

# Random Graphs with a Fixed Degree Sequence.

Let  $\underline{d} = (d_1, d_2, \dots, d_n)$

where  $d_1 + d_2 + \dots + d_n = 2M$  is even.

Let  $G_{n, \underline{d}} = \left\{ \begin{array}{l} \text{simple graphs with vertex set } [n] \\ \text{such that degree } d(i) = d_i, i \in [n] \end{array} \right\}$

$G_{n, \underline{d}}$  is chosen randomly from  $G_{n, \underline{d}}$ .

We assume that  $d_1, d_2, \dots, d_n \geq 1$  and that  $\sum d_i (d_i - 1) = \Omega(n)$ .

## Configuration model

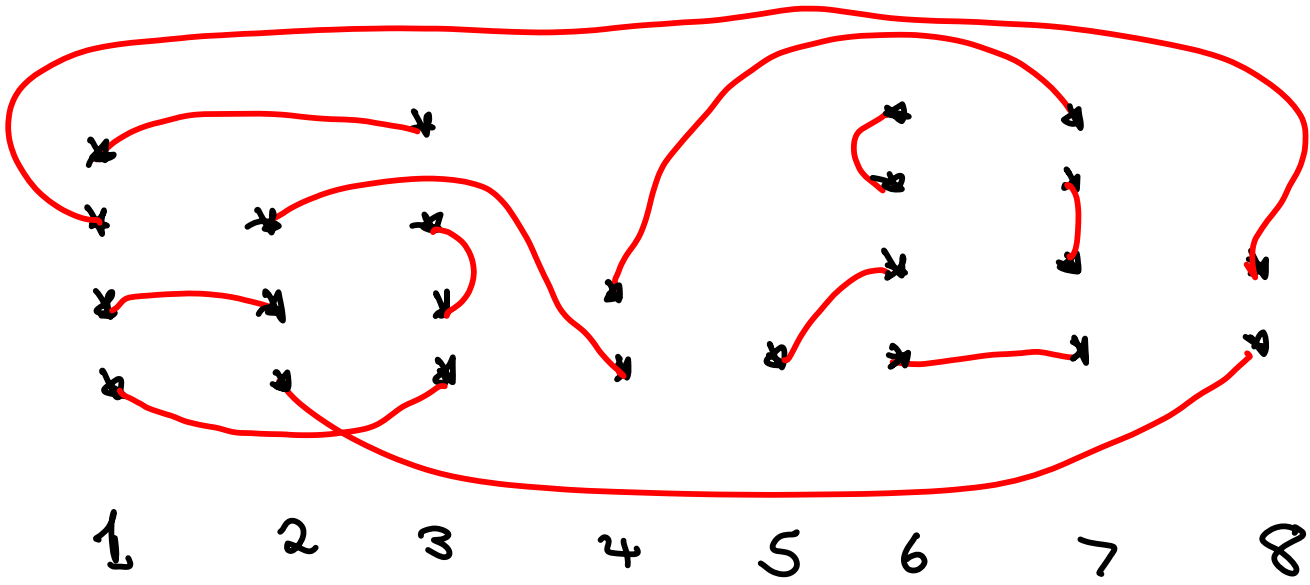
Let  $W_1, W_2, \dots, W_n$  be a partition of  $W$ , where  
 $|W_i| = d_i$  for  $1 \leq i \leq n$ .

For  $x, y \in W$  define  $\phi(x)$  by  $x \in W_{\phi(x)}$ .

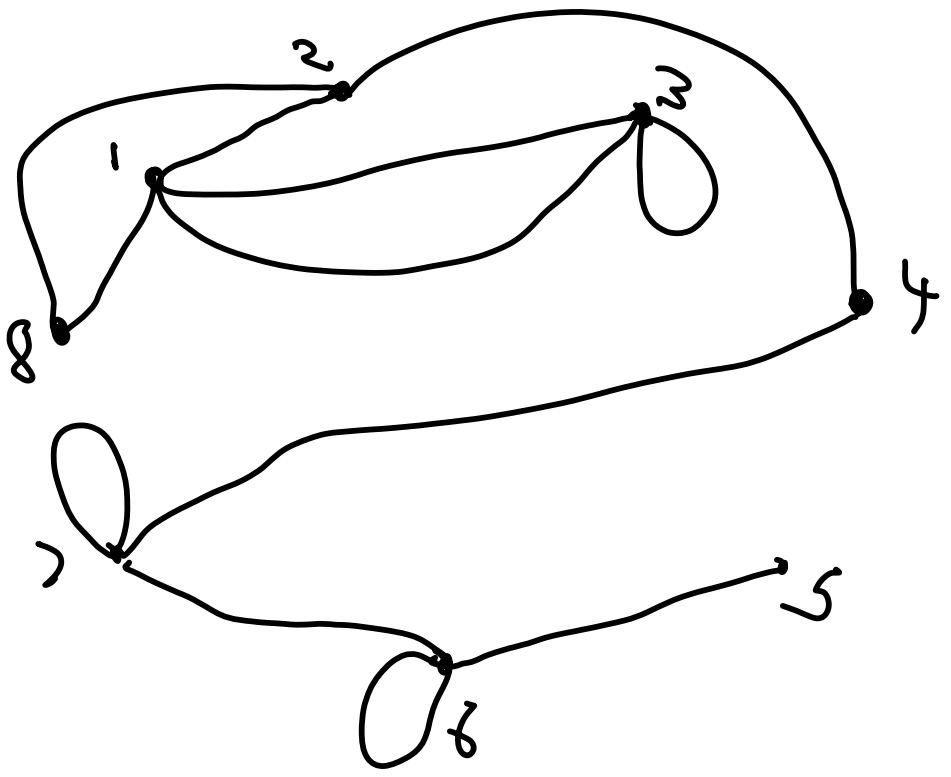
Let  $F$  be a partition of  $W$  into  $m$  parts  
(a configuration).

Given  $F$  we define the (multi) graph  $\gamma(F)$

$$\gamma(F) = ([n], \{(\phi(x), \phi(y)) : (x, y) \in F\})$$



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## Lemma

If  $G \in \mathcal{G}_{n, \underline{d}}$ , then

$$|\gamma^{-1}(G)| = \prod_{i=1}^n d_i!$$

## Proof

Arrange the edges of  $G$  in lexicographic order. Now go through the sequence of  $2m$  symbols, replacing each  $i$  by a new member of  $W_i$ . We get all  $F$  for which  $\gamma(F) = G$ .

□

## Corollary

If  $F$  is chosen uniformly at random from  $\Omega$  (the set of configurations) and  $G_1, G_2 \in \mathcal{G}_{n,d}$  then

$$P_r(\chi(F) = G_1) = P_r(\chi(F) = G_2).$$

so we can choose a random  $F$  and accept  $\chi(F)$  iff there are no loops or multiple edges.

The next question is: What is

$P_r(\gamma|F)$  is simple)

(i)

Since

$$|\Omega| = \frac{(2m)!}{m! 2^m}$$

(ii)

Take  $d_i$  "distinct" copies of  $i$  for  $i=1, 2, \dots, n$  and take a permutation of these  $2m$  symbols. Read off  $F$ , pair by pair.

Each distinct  $F$  arises in  $m! 2^m$  ways.

(i) & (ii) will tell us how big is  $G_{n,d}$ .

# Alternative Construction of F

Begin

$$U = W; F = \emptyset;$$

For  $i = 1, 2, \dots, m$  do

Begin

Choose  $x$  *arbitrarily* from  $U$ ;

Choose  $y$  *randomly* from  $U \setminus \{x\}$ ;

$$F := F \cup \{(x, y)\};$$

$$U := U \setminus \{x, y\}$$

End

Each  $F$  arises with probability  $\frac{1}{(2m-1)(2m-3)\dots 1}$   
 $= |\Omega|^{-1}$ .

Let  $\Delta = \max \{d_1, d_2, \dots, d_n\}$ .

Notation:  $F \stackrel{\text{ran}}{\in} \Omega \equiv F$  is chosen uniformly from  $\Omega$ .

Lemma

Assume that  $\Delta \leq n^{1/6}$  and  $F \stackrel{\text{ran}}{\in} \Omega$ . Then whp

(a)  $\gamma(F)$  has no double loops

(b)  $\gamma(F)$  has  $\leq \Delta \log n$  loops

(c)  $\gamma(F)$  has no triple edges.

(d)  $\gamma(F)$  has no adjacent double edges.

(e)  $\gamma(F)$  has  $\leq \Delta^2 \log n$  double edges.



Proof

(a)

$P_r(F \text{ contains a double loop})$

$$\sum_{i=1}^n \binom{d_i}{4} \cdot 3 \cdot \left( \frac{1}{2m-3} \right)^2$$

$$\leq n \Delta^4 m^{-2}$$

$$= O(1),$$

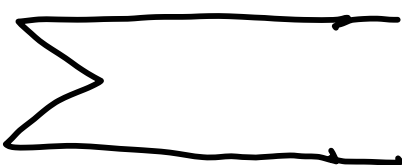
(b)

Let  $k_1 = \Delta \log n$ .

$$Pr(F \text{ has } \geq k_1 \text{ loops}) \leq o(1) +$$

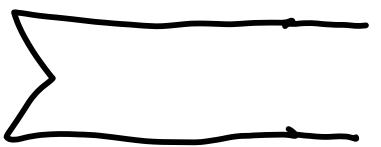
$$\sum_{l=1}^n \binom{d_i}{2}^{x_i} \left( \frac{1}{2m - 2k_1} \right)^{k_1}$$

$x_1 + \dots + x_n = k_1$   
 $x_i = 0, 1$



$$\leq o(1) + \left( \frac{\Delta}{2m} \right)^{k_1} \sum_{l=1}^n d_i^{x_i}$$

$x_1 + x_2 + \dots + x_n = k_1$



$$\leq o(1) + \left( \frac{\Delta}{2m} \right)^{k_1} \frac{(d_1 + \dots + d_n)^{k_1}}{k_1!} \leq o(1) + \left( \frac{\Delta e}{k_1} \right)^{k_1} = o(1).$$

(c)

$P_2(F \text{ contains a triple edge})$

$$\approx \sum_{1 \leq i < j \leq n} \binom{d_i}{3} \binom{d_j}{3} \cdot 6 \cdot \left( \frac{1}{2m-5} \right)^3$$

$$\approx \Delta^4 \left( \sum d_i \right)^2 n^{-3}$$

$$= o(1).$$


(d)

$P_r(F \text{ contains 2 adjacent double edges}) \leq$

$$\sum_{i=1}^n \binom{d_i}{2} \left( \frac{\Delta}{2m-8} \right)^2$$

$$\leq \frac{\Delta^3}{(2m-8)^2} \sum_{i=1}^n d_i$$

$$= o(1).$$

  $P_r \leq \frac{\Delta}{2m-8}$   
that  $a, b$   
are in same  $W_{ij}$ .

(e) Let  $k_2 = \Delta^2 \log n$ .

$P_r(F \text{ has } \geq k_2 \text{ double edges})$

$$o(1) + \sum_{\substack{\alpha_1 + \dots + \alpha_n = k_2 \\ \alpha_i = 0, 1}} \prod_{i=1}^n \left[ \binom{d_i}{2} \frac{\Delta}{2m - 4k_2} \right]^{\alpha_i}$$

$$o(1) + \left( \frac{\Delta^2}{m} \right)^{k_2} \sum_{\substack{\alpha_1 + \dots + \alpha_n = k_2 \\ \alpha_i = 0, 1}} \prod_{i=1}^n d_i^{\alpha_i}$$

$$\leq o(1) + \left( \frac{\Delta^2}{m} \right)^{k_2} \frac{(2m)^{k_2}}{k_2!}$$

$$= o(1).$$



# Switchings

Let now

$$\Omega_{i,j} = \{ F \in \Omega : F \text{ has } i \text{ loops, } j \text{ double edges and no double loops or triple edges and no vertex incident with 2 double edges} \}$$

Lemma (Sutherland) Let  $M_1 = 2M$  and  $M_2 = \sum_i d_i(d_i - 1)$

For  $i \leq k_1$ , and  $j \leq k_2$

$$\frac{|\Omega_{i-1, j}|}{|\Omega_{i, j}|} = \frac{2iM_1}{M_2} \left( 1 + \mathcal{O}\left(\frac{\Delta^3}{\hbar}\right) \right)$$

$$\frac{|\Omega_{0, j-1}|}{|\Omega_{0, j}|} = \frac{4jM_1^2}{M_2^2} \left( 1 + \mathcal{O}\left(\frac{\Delta^3}{\hbar}\right) \right).$$

Corollary

$$\frac{|\Omega_{0,0}|}{|\Omega|} = (1 + o(1)) e^{-\lambda(\lambda+1)}$$

where  $\lambda = \frac{M_2}{2M_1}$ .

Thus

$$|G_{n,d}| \approx e^{-\lambda(\lambda+1)} \frac{1}{\prod_{L=1}^n d_L!} \frac{(2m)!}{m! 2^m}$$



Proof It follows from the switching lemma that  $i \leq k_1$  and  $j \leq k_2$  implies

$$\frac{|\Omega_{i,j}|}{|\Omega_{0,0}|} = (1 + o(1)) \frac{\lambda^{i+2j}}{i! j!}$$

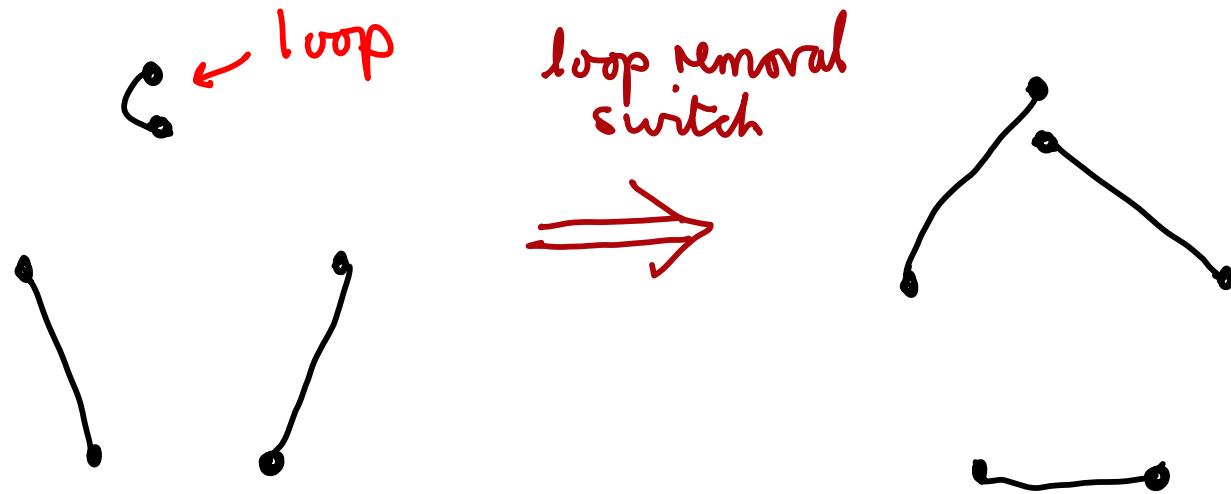
Therefore

$$\begin{aligned} (1 + o(1)) |\Omega| &= (1 + o(1)) |\Omega_{0,0}| \sum_{i=0}^{\lambda_1} \sum_{j=0}^{\lambda_2} \frac{\lambda^{i+2j}}{i! j!} \\ &= (1 + o(1)) |\Omega_{0,0}| e^{\lambda(\lambda+1)}. \end{aligned}$$

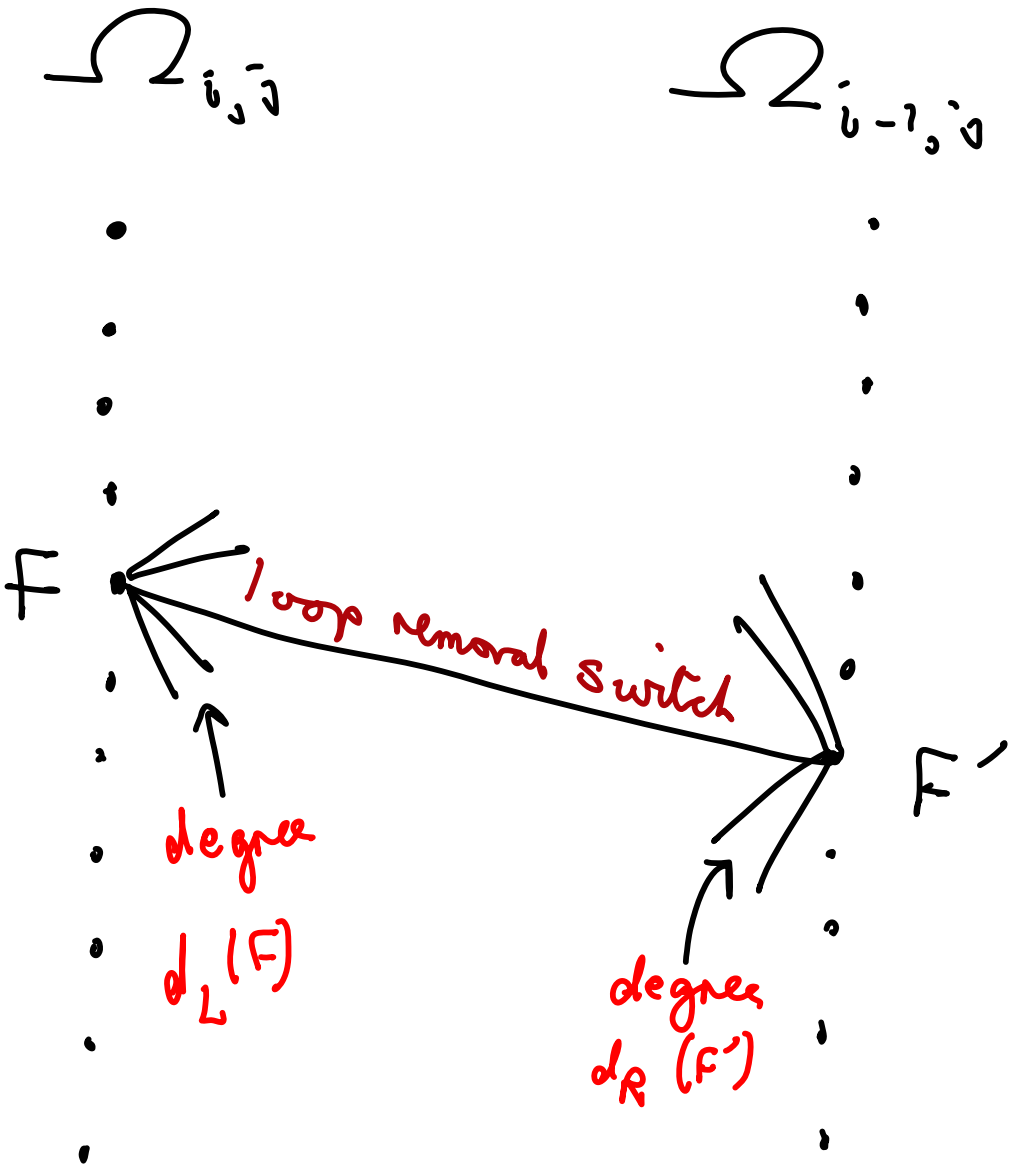


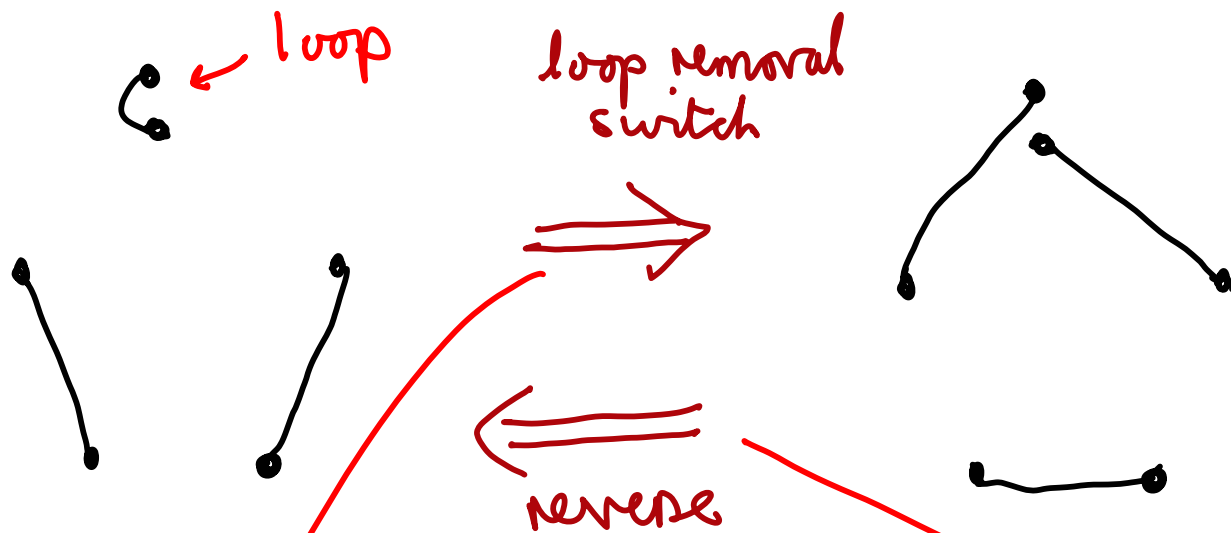
# Proof of switching lemma

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In general this operation takes a member  $F$  of  $\Omega_{i,j}$  to a member  $F'$  of  $\Omega_{i-1,j}$  unless it creates new loops or multiple edges.





$$\begin{aligned} \# \text{choices} &\leq i \times M_1^2 \\ &\geq i \times M_1^2 - \tilde{O}(i M_1 \Delta^2) \end{aligned}$$

$$\begin{aligned} \# \text{choices} &\leq M_1 M_2 / 2 \\ &\geq M_1 M_2 / 2 - \tilde{O}((M_1 + M_2) \Delta^3) \end{aligned}$$

Now

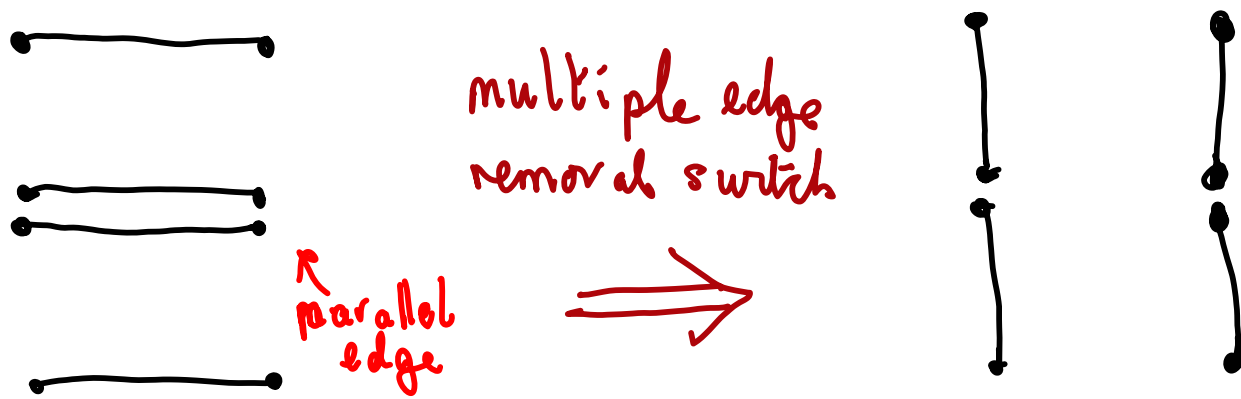
$$\sum_{F \in \Omega_{i,j}} d_L(F) = \sum_{F' \in \Omega_{i-1,j}} d_R(F')$$

$$\begin{aligned} & \approx iM_1^2 |\Omega_{i,j}| \\ & \geq iM_1^2 |\Omega_{i,j}| \\ & \times (1 - \tilde{O}(i\Delta^2/M_1)) \end{aligned}$$

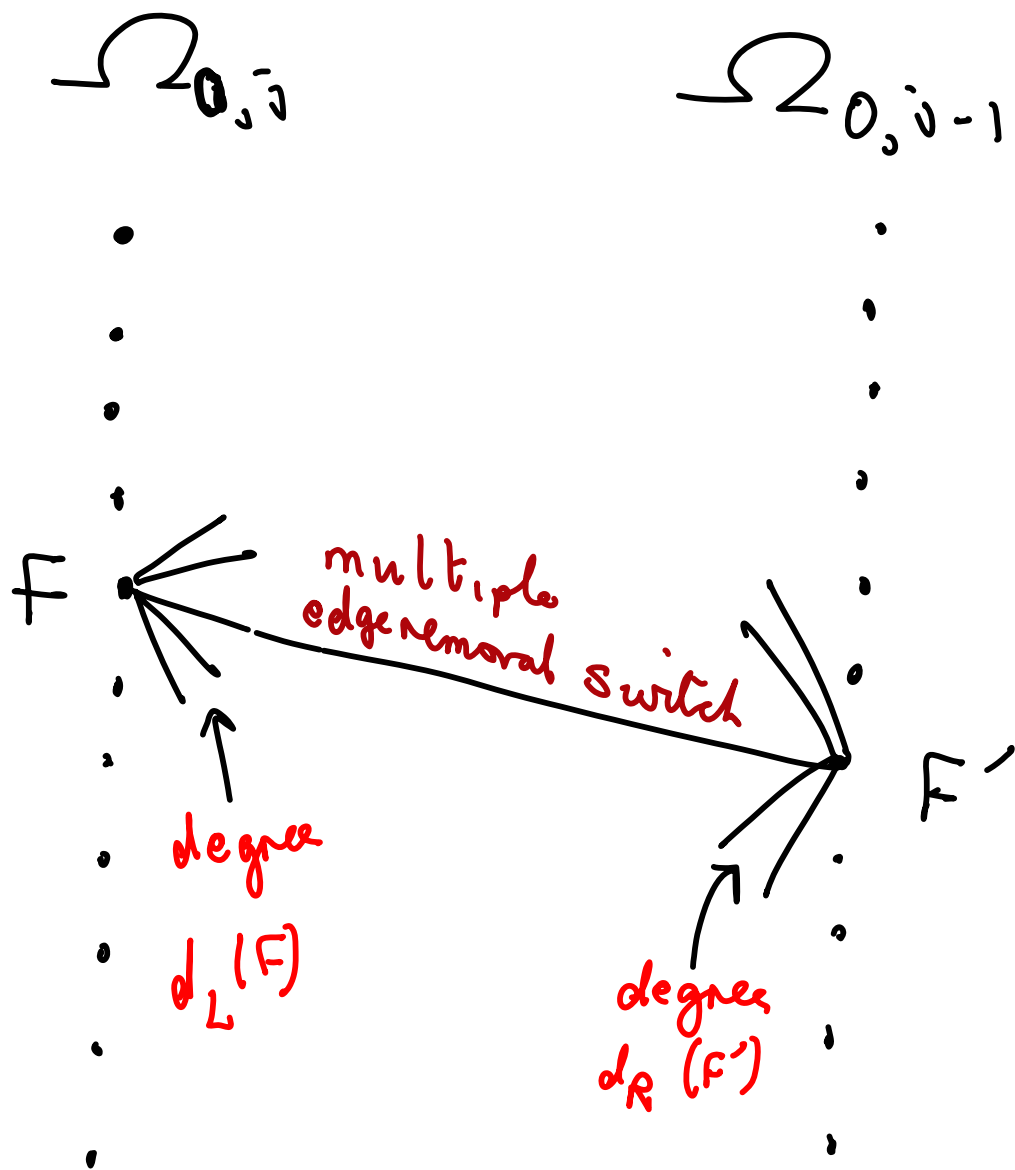
$$\begin{aligned} & \leq \frac{1}{2} M_1 M_2 |\Omega_{i-1,j}| \\ & \geq \frac{1}{2} M_1 M_2 |\Omega_{i-1,j}| \\ & \times (1 - \tilde{O}\left(\frac{\Delta^3}{M_1} + \frac{\Delta^3}{M_2}\right)) \end{aligned}$$

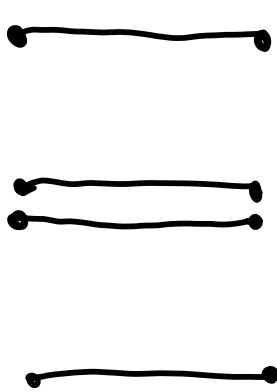
So

$$\frac{|\Omega_{i-1,j}|}{|\Omega_{i,j}|} = \frac{2iM_1}{M_2} \left(1 + \tilde{O}\left(\frac{\Delta^3}{M_1}\right)\right)$$



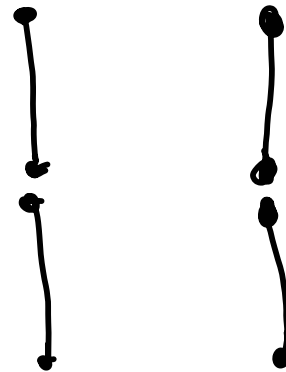
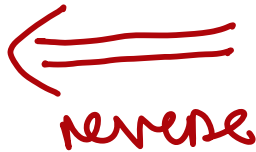
In general this operation takes a member  $F$  of  $\Omega_{i,j}$  to a member  $F'$  of  $\Omega_{i,j-1}$  unless it creates new loops or multiple edges.





↑ parallel edge

multiple edge removal switch



$$\begin{aligned} \# \text{choices} &: \approx j \times M_1^2 \\ &\geq j \times M_1^2 - \tilde{O}(j M_1 \Delta^2) \end{aligned}$$

$$\begin{aligned} \# \text{choices} &\leq M_2^2/4 \\ &\geq M_2^2/4 - \tilde{O}(M_2 \Delta^3) \end{aligned}$$



Now

$$\sum_{F \in \Omega_{0,j}} d_L(F) = \sum_{F' \in \Omega_{0,j-1}} d_R(F')$$

$$\leq j M_1^2 |\Omega_{0,j}|$$

$$\geq j M_1^2 |\Omega_{0,j}| \times (1 - \tilde{O}\left(\frac{\Delta^2}{M_1}\right))$$

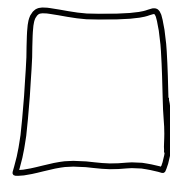
$$\leq \frac{1}{4} M_2^2 |\Omega_{0,j-1}|$$

$$\leq \frac{1}{4} M_2^2 |\Omega_{0,j-1}|$$

$$\times \left(1 - \tilde{O}\left(\frac{\Delta^3}{M_2}\right)\right)$$

So

$$\frac{|\Omega_{0,j-1}|}{|\Omega_{0,j}|} = \frac{4j M_1^2}{M_2^2} \left(1 + \tilde{O}\left(\frac{\Delta^3}{M_2}\right)\right)$$



So  $\forall F \in \Omega,$

$$P_r(\chi(F) \text{ is simple}) \approx e^{-\lambda(\lambda+1)}$$

where  $\lambda = \frac{\sum d_i(d_i-1)}{2 \sum d_i}$

So for any (multi) graph property  $\mathcal{P}$

$$P_r(G_{n,d} \in \mathcal{P}) \leq (1+o(1)) e^{-\lambda(\lambda+1)} P_r(\chi(F) \in \mathcal{P})$$

assuming  $\Delta \leq n^{1/3}$  [Not best known.]

$$\Pr(G_{n,d} \in \mathcal{P}) \leq (1+o(1)) e^{\lambda(\lambda+1)} \Pr(\mathcal{Y}(F) \in \mathcal{P})$$

This is particularly useful if  $\lambda = O(1)$

e.g. random  $r$ -regular graphs where

$r$  is a constant. Here  $\lambda = \frac{r-1}{2}$ .

## Theorem

Let  $G_{n,r}$  denote a random  $r$ -regular graph,  $r \geq 3$  constant, vertex set  $[n]$ .

Then whp  $G_{n,r}$  is  $r$ -connected.

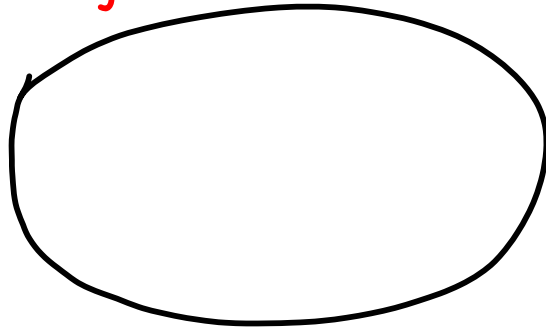
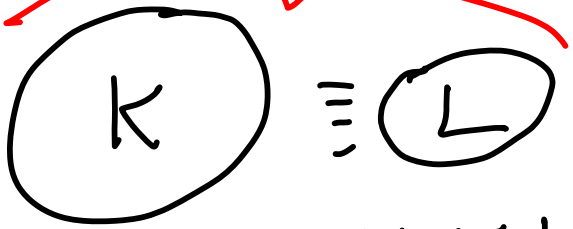
## Corollary

If  $n$  is even then whp  $G_{n,r}$  has a perfect matching.

[An  $r$ -edge connected,  $r$ -regular graph, with  $n$  even, has a perfect matching.]

Proof

$$\left. \begin{array}{l} \# \text{edges} \leq k+l \\ \text{incident with} \\ k \end{array} \right\} \geq \frac{r_{k+l}}{2}$$



$$l = |L| \leq r-1$$

$$L = N(K)$$

$r_0$

$$(i) \cdot \frac{r}{r-2} \leq k = |K| \leq n e^{-10}$$

$$P_r(\exists K, L) \leq \sum_{k,l} \binom{n}{k} \binom{n}{l} \left( \frac{r_k}{2} \right) \left( \frac{r_{k+l}}{r_n} \right)^{\frac{r_{k+l}}{2}}$$

$$\leq \sum_{k,l} n^{-\left(\frac{r}{2}-1\right)k + \frac{l}{2}} \frac{e^{k+l}}{k^k l^l} 2^{rk} (k+l)^{\frac{r_{k+l}}{2}}$$

$$\leq \sum_{k,l} n^{-(\frac{r}{2}-1)k + \frac{l}{2}} \frac{e^{k+l}}{k^k l^l} 2^{rk} (k+l)^{\frac{r(k+l)}{2}}$$

$$\left(\frac{k+l}{e}\right)^{l/2} \leq e^{k/2} \quad \left(\frac{k+l}{k}\right)^{k/2} \leq e^{l/2}$$

$$(k+l)^{rk/2} \leq k^{rk/2} e^{lr/2}$$

$$\frac{r-1}{2k} < \frac{r}{2} - 1$$

$$k > \frac{r-1}{r-2}$$

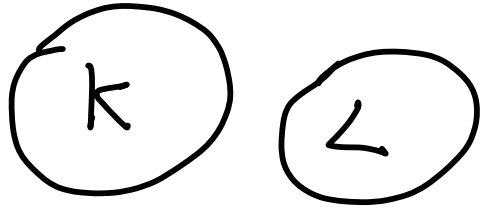
$$\leq C_r \sum_{k=r}^{n/10} \sum_{l=0}^{r-1} n^{-(\frac{r}{2}-1)k + \frac{l}{2}} e^{3k/2} k^{(r-2)k/2}$$

$\uparrow$   
 constant

$$= C_r \sum_{k=r_0}^{n/10} \sum_{l=0}^{r-1} \left( n^{-(\frac{r}{2}-1) + \frac{l}{2k}} e^{3/2} k^{\frac{r}{2}-1} 2^r \right)^k$$

$$= o(1).$$

$$(ii) 2 \leq k \leq \frac{r}{r-2} \leq 3$$



KVL contains

$$\frac{rk}{2} + l \geq k + l + 1 \text{ edges}$$

$P_r(\exists S: s=|S| \leq r-1+3, \text{ contains } s+1 \text{ edges})$

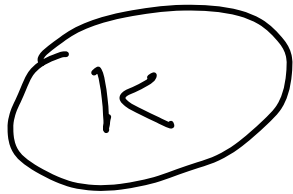
$$\leq \sum_{s=4}^{r+2} \binom{n}{s} \binom{rs/2}{s+1} \left(\frac{rs}{rn}\right)^{s+1}$$

$$\leq \sum_{s=4}^{r+2} n^s \cdot 2^{rs/2} \cdot s^{s+1} \cdot n^{-s-1}$$

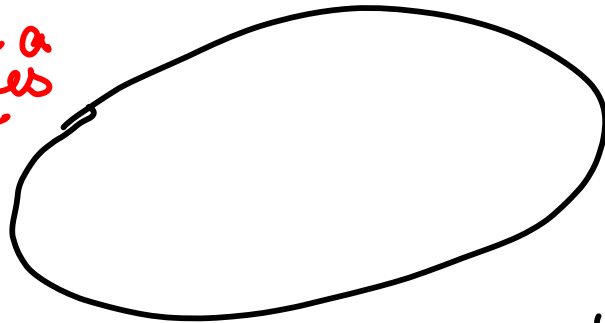
$$= o(1)$$

$$(iii) \quad ne^{-10} < k \leq \frac{n}{2}$$

$$\phi(2m) = \frac{(2m)!}{m! 2^m} \approx 2^{\frac{1}{2}} \left(\frac{2m}{e}\right)^m$$



$r$  edges



$$P_r(\exists K, L) \leq \sum_{k, l, a} \binom{n}{k} \binom{n}{l} \binom{r}{a} \frac{\phi(rk + rl - a) \phi(r(n-k-l) + a)}{\phi(rn)}$$

$$\leq C_r \sum_{k, l, a} \left(\frac{ne}{k}\right)^k \left(\frac{ne}{l}\right)^l \frac{\binom{rk + rl - a}{r} \binom{r(n-k-l) + a}{r}}{(rn)^{rn}}$$

$$\leq C_r \sum_{k, l, a} \left(\frac{ne}{k}\right)^k \left(\frac{ne}{l}\right)^l e^{O(1)} \left(\frac{k}{n}\right)^{rk} \left(1 - \frac{k}{n}\right)^{r(n-k)}$$



$$= C_r \sum_{k, l, a} \left(\frac{ne}{k}\right)^k \left(\frac{ne}{l}\right)^l e^{O(l)} \left(\frac{k}{n}\right)^{rk} \left(1 - \frac{k}{n}\right)^{r(n-k)}$$

$$\leq C_r \sum_{k, l, a} \left( \left(\frac{k}{n}\right)^{r-1} \cdot e^{1-r/2} \cdot n^{\frac{r}{k}} \right)^k$$

$$= O(1)$$

