

Separation of largest degrees,

Graph isomorphism and

edge coloring.

Lemma

Let $k = (n-1)p + x\sqrt{(n-1)pq}$, p constant, $q=1-p$,
where $x \leq (\log n)^2$ (for convenience).

Then

$$B_k = \binom{n-1}{k} p^k (1-p)^{n-1-k} = (1+o(1)) \sqrt{\frac{1}{2\pi npq}} e^{-x^2/2}.$$

Proof

Stirling's Formula gives

$$B_k = (1+o(1)) \sqrt{\frac{1}{2\pi npq}} \left(\frac{(n-1)p}{k} \right)^{\frac{k}{n-1}} \left(\frac{(n-1)q}{n-1-k} \right)^{1-\frac{k}{n-1}}.$$

New

$$\left(\frac{k}{(n-1)p} \right)^{k/n-1} = \left(1 + \alpha \sqrt{\frac{q}{p(n-1)}} \right)^{k/n-1}$$

$$= \exp \left\{ \left(\alpha \sqrt{\frac{q}{p(n-1)}} - \frac{\alpha^2}{2} \frac{q}{p(n-1)} + O(n^{-3/2}) \right) \left(p + \alpha \sqrt{\frac{pq}{n-1}} \right) \right\}$$

$$= \exp \left\{ \alpha \sqrt{\frac{pq}{n-1}} + \frac{\alpha^2}{2} \frac{q}{n-1} + O(n^{-3/2}) \right\}$$

$$\left(\frac{n-1-k}{(n-1)q} \right)^{1-k/n-1} = \left(1 - \alpha \sqrt{\frac{p}{q(n-1)}} \right)^{1-k/n-1}$$

$$= \exp \left\{ - \left(\alpha \sqrt{\frac{p}{q(n-1)}} + \frac{\alpha^2}{2} \cdot \frac{p}{q(n-1)} + O(n^{-3/2}) \right) \left(q - \alpha \sqrt{\frac{pq}{n-1}} \right) \right\}$$

$$= \exp \left\{ - \alpha \sqrt{\frac{pq}{n-1}} + \frac{\alpha^2}{2} \frac{p}{n-1} + O(n^{-3/2}) \right\}$$

So

$$\left(\frac{k}{(n-1)p} \right)^{\frac{k}{n-1}} \left(\frac{n-1-k}{(n-1)q} \right)^{1 - \frac{k}{n-1}} =$$

$$\exp \left\{ \frac{\partial c^2}{2(n-1)} + O(n^{-3/2}) \right\}$$

Substituting into

$$(1+o(1)) \sqrt{\frac{1}{2\pi n p q}} \left(\frac{(n-1)p}{k} \right)^{\frac{k}{n-1}} \left(\frac{(n-1)q}{n-1-k} \right)^{1 - \frac{k}{n-1}}^{n-1}$$

gives required expression.

□

Lemma

Let $\epsilon = \frac{1}{10}$ and p be constant

$$k_{\pm} = (n-1)p \pm (1 \pm \epsilon) \sqrt{2(n-1)pq \log n}.$$

Then whp

$$(i) \quad \Delta(G_{n,p}) \leq k_{+}$$

(ii) There are $\Omega(n^{2\epsilon(1-\epsilon)})$ vertices of degree at least k_{-} .

(iii) ~~\exists~~ $u \neq v$ such that $d(u), d(v) \geq k_{-}$
and $|d(u) - d(v)| \leq 10$.

We first prove that as $x \rightarrow \infty$

$$\frac{1}{x} e^{-x^2/2} \left(1 - \frac{1}{x^2}\right) \leq \int_x^\infty e^{-y^2/2} dy \leq \frac{1}{x} e^{-x^2/2}. \quad (***)$$

Proof

$$\int_x^\infty e^{-y^2/2} dy = - \int_x^\infty \frac{1}{y} (e^{-y^2/2})' dy$$

$$= - \left[\frac{1}{y} e^{-y^2/2} \right]_x^\infty - \int_x^\infty \frac{1}{y^2} e^{-y^2/2} dy$$

$$= \frac{1}{x} e^{-x^2/2} + \left[\frac{1}{y^3} e^{-y^2/2} \right]_x^\infty + 3 \int_x^\infty \frac{1}{y^4} e^{-y^2/2} dy \quad \square$$

(i) Let X be the number of vertices of degree k .

$$E(X_k) = (1+o(1)) \sqrt{\frac{n}{2\pi pq}} \exp\left\{-\frac{1}{2} \left(\frac{k - (n-1)p}{\sqrt{(n-1) pq}}\right)^2\right\}$$

assuming that $k \leq k_2 = (n-1)p + (\log n)^2 \sqrt{(n-1) pq}$.

But if $k > k_2$ then

$$E(X_k) \leq E(X_{k_2}) \quad - \text{binomial} \rightarrow \text{after mean}$$

$$\approx n \exp\left\{-\Omega((\log n)^4)\right\}$$

$$= o(1).$$

So if $Y_k = X_k + X_{k+1} + \dots$

$$E(Y_k) \approx \sum_{l=k}^{k_L} \sqrt{\frac{n}{2\pi pq}} \exp\left\{-\frac{1}{2} \left(\frac{l - (n-1)p}{\sqrt{(n-1) pq}}\right)^2\right\}$$

$$\approx \sum_{l=k}^{\infty} \sqrt{\frac{n}{2\pi pq}} \exp\left\{-\frac{1}{2} \left(\frac{l - (n-1)p}{\sqrt{(n-1) pq}}\right)^2\right\}$$

$$\approx \sqrt{\frac{n}{2\pi pq}} \int_{\lambda=k}^{\infty} \exp\left\{-\frac{1}{2} \left(\frac{\lambda - (n-1)p}{\sqrt{(n-1) pq}}\right)^2\right\} d\lambda$$

If $k = (n-1)p + x \sqrt{(n-1)pq}$ then

$$\sqrt{\frac{n}{2\pi pq}} \int_{\lambda=k}^{\infty} \exp \left\{ -\frac{1}{2} \left(\frac{\lambda - (n-1)p}{\sqrt{(n-1)pq}} \right)^2 \right\} d\lambda$$

$$= \sqrt{\frac{n}{2\pi pq}} \cdot \sqrt{(n-1)pq} \cdot \int_{y=x}^{\infty} e^{-y^2/2} dy$$

$$\approx \frac{n}{\sqrt{2\pi}} \cdot \frac{1}{x} \cdot e^{-x^2/2}$$

When $k = k_+$, $x = (1+\epsilon) \sqrt{2 \log n}$ and (i) follows.

When $k = k_-$, $x = (1 - \epsilon) \sqrt{2 \log n}$

and $E(Y_{k_-}) = \Omega(n^{2\epsilon(1-\epsilon)}) \rightarrow \infty$.

We use the second moment method to show concentration.

$$E(Y_k(Y_{k-1} - 1)) = n(n-1) \sum_{k \leq k_1, k_2 \leq k_-} P_r(d(1) = k_1 \wedge d(2) = k_2)$$

$$= n(n-1) \left[\sum_{k_1, k_2} P(\hat{d}(1) = k_1 - 1 \wedge \hat{d}(2) = k_2 - 1) + (1-p) P(\hat{d}(1) = k_1 \wedge \hat{d}(2) = k_2) \right]$$

where $\hat{d} = \# \text{nbrs in } \{3, 4, \dots, n\}$.

$$= n(n-1) \sum_{k_1, k_2} \left[p P(\hat{d}(1) = k_1 - 1) P(\hat{d}(2) = k_2 - 1) + (1-p) P(\hat{d}(1) = k_1) P(\hat{d}(2) = k_2) \right]$$

$$\frac{P(\hat{d}(1) = k_1 - 1)}{P(\hat{d}(1) = k_1)} = \frac{\binom{n-2}{k_1-1} (1-p)}{\binom{n-2}{k_1} p} = \frac{k_1 (1-p)}{(n-2-k_1) p} = 1 + \tilde{O}(n^{-1/2}).$$

$$= n(n-1) \sum_{k_1, k_2} \left[P(\hat{d}(1) = k_1) P(\hat{d}(2) = k_2) (1 + \tilde{O}(n^{-1/2})) \right]$$

$$= n(n-1) \sum_{k_1, k_2} \left[P(\hat{d}(1) = k_1) P(\hat{d}(2) = k_2) (1 + \tilde{O}(n^{-1/2})) \right]$$

$$\frac{P(\hat{d}(1) = k_1)}{P(d(1) = k_1)} = \frac{\binom{n-2}{k_1}}{\binom{n}{k_1}} (1-p)^{-2} = 1 + \tilde{O}(n^{-1/2})$$

$$= n(n-1) \sum_{k_1, k_2} \left[P(d(1) = k_1) P(d(2) = k_2) (1 + \tilde{O}(n^{-1/2})) \right]$$

$$= E(Y_k) (E(Y_k) - 1) (1 + \tilde{O}(n^{-1/2}))$$

So, with $k = k_-$,

$$P_r(Y_{k_-} \leq \frac{1}{2} E(Y_{k_-}))$$

$$\leq \frac{E(Y_{k_-}(Y_{k_-} - 1)) + E(Y_{k_-}) - E(Y_{k_-})^2}{E(Y_{k_-})^2 / 4}$$

$$= O\left(\frac{1}{n^{2\epsilon(1-\epsilon)}}\right)$$

$$= o(1).$$

This completes the proof of the second part.

$$P_r(\neg(iij)) \leq o(i) + \binom{n}{2} \sum_{k_1=k_2}^{k_L} \sum_{|k_2-k_1| \leq 10} P_r(d(1)=k_1 \wedge d(2)=k_2)$$

$$= o(i) + \sum_{k_1, k_2} \binom{n}{2} \left[p P(\hat{d}(1)=k_1-1) P(\hat{d}(2)=k_2-1) + (1-p) P(\hat{d}(1)=k_1) P(\hat{d}(2)=k_2) \right]$$

Now

$$\sum_{k_1, k_2} P(\hat{d}(1)=k_1-1) P(\hat{d}(2)=k_2-1)$$

$$\leq 21 (1 + \tilde{O}(n^{-1/2})) \sum_{k_1} P_r(\hat{d}(1)=k_1-1)^2$$

and

$$\sum_{k_1} P_r(\hat{d}^{(1)} = k_1 - 1)^2 \lesssim \frac{1}{2\pi p q n} \int_{y=x}^{\infty} e^{-y^2} dy,$$

where $x = \frac{k_1 - (n-1)p}{\sqrt{(n-1)pq}} \lesssim (1-\epsilon)\sqrt{2} \log n$

$$= \frac{1}{\sqrt{8\pi p q n}} \int_{z=x\sqrt{2}}^{\infty} e^{-z^2/2} dz$$

$$\lesssim \frac{1}{\sqrt{8\pi p q n}} \cdot \frac{1}{x\sqrt{2}} \cdot e^{-2(1-\epsilon)^2}$$

We get a similar bound for $\sum_{k_1} P_r(\hat{d}^{(1)} = k_1)^2$.

Thus

$$\begin{aligned} P_r(\rightarrow (iii)) &= o\left(n^{2-1-2(1-\epsilon)^2}\right) \\ &= o(1). \end{aligned}$$

□

Edge Colouring

The **Chromatic Index** $\chi'(G)$ of graph G is the minimum number of colors that can be used to color the edges of G so that if 2 edges share a vertex, they have a different color.

Vizing's Theorem states that

$$\Delta(G) \leq \chi'(G) \leq \Delta(G) + 1.$$

Also, if there is a unique vertex of maximum degree, then $\chi'(G) = \Delta(G)$.

So $\chi'(G_{n,p}) = \Delta(G_{n,p})$ whp.

Graph Isomorphism

In this section we describe a procedure for ordering the vertices of a graph G .

If it succeeds then it is possible

to quickly tell if $G \cong H$, for **any** H .

Algorithm

Input G . Parameter L .

Step 1

Re-label vertices so that degrees satisfy

$$d_G(v_1) \geq d_G(v_2) \geq \dots \geq d_G(v_n)$$

If $\exists i \leq L$ such that $d_G(v_i) = d_G(v_{i+1})$: **FAIL**

Step 2

For $i > L$ let:

$$X_i = \{ j \in \{1, 2, \dots, L\} : (v_i, v_j) \in G \}$$

Re-label vertices so that these sets satisfy

$$X_{L+1} \supseteq X_{L+2} \supseteq \dots \supseteq X_n \quad \text{— lexicographic ordering.}$$

If $\exists i > L$ such that $X_i = X_{i+1}$: **FAIL**.

Suppose now that the above algorithm succeeds for G .

Given an n -vertex graph H we run the algorithm on H .

(i) If algorithm fails $G \not\cong H$.

(ii) Suppose ordering of $V(H)$ is w_1, w_2, \dots, w_n . Then

$G \cong H \iff v_i \rightarrow w_i$ is an isomorphism.

Claim

Let $\rho = p^2 + q^2$ and $L = 3 \log_{1/\rho} n$.

Then whp the algorithm succeeds on $G = G_{n,p}$.

Proof

We have already proved that Step 1 succeeds whp.

We must now show that $X_i \neq X_j \neq i, j$ whp but there is slight problem because edges (v_i, v_j) are conditioned due to us knowing v_i has a high degree.

Fix i, j and let $\widehat{G}_{i,j} = G \setminus \{i, j\}$.

Now if i, j are not high degree vertices then the L largest degree vertices in $G, \widehat{G}_{i,j}$ will coincide, whp.

This is because there is whp, a gap ≥ 10 between high vertex degrees in G .

Thus

$$P_r(\text{Step 2 fails}) \leq$$

$$o(1) + \sum_{1 \leq i < j \leq n} P_r(i, j \text{ have same nbors among } L \text{ largest degree vertices in } \widehat{G}_{i,j})$$

$$= o(1) + \binom{n}{2} \rho^L$$

$$= o(1).$$



Automorphisms

It follows from the previous section that
whyp, $G_{n,p}$ has no non-trivial automorphisms.

For \downarrow $\sigma: [n] \rightarrow [n]$ is an automorphism, then

(i) $\sigma(v_i) = v_i, 1 \leq i \leq 2$

where v_i is the vertex with the i th largest degree.

(ii) $\sigma(v) = v$ for $v \notin \{v_1, v_2, \dots, v_2\}$.

This is because all of the sets X_v are
distinct.