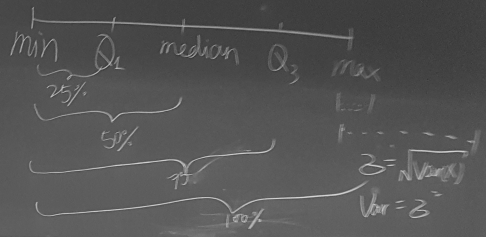


mean = average



### Knapsack Problem

Optimization:

Maximize  $p_1 x_1 + p_2 x_2 + \dots + p_n x_n$   
 Subject to  $w_1 x_1 + w_2 x_2 + \dots + w_n x_n \leq W$   
 $x_1, x_2, \dots, x_n \geq 0$  and are integers

Assumptions that  
 (i)  $0 < w_j \leq W$  for all  $j$   
 (ii)  $p_j > 0$  for all  $j$

- (i) Scout  
 $W$  = volume of knapsack  
 $p_j$  = value of one of item  $j$   
 $w_j$  = volume of " " " " " "
- (ii) Cargo Planes  
 $W$  = weight limit  
 $p_j$  = profit/item  
 $w_j$  = weight/item

- (iii)  $W$  = money to invest  
 $p_j$  = expected profit per unit of stock  
 $w_j$  = cost per unit of stock

Look at a dynamic programming formulation

Knapsack Problem

- ①  $1 \leq w_j \leq W$  for all  $j$
- ②  $p_j > 0$  for all  $j$

Optimization:

Maximize  $p_1 x_1 + p_2 x_2 + \dots + p_n x_n$

Subject to  $w_1 x_1 + w_2 x_2 + \dots + w_n x_n \leq W$   
 $x_1, x_2, \dots, x_n \geq 0$  and integers

① Scout

$W$  = volume of knapsack

$p_j$  = value of one of item  $j$

$w_j$  = volume of item  $j$

② Greedy Rules

$W$  = weight limit

$p_j$  = profit of item  $j$

$w_j$  = weight of item  $j$

(ii)  $W$  = money limit

$p_j$  = original profit per unit of item  $j$

$w_j$  = cost per unit of item  $j$

Look at a dynamic programming formulation

$f(r, w)$  = maximum obtainable, replacing  $n$  by  $r$  and  $W$  by  $w$ .

$$= \max_{0 \leq x_r \leq \lfloor \frac{w}{w_r} \rfloor} [p_r x_r + f(r-1, w - w_r x_r)]$$

#Steps to compute  $f(n, W) = O(nW^2)$

Reduce time to  $O(nW)$

Instead of deciding how many  $x_r$ , decide whether  $x_r = 0$  or  $\geq 1$

$$f(r, w) = \max \begin{cases} f(r-1, w) & : x_r = 0 \\ f(r, w - w_r) + p_r & : x_r \geq 1 \end{cases}$$

#Steps =  $O(nW)$

$f(r, w) = \text{maximum of items, up to } w$   
 Why  $w$ .

$$= \max_{0 \leq x_r \leq \lfloor \frac{w}{w_r} \rfloor} [p_r x_r + f(r-1, w - w_r x_r)]$$

#steps to compute  $f(n, W) = O(nW^2)$

Instead of deciding how big  $x_r$ ,  
 decide whether  $x_r = 0$  or  $\geq 1$

$$f(r, w) = \max \begin{cases} f(r-1, w) & : x_r = 0 \\ f(r, w - w_r) + p_r & : x_r \geq 1 \end{cases}$$

#steps =  $O(nW)$

Maximize  $x_1 + 3x_2 + 6x_3 + 13x_4$

st  $w_1 + 2w_2 + 3w_3 + 4w_4 \leq 10$

Optimization  
 Subset

Assume that

w	r=1	r=2	r=3	r=4
0	0	0	0	0
1	1	0	0	0
2	2	3	0	0
3	3	6	0	0
4	4	9	6	0
5	5	12	9	13
6	6	15	12	16
7	7	18	15	19
8	8	21	18	22
9	9	24	21	25
10	10	27	24	28

Maximize  $x_1 + 3x_2 + 6x_3 + 13x_4$

Does not mean  
 exactly optimal  
 in this case