in Figures 1 and 2. It is clear that the HKM direction is almost universally faster than NT, but the exceptions do show some patterns. The NT direction was faster on the `nb` and `nql` problems, which all have a large number of low-dimensional second-order cone constraints. (The reason for this behavior is not hard to understand. By comparing the formula in (23) for the HKM direction with (29) for the NT direction, it is clear that more computation is required to assemble the Schur complement matrix and more low-rank updating is necessary for the former direction, and these computations can dominate the work when the dimensions of each second-order cone is small.) Also, the NT direction is faster on the semidefinite problem `copo23`, again with many small blocks, and on the `sched` problems, which have small as well as large second-order blocks. The HKM direction is *much* faster on `maxG32` and considerably faster on the other `maxG` and the `qpG` problems. Because there is a class of problems on which the NT direction is faster, we feel it is worthwhile

to keep both options.

5.4 Homogeneous vs infeasible interior-point methods

Version 2.2 also allowed the user to employ homogeneous self-dual algorithms instead of the usual infeasible interior-point methods. However, this option almost always took longer than the default choice, and so it has been omitted from the current release. One theoretical advantage of the homogeneous self-dual approach is that it is oriented towards either producing optimal primal and dual solutions or generating a certificate of primal or dual infeasibility, while the infeasible methods strive for optimal solutions only, but detect infeasibility if either the dual or primal iterates diverge. However, we have observed no advantage to the homogeneous methods when applied to infeasible problems. We should mention, however, that our current version does not detect infeasibility in the problem `filtinf1`, but instead stops with a primal near-feasible solution and a dual feasible solution when it encounters numerical problems.

5.5 Presentation of problems

We note that `qpG11` and `qpG112`, and similarly `qpG51` and `qpG512`, are basically the same problem, but in the second version the linear variables are explicitly identified, rather than being part of a large sparse semidefinite block. The improvement in running time is dramatic: a factor of at least four (recall that `qpG512` is solved using the slower machine, `qpG51` by the faster). It is thus crucial to present problems to the algorithms correctly. We could add our own preprocessor to detect this structure, but believe users are aware of linear variables present in their problems. Unfortunately the versions of `qpG11` and `qpG51` in SDPLIB do not show this structure explicitly.

We also remark that the computation of the Lovasz theta function for a graph can be expressed as a semidefinite programming problem in two ways, and one of these is much more compact than the other, requiring a linear constraint only for each edge of the graph rather than also for each node, and so the problems `thetaG11n` and `thetaG51n` are much easier to solve than `thetaG11` and `thetaG51`, here by a factor up to two.

Finally, version 2.2 of SDPT3 included specialized routines to compute the Schur complement matrices directly for certain classes of problem (e.g., maxcut problems). In earlier versions of SDPT3, these specialized routines had produced dramatic decreases in solution times, but for version 2.2, these gains were marginal, since our general sparse matrix routines provided almost as much speedup. We have therefore dropped these routines in version 3.0.

	HKM						NT					
Problem	Itn	log $\langle x, z\rangle$	err1	err3	err5	time	Itn	log $\langle x, z\rangle$	err1	err3	err5	time
bm1	18	-5	7-8	3-13	7-6	3314	17	-1	6-6	2-12	9-2	9640
copo14	15	-9	1-10	5-15	5-9	913	15	-10	8-10	5-15	2-8	901
copo23*	18	-9	8-10	8-15	0	32928	17	-9	3-10	1-14	2-8	31427
hamming-7-5-6	10	-8	2-15	0	2-8	965	10	-8	2-15	0	2-8	968
hamming-9-8	12	-7	4-15	0	9-8	2843	12	-7	4-15	0	9-8	3575
hinf12	42	-7	4-8	2-10	0	17	39	-7	5-8	2-10	0	22
hinf13	23	-2	2-4	5-13	0	13	22	-3	3-4	5-13	0	16
minphase	31	-7	2-8	0	0	19	34	-4	5-8	0	0	26
torusg3-8	15	-2	4-10	8-16	1-2	362	14	-1	4-9	7-16	2-1	1249
toruspm3-8-50	13	-6	5-10	6-16	3-6	316	14	-7	1-9	6-16	2-7	1230
truss5	19	-5	2-6	8-15	0	54	18	-4	2-6	9-15	9-5	57
truss8	21	-3	1-5	9-15	3-5	444	21	-4	1-5	1-14	0	481

Table 3: Computational results on SDP problems in the DIMACS Challenge test set using SDPT3-2.2. These were performed on a 170MHz Sun Ultra 170 with 512MB of memory, except for the problems marked with an asterisk which were on (1 processor of) a 420MHz Sun Ultra 420R with 2GB of memory.

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	HKM						NT					
Problem	Itn	log $\langle x, z \rangle$	err1	err3	err5	time	Itn	log $\langle x, z \rangle$	err1	err3	err5	time
bm1	19	-6	1-6	4-13	1-6	3027	15	-1	1-5	3-13	2-1	7901
copo14	15	-10	8-11	5-15	3-9	80	14	-9	5-11	5-15	1-8	77
copo23	18	-9	5-10	9-15	0	3644	17	-10	4-10	1-14	2-8	3463
filter48-socp	43	-6	1-6	1-13	4-5	136	50	-4	1-6	9-14	8-4	148
filtinf1	26	-1	3-5	1-12	1-1	89	26	-1	2-5	4-12	1-1	88
hamming-7-5-6	9	-7	6-15	0	3-7	192	9	-7	3-15	0	3-7	195
hamming-9-8	11	-6	7-15	8-14	2-6	725	11	-6	7-15	8-14	2-6	1295
hinf12	43	-8	2-8	3-10	0	14	39	-8	2-8	2-10	0	10
hinf13	26	-5	9-5	8-13	0	11	24	-4	1-4	7-13	0	7
minphase	35	-7	8-9	1-12	0	18	36	-4	2-8	6-13	0	19
nb	14	-5	8-6	2-9	5-5	122	14	-5	8-6	2-9	5-5	94
nb-L1	17	-5	7-5	6-11	1-5	212	17	-5	8-5	6-11	2-5	168
nb-L2	14	-8	2-8	3-11	2-8	202	12	-8	3-9	2-9	1-8	146
nb-L2-b	13	-8	5-9	2-11	9-9	120	12	-9	3-8	2-10	0	86
nql30	12	-3	7-6	3-7	2-3	28	12	-3	6-6	3-7	2-3	26
nql60	14	-4	7-5	2-9	0	211	14	-4	3-1	2-8	0	174
nql180*	9	-1	4-4	5-5	2-1	2117	11	-1	7-4	5-6	0	1687
qssp30	13	-1	3-5	2-6	3-1	141	14	-2	5-5	3-6	3-2	142
sched-50-50-orig	34	-2	6-4	1-10	0	59	31	-2	4-4	3-9	0	50
sched-50-50-scaled	22	-4	1-4	1-10	6-5	40	21	-4	3-5	4-10	5-5	35
sched-100-50-orig	50	+1	3-3	2-10	0	185	50	+5	2-3	7-2	4+5	171
sched-100-50-scaled	25	-2	4-4	6-11	2-2	95	24	-2	4-4	4-11	1-2	86
sched-100-100-orig	44	-2	2-1	1-10	0	298	42	+2	4+0	1-9	1+8	268
sched-100-100-scaled	22	-1	3-2	2-14	0	163	20	-1	3-2	3-14	0	140
sched-200-100-orig	48	-1	5-3	2-9	0	912	50	+8	5+3	5+3	3+7	883
sched-200-100-scaled	28	-1	3-3	5-9	0	556	27	-2	3-3	1-9	0	500
torusg3-8	15	-1	2-11	7-16	2-1	276	14	-1	2-10	8-16	4-1	1082
toruspm3-8-50	14	-6	2-11	6-16	2-6	263	13	-6	3-11	6-16	5-6	1001
truss5	16	-5	5-7	8-15	0	17	15	-5	5-7	1-14	2-5	20
truss8	16	-5	2-6	8-15	0	96	15	-4	3-6	9-15	0	115

Table 4: Computational results on DIMACS Challenge problems using SDPT3-3.0. These were performed on a 170MHz Sun Ultra 170 with 512MB of memory, except for the problems marked with an asterisk which were on (1 processor of) a 420MHz Sun Ultra 420R with 2GB of memory.

	HKM						NT					
Problem	Itn	log $\langle x, z\rangle$	err1	err3	err5	time	Itn	log $\langle x, z\rangle$	err1	err3	err5	time
arch8	19	-7	4-7	7-13	8-8	144	23	-6	3-8	4-13	7-7	199
control7	22	-4	1-6	2-9	6-5	771	22	-5	1-6	2-9	0	835
control10	24	-4	2-6	6-9	0	4045	24	-4	3-6	6-9	0	4592
control11	23	-4	2-6	6-9	0	5982	24	-4	3-6	6-9	0	7181
gpp250-4	16	-6	5-8	8-14	0	91	16	-4	2-7	8-14	0	205
gpp500-4	15	-5	6-8	1-14	3-5	575	15	-4	3-8	2-14	3-4	1549
hinf15	22	-2	3-4	1-12	0	22	22	-2	3-4	9-13	0	27
mcp250-1	13	-6	3-10	5-16	2-6	42	15	-6	9-11	4-16	8-7	138
mcp500-1	14	-6	3-10	5-16	1-6	257	16	-6	5-10	5-16	1-6	1182
qap9	15	-5	4-8	3-13	0	141	15	-5	5-8	3-13	0	147
qap10	14	-5	4-8	9-13	0	297	14	-5	4-8	9-13	0	302
ss30	18	-5	2-7	2-13	5-6	348	25	-6	3-8	2-13	4-6	720
theta3	14	-7	2-11	6-15	1-7	360	14	-8	1-10	6-15	3-8	381
theta4	15	-7	1-10	9-15	2-7	1866	15	-8	2-10	8-15	3-8	1911
theta5	15	-7	1-10	1-14	3-7	6519	14	-6	2-10	1-14	5-7	6151
theta6	15	-7	2-10	1-14	2-8	20280	15	-7	2-10	1-14	6-8	20392
truss7	23	-3	4-6	2-13	0	17	22	-5	7-6	2-13	0	23
truss8	21	-3	1-5	1-14	3-5	444	21	-4	1-5	9-15	0	477
equalG11	18	-7	3-10	3-16	5-7	2625	17	-4	2-9	2-16	7-5	7001
equalG51	20	-6	7-9	5-16	0	5748	16	-2	4-8	5-16	3-2	12127
equalG32*	19	-5	1-9	1-16	1-5	19425	15	-1	6-8	8-17	7-2	50465
maxG11	14	-6	4-10	7-16	3-6	943	14	-6	8-10	7-16	1-6	4087
maxG51	16	-5	3-10	4-16	1-5	3155	16	-6	2-9	4-16	2-6	11257
maxG32*	15	-5	2-9	1-15	7-6	6165	15	-6	4-9	1-15	2-6	35362
qpG11	14	-5	2-10	0	2-5	4839	16	-5	2-11	0	2-5	20717
qpG112	15	-6	4-10	0	2-6	1142	15	-6	1-9	0	2-6	4543
qpG51*	21	-4	4-10	0	1-4	7790	24	-3	6-9	0	1-3	32103
qpG512	17	-4	4-9	0	1-4	3462	25	-4	2-9	0	5-5	18182
thetaG11	19	-6	1-7	8-14	0	6173	17	-6	9-8	2-14	0	8902
thetaG11n	15	-6	4-12	0	1-6	3113	15	-6	4-12	0	1-6	6742
thetaG51*	33	-6	4-8	1-14	4-6	109681	30	-5	5-8	8-14	9-6	102457
thetaG51n*	20	-7	2-9	2-14	0	37692	22	-7	2e-9	3e-14	0	46453

Table 5: Computational results on SDPLIB problems using SDPT3-2.2. These were performed on a 170MHz Sun Ultra 170 with 512MB of memory, except for the problems marked with an asterisk which were on (1 processor of) a 420MHz Sun Ultra 420R with 2GB of memory.

	HKM						NT					
Problem	Itn	log $\langle x, z \rangle$	err1	err3	err5	time	Itn	log $\langle x, z \rangle$	err1	err3	err5	time
arch8	18	-6	1-7	5-13	8-7	96	23	-6	6-8	4-13	9-7	151
control7	23	-5	7-7	2-9	4-6	380	23	-4	4-7	2-9	6-5	379
control10	25	-5	1-6	6-9	0	1850	24	-4	7-7	6-9	1-4	2031
control11	24	-5	1-6	6-9	0	2751	24	-4	1-6	6-9	6-5	3167
gpp250-4	15	-5	6-8	2-12	0	77	15	-4	2-8	1-13	3-4	169
gpp500-4	16	-6	3-8	4-14	0	520	15	-3	2-7	8-14	6-4	1450
hinf15	25	-4	9-5	2-12	0	15	24	-4	1-4	2-12	0	11
mcp250-1	13	-6	4-12	4-16	6-7	32	14	-6	1-11	4-16	7-7	97
mcp500-1	15	-7	2-11	5-16	2-7	185	16	-7	3-11	5-16	4-7	1077
qap9	15	-5	5-8	7-13	0	51	15	-5	5-8	4-13	0	53
qap10	14	-5	4-8	7-13	0	101	14	-5	4-8	1-12	0	101
ss30	17	-5	7-7	2-13	5-6	227	22	-5	1-7	2-13	2-5	428
theta3	14	-8	6-11	1-14	4-8	105	14	-8	2-10	2-14	4-8	121
theta4	15	-7	2-10	2-14	2-7	423	14	-7	3-10	2-14	8-8	429
theta5	15	-7	3-10	3-14	2-7	1380	14	-8	4-10	3-14	4-8	1354
theta6	14	-6	2-10	3-14	6-7	3643	14	-7	6-10	3-14	1-7	3767
truss7	23	-4	2-6	2-13	0	12	21	-3	2-5	2-13	0	17
truss8	16	-4	2-6	9-15	0	96	15	-3	3-6	1-14	0	113
equalG11	17	-6	2-10	3-16	2-6	2046	18	-6	1-8	1-16	3-6	7032
equalG51	20	-6	1-8	4-16	3-7	4776	16	-1	1-7	9-16	1-1	12103
equalG32*	20	-5	4-10	2-16	1-5	17588	14	-1	2-7	7-15	1-1	47067
maxG11	14	-6	2-11	7-16	4-6	562	14	-6	5-11	7-16	8-7	3334
maxG51	17	-5	4-11	3-16	2-5	2072	16	-5	7-11	4-16	4-6	7981
maxG32*	15	-5	9-11	1-15	6-6	3578	15	-6	2-10	1-15	2-6	30521
qpG11	14	-5	1-11	0	1-5	4220	15	-5	1-10	0	3-5	18280
qpG112	15	-7	1-11	0	4-7	602	14	-5	9-11	0	1-5	3382
qpG51*	21	-4	5-11	0	1-4	6850	24	-4	8-10	0	3-4	31682
qpG512	19	-4	3-10	0	2-5	2086	27	-4	1-9	0	2-4	13993
thetaG11	18	-7	1-9	1-14	4-7	2103	18	-7	5-10	2-14	2-7	5289
thetaG11n	15	-7	1-12	2-13	4-7	1559	15	-7	1-12	2-13	4-7	4606
thetaG51*	33	-6	6-8	1-14	2-6	19426	30	-5	1-8	2-13	7-6	22299
thetaG51n*	19	-5	2-9	3-13	0	5941	22	-5	2-9	3-13	0	11541

Table 6: Computational results on SDPLIB problems using SDPT3-3.0. These were performed on a 170MHz Sun Ultra 170 with 512MB of memory, except for the problems marked with an asterisk which were on (1 processor of) a 420MHz Sun Ultra 420R with 2GB of memory.

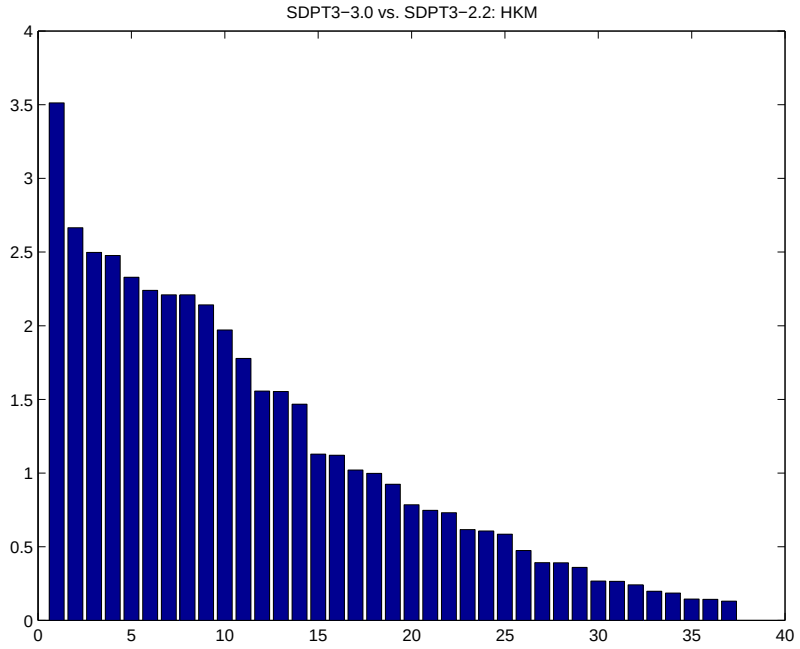

Figure 1. Comparing SDPT3-3.0 and SDPT3-2.2 using the HKM search direction. Bars above the axis demonstrate a win for 3.0.

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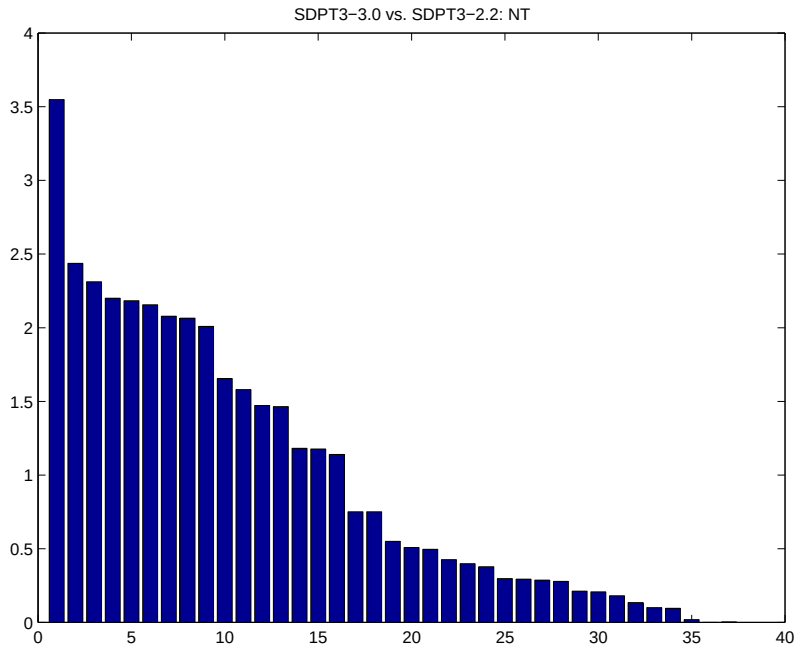

Figure 2. Comparing SDPT3-3.0 and SDPT3-2.2 using the NT search direction. Bars above the axis demonstrate a win for 3.0.

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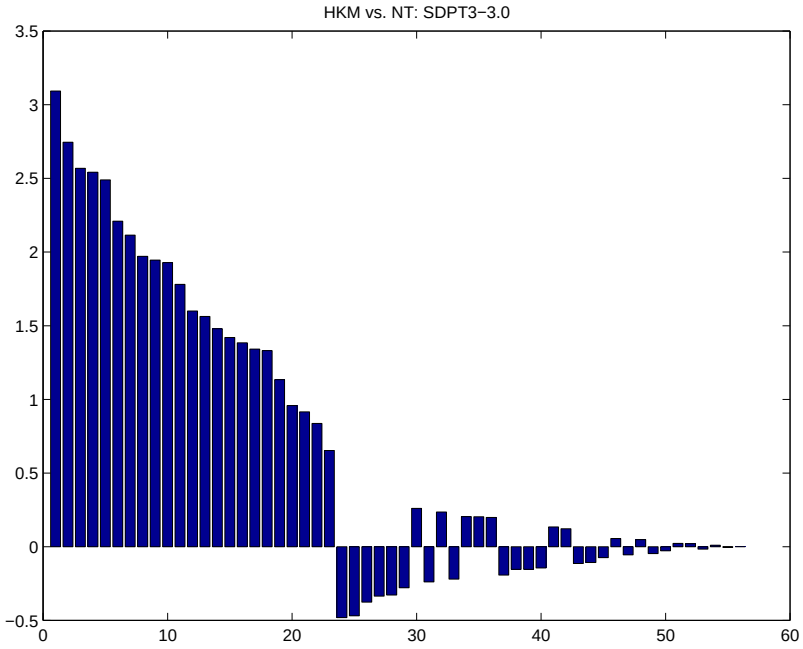

Figure 3. Comparing the HKM and NT search directions. Bars above the axis demonstrate a win for the HKM direction.

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