

## SURVIVAL TIME OF A RANDOM GRAPH

A. M. FRIEZE

*Received September 3, 1987*

Let  $V_n = \{1, 2, \dots, n\}$  and  $e_1, e_2, \dots, e_N, N = \binom{n}{2}$  be a random permutation of  $V_n^{(2)}$ . Let  $E_t = \{e_1, e_2, \dots, e_t\}$  and  $G_t = (V_n, E_t)$ . If  $\Pi$  is a monotone graph property then the hitting time  $\tau(\Pi)$  for  $\Pi$  is defined by  $\tau = \tau(\Pi) = \min \{t: G_t \in \Pi\}$ . Suppose now that  $G_t$  starts to deteriorate i.e. loses edges in order of age,  $e_1, e_2, \dots$ . We introduce the idea of the *survival time*  $\tau' = \tau'(\Pi)$  defined by

$$\tau' = \max \{u: (V_n, \{e_u, e_{u+1}, \dots, e_T\}) \in \Pi\}.$$

We study in particular the case where  $\Pi$  is  $k$ -connectivity. We show that

$$1) \lim_{n \rightarrow \infty} \Pr(\tau' \cong an) = e^{-2a} \text{ for } a \in \mathbb{R}^+$$

$$2) \lim_{n \rightarrow \infty} \frac{1}{n} E(\tau') = \frac{1}{n}$$

i.e.  $\tau'/n$  is asymptotically negative exponentially distributed with mean  $\frac{1}{2}$ .

### 1. Introduction

Let  $V_n = \{1, 2, \dots, n\}$  and  $e_1, e_2, \dots, e_N, N = \binom{n}{2}$  be a random permutation of  $V_n^{(2)}$ , the edge of set of the complete graph  $K_n$ . If  $G_t = (V_n, E_t)$  where  $E_t = \{e_1, e_2, \dots, e_t\}$  then the Markov chain  $\mathcal{G} = (G_t)_{t=0}^N$  is the central object of study in the theory of random graphs. If  $\Pi$  is a monotone increasing graph property then one wishes to establish asymptotic properties of the distribution of the *hitting time*

$$\tau(\Pi) = \min \{t: G_t \in \Pi\}.$$

Erdős and Rényi [4], [5] showed this to be an interesting problem and hundreds of papers have been written on this subject since the early papers — see the recent encyclopaedic text of Bollobás [2] or the introductory text of Palmer [7]. Erdős and Rényi thought of  $\mathcal{G}$  as describing the *evolution* of some *living organism*. Suppose we pursue this analogy and think of  $\tau(\Pi)$  as the time when this organism is *fully grown*. Going the way of all flesh our graph will start to *deteriorate*, say by

losing edges. We will assume that older edges disappear first. We propose to study the progress of the graph

$$H_u = (V_n, \{e_u, e_{u+1}, \dots, e_\tau\})$$

for various properties  $\Pi$ . We define the *survival time*  $\tau(\Pi)$  for the process  $\mathcal{G}$  by

$$\tau'(\Pi) = \max \{u: H_u \in \Pi\}.$$

The property discussed in this paper is  $k$ -connectivity. One can imagine many other properties worthy of study and we list some open problems at the end of the paper.

Our main result can be expressed as follows: let  $k \geq 1$  be a fixed integer. Let  $D_k$  denote the property of having minimum degree at least  $k$  and let  $C_k$  denote the property of being  $k$ -connected

**Theorem**

*If  $\Pi$  is  $C_k$  or  $D_k$  then*

- (i)  $\lim_{n \rightarrow \infty} \Pr(\tau'(\Pi) \geq an) = e^{-2a}$  for  $a \in \mathbb{R}^+$
- (ii)  $\lim_{n \rightarrow \infty} \frac{1}{n} E(\tau'(\Pi)) = \frac{1}{2}$
- (iii)  $\lim_{n \rightarrow \infty} \Pr(\tau'(C_k) = \tau'(D_k)) = 1. \blacksquare$

**2. Preliminaries**

We need some probability inequalities related to the  $k$ -connectivity of a random graph. The main thrust of the assertions are from Erdős and Rényi [6] but the calculations giving the precise bounds used are left to an appendix.

Let  $Z_m^{(r)}$  be the set of vertices of degree  $r$  in  $G_m$  and let  $z_m^{(r)} = |Z_m^{(r)}|$ .

Let  $S \subseteq V_n$  be a *non-trivial* separator of  $G_m$  if  $G_m[V_n - S]$  is not connected but has no isolated vertices.

**Lemma 1**

*Let*

$$(2.1) \quad m = \frac{1}{2} n \log n + \frac{1}{2} (k-1) n \log \log n - \frac{1}{2} \omega n \quad \text{be integer.}$$

(a) *Suppose that in (2.1)  $\omega = \omega(n) \rightarrow \infty$  but  $\omega = o(\log \log n)$ . In what follows  $\alpha = \alpha(k)$  is some 'constant' (depending only on the constant  $k$ ) whose exact value is unimportant.*

$$(2.2a) \quad \Pr(\delta(G_m) < k-1) \leq \frac{\alpha e^\omega}{\log n}$$

*where  $\delta(G)$  denotes the minimum degree of graph  $G$ .*

$$(2.2b) \quad \Pr(\delta(G_m) \geq k) \leq \alpha e^{-\omega}$$

$$(2.2c) \quad \Pr\left(\left|z_m^{(k-1)} - \frac{e^\omega}{(k-1)!}\right| \geq \frac{\varepsilon e^\omega}{(k-1)!}\right) \leq \frac{\alpha e^{-\omega}}{\varepsilon^2} \quad \text{for } 0 < \varepsilon < 1$$

(2.2d)  $\Pr(G_m \text{ has a non-trivial separator of size } \leq k-1) \leq \alpha \frac{(\log n)^2}{\sqrt{n}}$

(2.2e)  $\Pr(\exists \omega: |\omega| \leq \log \log n, m(\text{as in (2.1)}) \text{ such that } G_m \text{ has a non-trivial separator } S \text{ of size } \leq k-1 \text{ for which } G_m[V_n - S] \text{ has a component of size } t, 3 \leq t \leq \frac{1}{2}|V_n - S|) = O\left(\frac{(\log n)^2}{\sqrt{n}}\right).$

(2.2f)  $\Pr(\exists v, w \in Z_m^{(k-1)}: \text{the distance from } v \text{ to } w \text{ in } G_m \text{ is 2 or less}) \leq \alpha \sqrt{\frac{\log n}{n}}.$

(b) If we relax the restriction  $\omega = o(\log \log n)$  in (a) to  $\omega n \leq m$  then we can still prove

(2.3)  $\Pr(\delta(G_m) \geq k) \leq \alpha e^{-\omega/2}.$

(c) Let now  $m = \frac{1}{2}n \log n + \frac{1}{2}(k-1)n \log \log n + \frac{1}{2}\omega n$  where  $\omega = \omega(n) \rightarrow \infty$ . Then

(2.4)  $\Pr(\delta(G_m) < k) \leq \max\{\alpha e^{-\omega}, n^{-2}\}.$  ■

The hitting time of  $C_k$  was discussed by Bollobás and Thomason [3]. That paper established

(2.5)  $\lim_{n \rightarrow \infty} \Pr(\tau(C_k) = \tau(D_k)) = 1.$

Let  $m_1 = \left[\frac{1}{2}n \log n + \frac{1}{2}(k-1)n \log \log n - \frac{1}{2}n \log \log \log n\right]$  and  $m^* = \tau(D_k)$ . It will be useful to think of  $G_m^*$  in the following terms (Bollobás [1]):

$$E_{m^*} = E_{m_1} \cup X \cup Y$$

where

$$X = \{e \in E_{m^*} - E_{m_1} : e \cap Z_{m_1}^{(k-1)} = \emptyset\}$$

and

$$Y = E_{m^*} - (E_{m_1} \cup X).$$

This is valid, conditional on an event of probability  $1 - O((\log \log n)^{-1})$  — see (2.2b) with  $m = m_1$  and  $\omega = \log \log \log n$ .

Now, in this case,

(2.6)  $X$  is a random  $|X|$ -subset of  $(V_n - Z_{m_1}^{(k-1)})^{(2)}$ .

Also

(2.7)  $\Pr(|X \cup Y| > n \log \log \log n) \leq \Pr(\delta(G_m) \leq k-1) \leq \frac{\alpha}{\log \log n}.$

where  $m_2 = m_1 + \lceil n \log \log \log n \rceil$ , by (2.4). Furthermore

(2.8)  $\Pr(\exists e \in Y: e \subset Z_{m_1}^{(k-1)}) \leq \frac{\log n}{n}$

and

(2.9)  $\Pr(\exists z \in Z_{m_1}^{(k-1)}: |\{e \in Y: z \in e\}| > 6 \log \log \log n) \leq \frac{1}{\log \log n}.$

Given  $|X \cup Y| \leq n \log \log \log n$  and  $|Z_{m_1}^{(k-1)}| = O(\log \log n)$  (use  $\varepsilon=1$  in (2.2c)) it is unlikely that the conditions in (2.8), (2.9) will be violated — we are after all adding at most  $n \log \log \log n$  random edges.

### 3. Proof of the Theorem

(i) and (iii).

Let  $a \in \mathbb{R}^+$  and  $u = [an]$ . We show first that

$$(3.1) \quad \lim_{n \rightarrow \infty} \Pr(\delta(H_u) = k) = e^{-2a}.$$

We aim in fact to prove

$$(3.2a) \quad |\Pr(\delta(H_u) = k) - e^{-2a}| = O\left(\frac{(\log \log \log n)^2}{\log \log n}\right)$$

where the hidden constant in the “big  $O$ ” notation may depend on  $k$ . Thus from now on, if after describing an event  $\mathcal{E}$  we write  $[P=1-o(1)]$ , we mean  $\Pr(\mathcal{E}) = 1 - O((\log \log \log n)^2 / \log \log n)$ .

For the proof of (i) and (iii) we only require that  $a$  be a constant. However to prove (ii) we need to allow  $a$  to grow with  $n$ . In what follows we will assume only that

$$(3.2b) \quad a \leq \log \log \log \log n.$$

Let now  $Z^{(k-1)} = Z_{m_1}^{(k-1)}$  and  $\hat{Z}^{(k-1)}$  be the set of vertices of degree  $k-1$  in the graph  $H' = (V_n, \{e_u, e_{u+1}, \dots, e_{m_1}\})$ . Note that  $H'$  has the same distribution as  $G_{m_1-u+1}$  and we may use Lemma 1 (a) with  $\omega = \log \log \log n + 2a + O(1/n)$ . (The  $O(1/n)$  term accounts for  $\omega_n$  integral.)

$G_{m_1}$  is obtained from  $H'$  by adding  $u-1$  random edges and so  $Z^{(k-1)} \subseteq \hat{Z}^{(k-1)}$ . By applying (2.2c) twice: once with  $\omega = \log \log \log n$  and  $\varepsilon = \omega^{-1}$  and once with  $\omega = \log \log \log n + 2a$ ,  $\varepsilon$  as before, we obtain

$$(3.3) \quad \Pr\left(\frac{1-\varepsilon}{1+\varepsilon} \leq e^{2\alpha} \frac{|Z^{(k-1)}|}{|\hat{Z}^{(k-1)}|} \leq \frac{1+\varepsilon}{1-\varepsilon}\right) \geq 1 - \frac{2\alpha(\log \log \log n)^2}{\log \log n}.$$

Now let  $G'_m = (V_n, E'_m)$  where  $E'_m = E_m - E_{u-1}$  and note that  $G'_m$  has the same distribution as  $G_{m-u+1}$ . Let  $\hat{m} = \min\{m: \delta(G'_m) \geq k\}$ , and let  $\hat{v}$  be the unique  $[P=1-o(1)]$ , see (2.8)) vertex of degree  $k-1$  in  $G'_{\hat{m}-1}$ . Then

$$\delta(H_u) = k \leftrightarrow \hat{v} \in Z^{(k-1)}.$$

Now  $\hat{v}$  is “close to” being a random element of  $\hat{Z}^{(k-1)}$  and so we can see from (3.3) that  $\Pr(\hat{v} \in Z^{(k-1)}) \approx e^{-2a}$ , but let us do this more carefully.

Consider now the set  $\{e_{m_1}, e_{m_1+1}, \dots, e_N\}$  and in particular the subset  $F$  of edges incident with one vertex in  $\hat{Z}^{(k-1)}$  and one vertex not in  $\hat{Z}^{(k-1)}$ . We know  $[P=1-o(1)]$ , see (2.8)) that  $e_{\hat{m}} \in F$ . For  $v \in \hat{Z}^{(k-1)}$  let  $d_v$  be the degree of  $v$  in  $G_{m_1}$ . We work on the assumption that  $\hat{Z}^{(k-1)}$  is an independent set in  $G_{m_1}$ ,  $[P=1-o(1)]$ . Now  $d_v = k-1$  for  $v \in Z^{(k-1)}$  and  $d_v \leq 6 \log \log n$  otherwise  $[P=1-o(1)]$ , see (2.9)). If  $d_v$  were the same for all  $v \in \hat{Z}^{(k-1)}$  then we could deduce that  $\Pr(\hat{v} \in Z^{(k-1)}) = |Z^{(k-1)}| / |\hat{Z}^{(k-1)}|$  and we would be done. This is nearly so and for each  $v \in \hat{Z}^{(k-1)}$

we randomly select a set  $F_v \subseteq F$  of  $n_1 = [n - 1 - 6 \log \log n]$  edges incident with  $v$ . Let  $F' = \bigcup_{v \in Z^{(k-1)}} F_v$ . Then

$$\Pr(\hat{v} \in Z^{(k-1)} | e_{\hat{m}} \in F') = \frac{|Z^{(k-1)}|}{|\hat{Z}^{(k-1)}|}$$

and

$$\Pr(e_{\hat{m}} \notin F') = \frac{1}{n_1} + O\left(\frac{\log n}{n}\right).$$

(The  $O\left(\frac{\log n}{n}\right)$  term above, accounts for  $e_{\hat{m}} \in \hat{Z}^{(k-1)}$ ). This completes the proof of (3.2a).

We show next that

$$(3.4) \quad \lim_{n \rightarrow \infty} \Pr(\tau'(D_k) = \tau'(C_k)) = 1.$$

Observe that if  $\tau'(C_k) < u < \tau'(D_k)$  then either  $\tau(D_k) \neq \tau(C_k)$  or there exists  $S \subseteq V_n$ ,  $|S| = k - 1$  such that

$$(3.5) \quad S \text{ is a non-trivial separator of } H_{u-1}.$$

Note next that (3.2) implies

$$\lim_{n \rightarrow \infty} \Pr(\tau'(D_k) \geq n \log \log \log \log n) = 0.$$

Thus if  $u_1 = [n \log \log \log \log n]$ ,  $\tau'(C_k) < u < u_1$ ,  $K = (V_n, e_{u_1}, e_{u_1+1}, \dots, e_{m_1})$  and  $S$  is as in (3.5), then either

(a)  $S$  is a non-trivial separator of  $H_{u-1}$  and  $H_{u-1}[V_n - S]$  has a component of size  $t$ ,  $3 \leq t \leq \frac{1}{2} |V_n - S|$ ,

or

(b)  $S$  is a non-trivial separator of  $K$

or

(c)  $\delta(K) \leq k - 2$ ,

or

(d) 2 vertices of  $K$  of degree  $k - 1$  share a common neighbour.

But  $K$  has the same distribution as  $G_{m_1 - u_1 + 1}$  and (2.2) shows that these 4 events all have probability tending to zero.

(i) and (iii) follow directly from (3.1) and (3.4).

(ii).

We use

$$E\left(\frac{1}{n} \tau'(II)\right) = \int_0^{(1/2)n} \Pr(\tau'(II) \geq nx) dx$$

and

$$\tau'(D_k) \geq \tau'(C_k) \text{ whenever } \tau(D_k) = \tau(C_k).$$

Let  $\lambda = \log \log \log \log n$  and  $\mathcal{E}$  denote the event " $\tau(D_k) = \tau(C_k)$ ". Then

$$(3.6) \quad \Pr(\bar{\mathcal{E}}) = O\left(\frac{(\log n)^2}{\sqrt{n}}\right).$$

This follows from (2.2d) and the proof of Theorem VII. 4 of [2]. Now

$$(3.7) \quad \begin{aligned} E\left(\frac{1}{n} \tau'(C_k)\right) &\cong \int_0^\lambda \Pr(\tau'(C_k) \cong nx) dx \cong \\ &\cong \int_0^\lambda e^{-2x} dx - O\left(\frac{\lambda(\log \log \log n)^2}{\log \log n}\right) = \text{by (3.2)} \\ &= \frac{1}{2} - o(1). \end{aligned}$$

Now a lower bound for  $E(\tau'(D_k))$ .

$$\begin{aligned} E\left(\frac{1}{n} \tau'(D_k)\right) &= E\left(\frac{1}{n} \tau'(D_k) | \mathcal{E}\right) \Pr(\mathcal{E}) \cong \\ &\cong E\left(\frac{1}{n} \tau'(C_k) | \mathcal{E}\right) \Pr(\mathcal{E}) = \\ &= E\left(\frac{1}{n} \tau'(C_k)\right) - E\left(\frac{1}{n} \tau'(C_k) | \bar{\mathcal{E}}\right) \Pr(\bar{\mathcal{E}}) \cong \\ &\cong \frac{1}{2} - o(1) - E\left(\frac{1}{n} \tau(C_k) | \bar{\mathcal{E}}\right) \Pr(\bar{\mathcal{E}}). \end{aligned}$$

But

$$E\left(\frac{1}{n} \tau(C_k) | \bar{\mathcal{E}}\right) \cong (\log n)^2 + n \Pr(\tau(C_k) \cong n(\log n)^2) / \Pr(\bar{\mathcal{E}}).$$

Thus, using (3.6) and  $\Pr(\tau(C_k) \cong n(\log n)^2) = o(1/n)$  (very crudely), we have

$$(3.8) \quad E\left(\frac{1}{n} \tau(C_k) | \bar{\mathcal{E}}\right) \Pr(\bar{\mathcal{E}}) = o(1)$$

and so

$$E\left(\frac{1}{n} \tau'(D_k)\right) \cong \frac{1}{2} - o(1).$$

Now for upper bounds:

$$(3.9) \quad \begin{aligned} E\left(\frac{1}{n} \tau'(D_k)\right) &= \int_0^\lambda \Pr(\tau'(D_k) \cong nx) dx + \int_\lambda^{(1/2)n} \Pr(\tau'(D_k) \cong nx) dx \cong \\ &\cong \int_0^\lambda e^{-2x} dx + O\left(\frac{\lambda(\log \log \log n)^2}{\log \log n}\right) + \int_\lambda^\infty 2\alpha e^{-x/2} dx + \int_\lambda^{(1/2)n} n^{-2} dx, \\ &= \frac{1}{2} + o(1). \end{aligned}$$

The first integral in (3.9) is approximated as in (3.7). For the second note that

$$(3.10) \quad \Pr(\tau'(D_k) \cong nx) \cong \Pr(\delta(G_{m^+}) < k) + \Pr(\delta(G_{m^-}) \cong k)$$

where

$$m^+ = \frac{1}{2} n \log n + \frac{1}{2} (k-1) n \log \log n + \frac{1}{2} nx$$

and

$$m^- = m^+ - nx.$$

Now use (2.3) with  $m=m^-$  and (2.4) with  $m=m^+$  in (3.10). Now an upper bound for  $E(\tau'(c_k))$ .

$$\begin{aligned} E\left(\frac{1}{n} \tau'(C_k)\right) &= E\left(\frac{1}{n} \tau'(C_k)|\mathcal{E}\right) \Pr(\mathcal{E}) + E\left(\frac{1}{n} \tau'(C_k)|\bar{\mathcal{E}}\right) \Pr(\bar{\mathcal{E}}) \cong \\ &\cong E\left(\frac{1}{n} \tau'(D_k)|\mathcal{E}\right) \Pr(\mathcal{E}) + E\left(\frac{1}{n} \tau(C_k)|\bar{\mathcal{E}}\right) \Pr(\bar{\mathcal{E}}) \cong \\ &\cong E\left(\frac{1}{n} \tau'(D_k)\right) + o(1) \cong \quad (\text{by (3.8)}) \\ &\cong \frac{1}{2} + o(1) \quad (\text{by (3.10)}). \end{aligned}$$

This completes the proof of our theorem.

#### 4. Revival Time

We gain some insight into the distribution of survival time by considering the *revival time*. Let  $u=\tau'(\Pi)$  and consider adding random edges to  $H_u$  until a graph with property  $\Pi$  is obtained. The number of edges added  $\tau''(\Pi)$  is called the revival time. It is straightforward to show that we can replace  $\tau'$  by  $\tau''$  in our theorem and obtain a valid result.

For example: if  $\Pi=D_k$  then  $H_u[P=1-o(1)]$  contains a unique vertex  $v$  of minimal degree  $k-1$ .  $\tau'' \cong an$  if  $v$  is not incident with any of the first  $an$  random edges added to  $H_u$ . The probability of this is approximately  $\left(1-\frac{2}{n}\right)^{an} \approx e^{-2a}$ . For  $\Pi=C_k$  we use the smaller likelihood of non-trivial separators.

It seems now that in general one may be able to guess the result for survival time by computing the revival time, which seems easier. One then needs a few asymptotic calculations for verification.

## 5. Open Questions

What is the survival time for the following properties:

- (1) having a cycle?
- (2) having a cycle of size  $k$ ?
- (3) having a path of length  $k$ ?
- (4) having a tree component of size  $k$ ?
- (5) having a clique of size  $k$ ?
- (6) having any fixed subgraph?
- (7) being non-planar?
- (8) having diameter  $k$ ?
- (9) having a vertex of degree  $k$ ?

One can also consider similar problems for random bipartite graphs, digraphs or subgraphs of the  $n$ -cube.

There are 2 conspicuous omissions from our list. These are perfect matchings and hamilton cycles. If  $H_k$  is the property of having  $\lfloor k/2 \rfloor$  edge disjoint hamilton cycles plus a further edge disjoint matching of size  $\lfloor n/2 \rfloor$ , if  $k$  is odd, then it seems fairly clear that we can add  $H_k$  to our theorem. The proof does not require any new ideas and would be rather long, too long for this paper.

## Appendix

As usual let  $G_p$ ,  $p = \frac{N}{m}$  denote the random graph in which edges are independently included with probability  $p$ . Let  $\Pi$  be any graph property. Then

$$(A0) \quad \Pr(G_p \in \Pi) = \sum_{m'} \Pr(G_{m'} \in \Pi) \Pr(G_p \text{ has } m' \text{ edges}).$$

Thus if  $\Pi$  is monotone i.e. if it is preserved either by adding edges or by deleting edges, then, for large  $n$

$$(A1) \quad \Pr(G_m \in \Pi) \leq 3 \Pr(G_p \in \Pi).$$

We can thus work mainly with  $G_p$  and multiply our estimates by 3. Our inequalities are only required to hold for  $n$  large.

**Proof of (2.2a).**

For  $k \geq 2$

$$\begin{aligned} \Pr(\delta(G_p) \leq k-2) &\leq n \sum_{t=0}^{k-2} \binom{n-1}{t} p^t (1-p)^{n-1-t} \leq \\ &\leq 2n \sum_{t=0}^{k-2} \frac{n^t}{t!} \left(\frac{\log n}{n}\right)^t \frac{(\log n)^{-(k-1)} e^\omega}{n} \leq \\ &\leq \frac{3e^\omega}{(k-2)! \log n}. \end{aligned}$$

Now use (A1).



**Proof of (2.2b).**

Let  $z_p^{(k-1)}$  be the number of vertices of degree  $k-1$  in  $G_p$ .

$$\begin{aligned} E(z_p^{(k-1)}) &= n \binom{n-1}{k-1} p^{k-1} (1-p)^{n-k} = \\ &= \left( 1 + O\left(\frac{(\log n)^2}{n}\right) \right) \left(\frac{np}{\log n}\right)^{k-1} \frac{e^\omega}{(k-1)!} = \\ &= \left( 1 + O\left(\frac{\log \log n}{\log n}\right) \right) \frac{e^\omega}{(k-1)!}. \end{aligned}$$

Preparing for the Chebyshev inequality

$$\begin{aligned} E(z_p^{(k-1)}(z_p^{(k-1)} - 1)) &= \\ &= n(n-1) \left( p \binom{n-2}{k-2} p^{k-2} (1-p)^{n-k} + (1-p) \binom{n-2}{k-1} p^{k-1} (1-p)^{n-k-1} \right)^2 = \\ &= n(n-1) \left( p \left( \frac{E(z_p^{(k-1)})}{n(n-1)p} \right)^2 + (1-p) \left( \frac{E(z_p^{(k-1)})}{n(1-p)} \frac{n-1}{n-k} \right)^2 \right) = \\ &\cong (1+2p)(E(z_p^{(k-1)}))^2 \end{aligned}$$

and so

$$\begin{aligned} \text{(A2)} \quad \text{Var}(z_p^{(k-1)}) &\cong E(z_p^{(k-1)}) + 2pE(z_p^{(k-1)})^2 \cong \\ &\cong 2E(z_p^{(k-1)}) \quad \text{as } \omega = o(\log \log n). \end{aligned}$$

Thus by the Chebycheff inequality

$$\Pr(z_p^{(k-1)} = 0) \cong \frac{2}{E(z_p^{(k-1)})} \cong 3(k-1)! e^{-\omega}.$$

Thus

$$\Pr(\delta(G_p) \cong k) \cong 3(k-1)! e^{-\omega}.$$

We can now use (A1).

**Proof of (2.2c).**

The property in question is not monotone (or even convex) [2] and so we cannot use (A1). We can however assert using (A2) that

$$\Pr\left(\left|z_p^{(k-1)} - \frac{e^\omega}{(k-1)!}\right| \cong \frac{\varepsilon e^\omega}{(k-1)!}\right) \cong \frac{2e^{-\omega}}{\varepsilon^2}.$$

It now follows from (A0) that there exists  $m', m - \sqrt{n \log n} \cong n' \cong n$ , such that

$$\Pr\left(\left|z_{m'}^{(k-1)} - \frac{e^\omega}{(k-1)!}\right| \cong \frac{\varepsilon e^\omega}{(k-1)!}\right) \cong \frac{2e^{-\omega}}{\varepsilon^2}.$$

As in the proof of (2.2a) we can deduce that

$$\Pr(\delta(G_{m'}) < k-1) \leq \frac{10e^\omega}{(k-2)! \log n}.$$

Since  $G_m$  is obtained from  $G_{m-m'}$  by adding  $m-m'$  random edges, we have that given  $\mathcal{E} = \{\delta(G_{m'}) \geq k-1 \text{ and } z_{m'}^{(k-1)} \text{ is 'close' to its mean}\}$ ,

$$\begin{aligned} \Pr(z_m^{(k-1)} \neq z_{m'}^{(k-1)} | \mathcal{E}) &= O\left(\frac{\sqrt{n \log n e^\omega}}{N-n}\right) = \\ &= O\left(\sqrt{\frac{\log n}{n}} e^\omega\right) \end{aligned}$$

and (2.2c) follows.

**Proof of (2.2d).**

Let  $\theta$  be the probability that  $G_p$  has a non-trivial separator of size  $s$ ,  $0 \leq s \leq k-1$ . Then

$$\theta \leq \sum_{s=0}^{k-1} \sum_{t=2}^{(1/2)(n-s)} \binom{n}{s} \binom{n-s}{t} t^{t-2} p^{t-1} (1-(1-p)^s) (1-p)^{t(n-s-t)}$$

(choose an  $s$ -set  $S$  for the separator, a  $t$ -set  $T$  for a small component, a spanning tree of  $T$ . Multiply by the probability that the edges of the tree exist and there is at least one  $v, T$  edge for each  $v \in S$  and no  $S, S \cup T$  edges.)

Thus

$$\begin{aligned} \theta &\leq \sum_{s=0}^{k-1} \sum_{t=2}^{(1/2)(n-s)} \binom{ne}{s} \binom{ne}{t} t^{t-2} p^{t-1} (tp)^s e^{-t(n-s-t)p} \leq \\ \text{(A3)} \quad &\leq 8 \sum_{s=0}^{k-1} \sum_{t=2}^{(1/2)(n-s)} \frac{1}{p} (npe^{-np} t^{s-2/t} e^{(s+t+1)p})^t = \\ &= O\left(\frac{\log n}{n}\right) \end{aligned}$$

(for small  $t$ , the "complex" term above is  $O\left(\left(\frac{e^\omega \log n}{n}\right)^t\right)$ . For larger  $t$ , it is  $O\left(\left(\frac{e^\omega \log n}{\sqrt{n}}\right)^t\right)$ . Unfortunately, we are not dealing with a monotone property. However (A0) implies

$$\Pr(G_m \text{ has a non-trivial separator}) \leq \theta \Pr(G_p \text{ has } m \text{ edges})^{-1}$$

and (2.2d) follows. (See [2] Theorem II.2.)

**Proof of (2.2e).**

We only have to consider the sum in (A3) for  $t \geq 3$ , which is then  $O\left(\left(\frac{\log n}{n}\right)^2\right)$ . We only have to multiply this by  $O\left(n \log \log n \times \frac{\sqrt{n}}{\sqrt{\log n}}\right)$  in order to account for the number of different values of  $m$  and the transition from  $G_p$  to  $G_m$ .

**Proof of (2.2f).**

We can clearly assume  $k \geq 2$ . Let  $\psi$  be the probability that there exist  $v, w \in Z_m^{(k-1)}$  at distance 2 or less from each other in  $G_p$ . Then

$$\psi \cong \sum_{r=2}^3 \frac{r!}{2} \binom{n}{r} p^{r-1} \left( \binom{n-r}{k-2} p^{k-2} (1-p)^{n-k-r+2} \right)^2 + 3 \binom{n}{3} p^3 \left( \binom{n-3}{k-2} p^{k-2} (1-p)^{n-k-1} \right)^2$$

(the first term above deals with paths of length  $r=2$  or 3 and the second term deals with triangles.) Thus  $\psi = O\left(\frac{1}{n}\right)$  and we can finish as in the proof of (2.2d).

**Proof of (2.3).**

The proof used for (2.2b) will be valid for  $p \geq p_0 = \frac{\log n}{4n}$ , say. All that is needed for smaller  $p$  is to show

$$\Pr(G_{p_0} \text{ has no isolated vertices}) = O\left(\frac{e^{np}}{n}\right).$$

This can be done using the Chebycheff inequality.

**Proof of (2.4).**

Non-trivial separators are handled as in the proof of (2.2d). The minimum degree calculation only requires the use of the expected number of vertices of degree  $k-1$  or less. ■

**References**

- [1] B. BOLLOBÁS, The evolution of sparse graphs, In *Graph Theory and Combinatorics*, Proc. Cambridge Combinatorial Conference in honour of Paul Erdős (B. Bollobás, Ed.), Academic Press (1984), 35—57.
- [2] B. BOLLOBÁS, *Random Graphs*, Academic Press, 1985.
- [3] B. BOLLOBÁS and A. THOMASON, Random graphs of small order, *Annals of Discrete Mathematics*.
- [4] P. ERDŐS and A. RÉNYI, On random Graphs I, *Publ. Math. Debrecen*, 6 (1959), 290—297.
- [5] P. ERDŐS and A. RÉNYI, On the evolution of random graphs, *Publ. Math. Inst. Hungar. Acad. Sci.*, 7 (1960), 17—61.
- [6] P. ERDŐS and A. RÉNYI, On the strength of connectedness of a random graph, *Acta Math. Acad. Sci. Hungar.*, 12 (1961), 261—267.
- [7] E. PALMER, *Graphical Evolution*.

A. M. Frieze

Department of Mathematics  
Carnegie Mellon University  
Pittsburgh, PA 15213