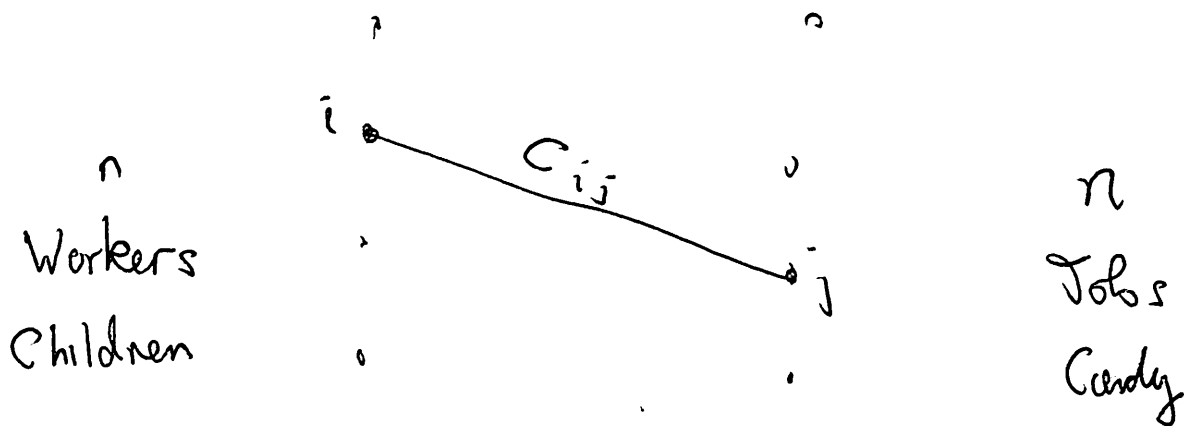


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Assignment Problem



C_{ij} = cost of worker doing job j

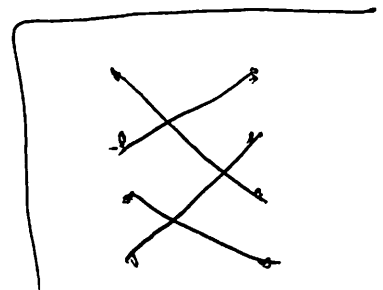
Problem: assign $i \rightarrow \pi(i)$ to

minimise
$$\sum_{i=1}^n C_{i, \pi(i)}$$

[In graph theory terms we see a minimum weight perfect matching]

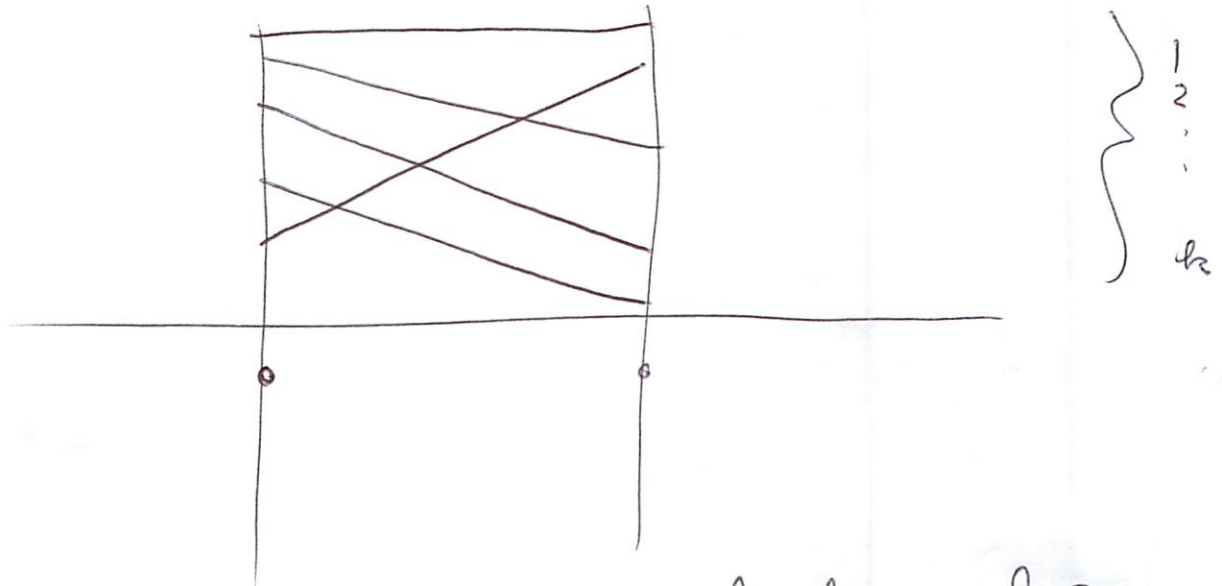
Reduce solution to a sequence

of shortest path problems



Suppose we have solved the

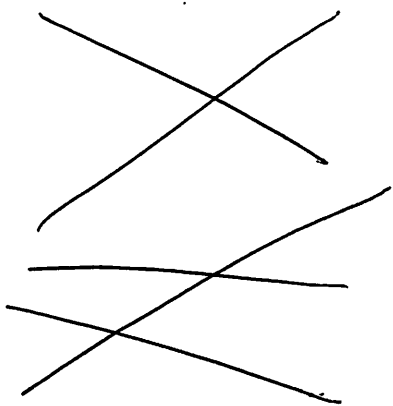
$k \times k$ problem:



Given solution to $k \times k$ problem,

solving $(k+1) \times (k+1)$ problem is a
shortest path problem.

Suppose solution to $k \times k$ problem is a matching M



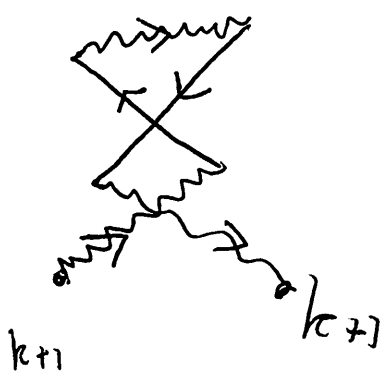
M' is any $(k+1)$ -matching

$k+1$ • • $k+1$

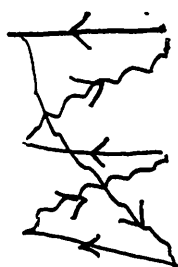
$$M' \oplus M = (M' / M) \cup (M / M')$$

Choose M' to maximize

$$\text{cost}(M') - \text{cost}(M) = \sum l(\text{zigzag}) - \sum l(\text{straight})$$



One of these



$l(C) \geq 0$

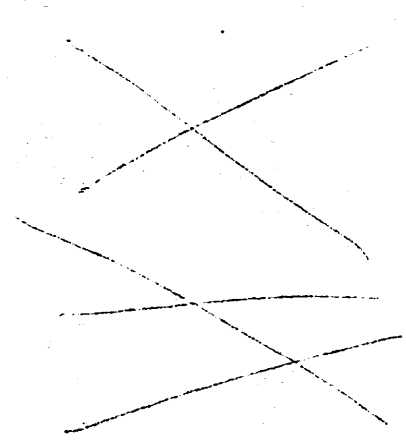
Several of these

Cost: $\frac{-c_{ij}}{+c_{ij}}$

Sum of lengths of paths + cycles =

If $l(C) < 0$ then it would give a better solution to $k \times k$ problem.

\mathbb{Z}^2 is a normal subgroup
 of M containing a and b



$(1+g)$ part of M
 normal

M

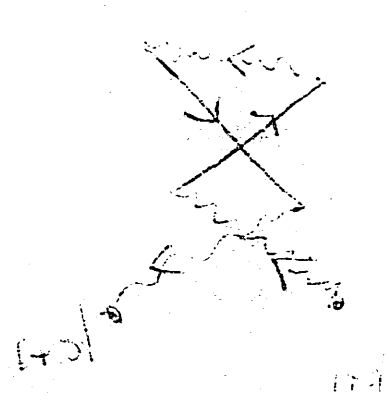
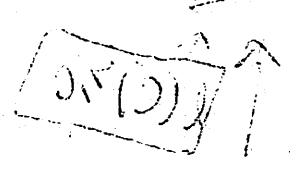
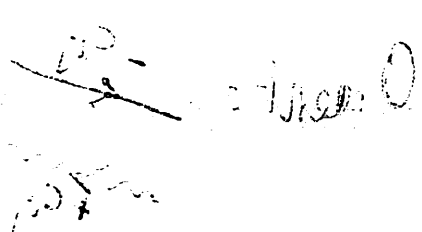
$(1+g)$

$(1+g)$

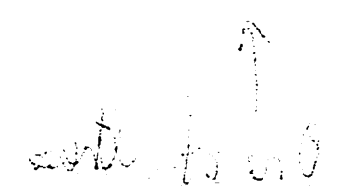
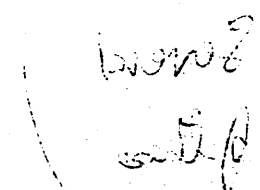
$$(M/M) \oplus (M/M) = M \oplus M$$

$$\rightarrow (M) \oplus (M) \oplus (M)$$

$$(N) \oplus (N) \oplus (N)$$



...



...

Integer Programming

Linear programming where you have to keep some variables integers. — Hard in general.

Assignment Problem:

Special

I.P.

$$\text{minimize } \sum_{i=1}^n \sum_{j=1}^n c_{ij} x_{ij}$$

$$\text{subject to } \sum_{i=1}^n x_{ij} = 1, \quad \forall j$$

$$\sum_{j=1}^n x_{ij} = 1, \quad \forall i$$

$$x_{ij} = 0 \text{ or } 1 \quad (*)$$

If \uparrow replace $(*)$

by $0 \leq x_{ij} \leq 1$

and use the

Simplex

algorithm then

$(*)$ holds in optimum