

Cell decomposition in valued fields

Yimu Yin
Department of Philosophy
Carnegie Mellon University

October 31, 2006

The topics of this talk are contained in the following three papers:

1. *The rationality of the Poincaré series associated to the p -adic points on a variety*, Jan Denef, Invent. Math., 77, pp. 1-23, 1984.
2. *Uniform p -adic cell decomposition and local zeta functions*, Johan Pas, J. reine angew. Math., 399, pp. 137-172, 1989.
3. *Rationality of p -adic Poincaré series: uniformity in p* , Angus MacIntyre, Annals of Pure and Applied Logic, 49, pp. 31-74, 1990.

Recall Macintyre's Theorem (generalized to finite extensions of \mathbb{Q}_p by Prestel and Roquette):

Theorem. 1. *The theory of p -adically closed fields of p -rank d admits quantifier elimination in Macintyre's language (with d many new constants).*

In 1984 Weispfenning gave a primitive recursive QE procedure for this theory, though it was done in a considerably expanded language.

The problem of uniformity: Develop a “natural” formalism in which the QE procedure is uniform for all

$\mathbb{Q}_2, \mathbb{Q}_3, \mathbb{Q}_5, \mathbb{Q}_7, \dots, \mathbb{Q}_{57}, \dots, \mathbb{Q}_{2^{232582657}-1}, \dots$

that is, independent of the choice of p .

Macintyre published a paper in 1990 that offered a solution to this problem. But the formalism there is very complicated. (Verdict: unnatural.)

Paul J. Cohen's Quantifier Elimination Procedure:

1. A primitive recursive decision procedure for \mathbb{Q}_p .
2. Elimination takes place inside a fixed henselian field, using only Hensel's Lemma.
3. He gives a procedure for "isolating" the roots of polynomials, and simultaneously reducing conditions on the n th root of a polynomial F to "simple" conditions on the coefficients of F .

Inspired by Cohen's work, Jan Denef subsequently developed the technique of p -adic cell decomposition. He used it to prove the following theorem:

Theorem. 2. *$P(T)$ is a rational function of T .*

Let $f_1(\bar{x}), \dots, f_r(\bar{x})$ be polynomials in m variables $\bar{x} = (x_1, \dots, x_m)$ with coefficients in \mathbb{Z}_p . For $n \in \mathbb{N}$ let $N^*(n)$ be the number of elements in the set

$$\{\bar{x} \pmod{p^n} : \bar{x} \in \mathbb{Z}_p^m \text{ and } \bigwedge_i (f_i(\bar{x}) = 0 \pmod{p^n})\},$$

and let $N(n)$ be the number of elements in the set

$$\{\bar{x} \pmod{p^n} : \bar{x} \in \mathbb{Z}_p^m \text{ and } \bigwedge_i (f_i(\bar{x}) = 0)\}.$$

Define the Poincaré series

$$P^*(T) = \sum_{n=0}^{\infty} N^*(n)T^n,$$

$$P(T) = \sum_{n=0}^{\infty} N(n)T^n.$$

Igusa proved that $P^*(T)$ is a rational function if $r = 1$. Meuser proved this for every r . Serre asked whether $P(T)$ is a rational function of T .

Let $|d\bar{x}|_p = |dx_1|_p \dots |dx_m|_p$ be the Harr measure on \mathbb{Q}_p^m with $|\mathbb{Z}_p^m|_p = 1$. Let

$$D = \{(\bar{x}, w) \in \mathbb{Z}_p^m \times \mathbb{Z}_p : \exists \bar{y} \in \mathbb{Z}_p^m (x = y \pmod{w} \wedge \bigwedge_i f_i(\bar{y}) = 0)\}.$$

For any positive real number s let

$$Z(s, p) = \int_D |w|_p^s |d\bar{x}|_p |dw|_p.$$

$$\begin{aligned}
Z(s, p) &= \sum_{n=0}^{\infty} \int_{\substack{D \\ \text{ord}_p(w)=n}} p^{-ns} |d\bar{x}|_p |dw|_p \\
&= \sum_{n=0}^{\infty} p^{-ns} \int_{\substack{(x, p^n) \in D \\ \text{ord}_p(w)=n}} |d\bar{x}|_p |dw|_p \\
&= \sum_{n=0}^{\infty} p^{-ns} \int_{(x, p^n) \in D} |d\bar{x}|_p \int_{\text{ord}_p(w)=n} |dw|_p \\
&= \sum_{n=0}^{\infty} p^{-ns} \frac{N(n)}{p^{nm}} \left(\frac{1}{p^n} - \frac{1}{p^{n+1}} \right) \\
&= \frac{p-1}{p} \sum_{n=0}^{\infty} N(n) (p^{-s} p^{-m-1})^n.
\end{aligned}$$

So to prove the rationality of $P(T)$, it is enough to prove that integrals of the form

$$Z(s, p) = \int_D |h(\bar{x})|_p^s |d\bar{x}|_p$$

is a rational function of p^{-s} , where h is a function and D a subset of \mathbb{Q}_p^n for some n , both of which are definable in a suitable language for \mathbb{Q}_p . (If h has no zero on \mathbb{Q}_p^m then s could be any real, otherwise we have to require $s > 0$.)

A suitable language (a.k.a. the Denef-Pas language)

We have 3 sorts: a field K , a residue field \overline{K} , and a valuation group $\Gamma \cup \{\infty\}$.

- a valuation $v : K \longrightarrow \Gamma \cup \{\infty\}$.
- an angular component $\overline{ac} : K \longrightarrow \overline{K}$.

Later we shall add more symbols to the language. For example, we want the Γ -sort to have the Presburger language at certain point.

What is an angular component?

1. $\overline{\text{ac}}(x) = 0$ iff $x = 0$;
2. $\overline{\text{ac}} : K^\times \longrightarrow \overline{K}^\times$ is a group homomorphism;
3. $\overline{\text{ac}} u = u + M$ for $u \in O \setminus M$.

Lemma. 3. For $a \neq b \in K$ with $v(a) = v(b)$, $\overline{\text{ac}} a = \overline{\text{ac}} b$ iff $v(a - b) > v(a) = v(b)$.

Axioms

1. $\text{char } K = 0$;
2. $\text{char } \overline{K} = 0$;
3. K is henselian;
4. v is a valuation, \overline{ac} is an angular component, and all other symbols are axiomatized in the standard way.

Call this theory $VF_{\overline{ac}}$, valued fields with angular component.

Definition. 4. A formula φ is simple if φ does not contain any K -quantifiers. A subset D of K^m or $K^m \times \overline{K}^n$ is simple if it is defined by a simple formula.

Definition. 5. A function $h : K^m \times \overline{K}^n \longrightarrow K$ is strongly definable if, for each simple formula $\varphi(t)$, there is a simple formula $\psi(x, \xi)$ such that

$$\varphi(h(x, \xi)) \leftrightarrow \psi(x, \xi).$$

What's a cell?

Let $(x, \xi) \in K^m \times \overline{K}^n$. Let C be a simple subset of $K^m \times \overline{K}^n$, which we shall call a parameter set. Let $b_1, b_2, c : C \rightarrow K$ be strongly definable functions. Let $\lambda \in \mathbb{N}$. Let \square_1, \square_2 be $<, \leq$ or no condition.

Definition. 6. For each $\xi \in \overline{K}^n$, the set

$$A(\xi) = \{(x, t) \in K^m \times K : (x, \xi) \in C, \\ v b_1(x, \xi) \square_1 \lambda v(t - c(x, \xi)) \square_2 v b_2(x, \xi), \\ \overline{ac}(t - c(x, \xi)) = \xi_1\}$$

is called a fiber.

Definition. 7. Suppose that if $\xi \neq \xi'$ then $A(\xi) \cap A(\xi') = \emptyset$.

$$A = \bigcup_{\xi \in \overline{K}^n} A(\xi)$$

is called a cell in $K^m \times K$ with center $c(x, \xi)$.

Cell decomposition

Let $f(x, t)$ be a polynomial of the form

$$g_d(x, \Delta)t^d + \dots + g_0(x, \Delta),$$

where $g_0(x, \Delta), \dots, g_d(x, \Delta)$ are strongly definable functions (Δ are extra parameters which will be omitted in the sequel).

Theorem. 8. *There is a finite partition of $K^m \times K$ into cells A such that:*

Write

$$f(x, t) = \sum_{i=0}^d a_i(x, \xi)(t - c(x, \xi))^i.$$

Let $A(\xi)$ be a fiber of A . Then for each $(x, \xi) \in A(\xi)$ we have

$$\begin{aligned} v f(x, t) &= v a_{i_0}(x, \xi)(t - c(x, \xi))^{i_0} \\ &= \min_{0 \leq i \leq d} v a_i(x, \xi)(t - c(x, \xi))^i \end{aligned}$$

and

$$\overline{ac} f(x, t) = \xi_{j_0},$$

where i_0, j_0 are fixed (i.e. do not depend on (x, ξ, t)).

For polynomials f_1, \dots, f_r :

Theorem. 9. *There is a finite partition of $K^m \times K$ into cells A such that:*

Let $A(\xi)$ be a fiber of A . Then for each $(x, t) \in A(\xi)$ and each $1 \leq i \leq r$ we have

$$v f_i(x, t) = v h_i(x, \xi)(t - c(x, \xi))^{\nu_i}$$

and

$$\overline{\text{ac}} f_i(x, t) = \xi_{\mu(i)},$$

where the h_i 's are strongly definable functions and $\nu_i \in \mathbb{N}$ and $1 \leq \mu(i) \leq n$ are fixed (i.e. do not depend on (x, ξ, t)).

Quantifier elimination

Theorem. 10. *The theory $VF_{\overline{ac}}$ admits elimination of K -quantifiers.*

Sketch of the proof. Need to consider formulas of the forms

$$\bigwedge_{i=1}^r \overline{ac} f_i(x, t) = \rho_i \wedge \bigwedge_{j=1}^s v g_j(x, t) = l_j.$$

After cell decomposition this is reduced to

$$(x, t) \in A(\xi) \wedge \bigwedge_{i=1}^r \xi_{\mu(i)} = \rho_i \\ \wedge \bigwedge_{j=1}^s (v h_j(x, \xi) + \nu_j v (t - c(x, \xi)) = l_j).$$

Introducing a new variable l for $v(t - c(x, \xi))$ we then add two conjuncts to the above

$$v(t - c(x, \xi)) = l \wedge \overline{ac}(t - c(x, \xi)) = \xi_1,$$

hence reduce the whole thing to

$$\exists t (v(t - c(x, \xi)) = l \wedge \overline{ac}(t - c(x, \xi)) = \xi_1),$$

which is true. □

Compute zeta functions

We now work with \mathbb{Q}_p . We expand the language so that the Γ -sort becomes a model of Presburger arithmetic.

Theorem. 11. *Let $D \subseteq \mathbb{Q}_p^m \times \mathbb{Q}_p$ be a simple subset defined by $\varphi(x, t)$. Consider*

$$Z(0, p) = \int_D |dx|_p |dt|_p.$$

Then there are simple formulas $\varphi_i(1 \leq i \leq s)$ such that for almost all p

$$Z(0, p) = \frac{1}{p} \sum_{i=1}^s \sum_{\xi_i \in \overline{\mathbb{Q}_p}^{n_i}} \sum_{l \in \mathbb{Z}} p^{-l} \int_{E_i(\xi_i, l)} |dx|_p,$$

where $E_i(\xi_i, l) \subseteq \mathbb{Q}_p^m$ is the set defined by $\varphi_i(x, \xi_i, l)$.

Lemma. 12. *Let $E \subseteq \mathbb{Z}^{m+1}$ be defined by a formula $\psi(l_1, \dots, l_m, n)$ that contains only Γ -variables. Suppose*

$$J(s) = \sum_E p^{-ns - l_1 - \dots - l_m}$$

is convergent for $s \in S$, with S an open subset of \mathbb{R} . Then there are polynomials $Q, R \in \mathbb{Z}[X, Y]$ such that for almost all P and all $s \in S$

$$J(s) = \frac{Q(p, p^{-s})}{R(p, p^{-s})}.$$

Lemma. 13 (Meuser's Lemma). *Let $L \subseteq \mathbb{Z}^m$ be defined by a finite system of linear inequalities in (k_1, \dots, k_m) with coefficients from \mathbb{Z} . Let $A_1(X), \dots, A_m(X) \in \mathbb{Z}[X]$ be linear. Suppose that*

$$J(s) = \sum_L p^{-\sum_{i=1}^m k_i A_i(s)}$$

is convergent for $s \in S$, with S an open subset of \mathbb{R} . Then $J(s)$ is a rational function of p^{-s} on S .

Theorem. 14. *Let $h(x)$ be a definable function and $D \subseteq \mathbb{Q}_p^m \times \mathbb{Q}_p$ a definable subset. Then there is a rational function $Q(T)/R(T)$ such that for almost all p*

$$Z(s, p) = \int_D |h(\bar{x})|_p^s |d\bar{x}|_p = \frac{Q(p^{-s})}{R(p^{-s})}.$$

By Denef's theorem which deals with each p separately, we conclude that $Z(s, p)$ is a rational function for all p and the degrees of numerators and denominators of these rational functions are bounded.