

10/14/15

Matrix A $m \times n$

$$S_A = \left\{ (p_1, p_2, \dots, p_m) : p_i \geq 0 \forall i \text{ \& } p_1 + \dots + p_m = 1 \right\}$$

$\left\{ \text{Randomized strategies for } A \right\}$

$$S_B = \left\{ (q_1, \dots, q_n) : q_j \geq 0 \forall j \text{ \& } q_1 + \dots + q_n = 1 \right\}$$

$$\text{Payoff: } \text{PAY}(p, q) = \sum_{i=1}^m \sum_{j=1}^n a_{ij} p_i q_j$$

In this situation $P_A = P_B$

$$P_A = \max_{P \in S_A} \min_{q \in S_B} \sum_{i=1}^m \sum_{j=1}^n a_{ij} P_i q_j$$

$$= \max_{P \in S_A} \left[\min_{q \in S_B} \sum_{j=1}^n \left(C_j(p) \right) q_j \right]$$

$$\sum_{i=1}^m C_{ij} P_i$$

Consider minimize $3q_1 + 4q_2 + 2q_3$
 s.t. $q_1 + q_2 + q_3 = 1$
 $q_1, q_2, q_3 \geq 0.$

$$P_A = \max_{P \in S_A} \min_{j=1 \dots m} C_j(p) \quad \text{--- Linear Program}$$

$$P_A = \max_{p \in S_A} \max_{\mu} \begin{aligned} &\mu \\ &\mu \leq c_1(p) \\ &\mu \leq c_2(p) \\ &\vdots \\ &\mu \leq c_n(p) \end{aligned}$$

$$P_A = \max_{\mu} \begin{aligned} &-\mu + \sum_{i=1}^m c_{i1} p_i \geq 0 \\ &-\mu + \sum_{i=1}^m c_{i2} p_i \geq 0 \\ &\vdots \\ &-\mu + \sum_{i=1}^m c_{in} p_i \geq 0 \end{aligned}$$

$$p_1 + p_2 + \dots + p_m = 1$$

$$p_1, p_2, \dots, p_m \geq 0.$$

$$P_B = \min_{q \in S_B} \max_{p \in S_A} \sum_{i=1}^m \sum_{j=1}^n a_{ij} p_i q_j$$

$$= \min_{q \in S_B} \left[\max_{p \in S_A} \sum_{i=1}^m \left(\sum_{j=1}^n a_{ij} q_j \right) p_i \right]$$

$d_j(q)$

$$= \min_{q \in S_B} \max \{ d_1(q), d_2(q), \dots, d_m(q) \}$$

min

η

$$\eta \geq d_1(q)$$

$$\eta \geq d_2(q)$$

\vdots

$$\eta \geq d_m(q)$$

LP

$$q_1 + \dots + q_n = 1$$

$$q_1, \dots, q_n \geq 0$$

max

$$\frac{P_A}{\sum}$$

$$-\sum + \sum_{i=1}^m a_{i1} P_i \geq 0$$

$$-\sum + \sum_{i=1}^m a_{in} P_i \geq 0$$

$$P_1 + \dots + P_m = 1$$

$$P_1, \dots, P_m \geq 0$$

min

$$\frac{P_B}{D}$$

$$-D + \sum_{j=1}^n a_{1j} q_j \leq 0$$

$$-D + \sum_{j=1}^n a_{mj} q_j \leq 0$$

$$q_1 + \dots + q_n = 1$$

$$q_1, \dots, q_n \geq 0$$

These two LP's are dual
to each other - check,

$$\Rightarrow P_A = P_B$$

Dominance

		B				
		4	3	6	5	4
A		3	2	5	4	3
		5	1	7	2	2

$1 < 2$ $3 < 2$
 ↑
 dominated by

2 is dominated by 1.

A plays 1 & B plays 2

Symmetric Game

A is $n \times n$ and $A^T = -A \Rightarrow \sum_{i=1}^n \sum_{j=1}^n a_{ij} p_i p_j = 0$

$\infty \times \infty$ matrix
 PAY (p, q)

0				
	0			
		0		
			0	
				0

$p_A \leq 0$
 $p_B \geq 0$

Both players play same

$a_{ii} = 0$
 $a_{ij} p_i p_j + a_{ji} p_j p_i$
 $p_A = p_B = 0$

Example

R, P, S

$$\begin{array}{c} R \\ P \\ S \end{array} \begin{array}{c} R \\ P \\ S \end{array} \begin{bmatrix} 0 & -1 & 1 \\ 1 & 0 & -1 \\ -1 & 1 & 0 \end{bmatrix}$$

Solution $p = (\frac{1}{3}, \frac{1}{3}, \frac{1}{3})$ — check LP^1_S
 $q = (\frac{1}{3}, \frac{1}{3}, \frac{1}{3})$

Non-Singular Games

A is non-singular

$$v = \frac{1}{\mathbf{1}^T A^{-1} \mathbf{1}} \quad \& \quad p = v A^{-T} \mathbf{1}$$
$$q = v A^T \mathbf{1}$$

if $p, q \geq 0$ then p, q solve the game.

$$P_A = v ?$$

Put $p = v A^{-T} \mathbf{1} \geq 0$

$$\mathbf{1} = \begin{bmatrix} 1 \\ \vdots \\ 1 \end{bmatrix}$$

$$\mathbf{1}^T p = v \mathbf{1}^T A^{-T} \mathbf{1} = \mathbf{1}$$

$$-s \mathbf{1} + \cancel{A}^T p \geq 0 \quad \leftarrow \text{Remaining Constraints}$$

? $s = v$

$$p = v A^{-T} \mathbf{1}$$

satisfies?

~~$$-v + v A^{-T} \mathbf{1} \geq 0$$~~

$$-v + v A^{-T} \mathbf{1} = 0$$