

9/11/13

$$C = \begin{bmatrix} 3 & 2 & 1 & 5 \\ 4 & 3 & 4 & 1 \\ 2 & 4 & 1 & 3 \\ 5 & 4 & 3 & 4 \end{bmatrix}$$

i	1	2	3	4
$\pi(i)$	4	3	2	1
y_i	10	8	8	10

$i=1$

$$c_{ij} + \alpha y_j$$

$$3 + \frac{1}{2} \cdot 10$$

$$2 + \frac{1}{2} \cdot 8$$

$$1 + \frac{1}{2} \cdot 8 \leftarrow 5 \quad *$$

$$5 + \frac{1}{2} \cdot 10$$

$$\underline{i=2}$$

$$4 + \frac{1}{2} \times 10$$

$$3 + \frac{1}{2} \times 8$$

$$4 + \frac{1}{2} \times 8$$

$$1 + \frac{1}{2} \times 10 = 6$$

New π

3 4 3 3

$$\underline{i=3}$$

$$2 + \frac{1}{2} \times 10$$

$$4 + \frac{1}{2} \times 8$$

$$1 + \frac{1}{2} \times 8 = 5$$

$$3 + \frac{1}{2} \times 10$$

$$C = \begin{bmatrix} 3 & 2 & 1 & 5 \\ 4 & 3 & 4 & 1 \\ 2 & 4 & 1 & 3 \\ 5 & 4 & 3 & 4 \end{bmatrix}$$

$$\underline{i=4}$$

$$5 + \frac{1}{2} \times 10$$

$$4 + \frac{1}{2} \times 8$$

$$3 + \frac{1}{2} \times 8 \quad \rangle$$

$$4 + \frac{1}{2} \times 10$$

Evaluate

$$y_1 = 1 + \frac{1}{2} y_3 = 2$$

$$y_2 = 1 + \frac{1}{2} y_4 = 2$$

$$y_3 = 1 + \frac{1}{2} y_3 = 2$$

$$y_4 = 3 + \frac{1}{2} y_3 = 4$$

$$C = \begin{bmatrix} 3 & 2 & 1 & 5 \\ 4 & 3 & 4 & 1 \\ 2 & 4 & 1 & 3 \\ 5 & 4 & 3 & 4 \end{bmatrix}$$

Then
Check for optimality

...

Optimality Conditions.

(1) Assume π satisfies optimality conditions.

$\hat{\pi}$ is any other policy.

$$s_i = \hat{y}_i - y_i$$

$$\hat{y}_i = c_{i, \hat{\pi}(i)} + \alpha \sum y_{\hat{\pi}(i)}$$

$$y_i \leq c_{i, \hat{\pi}(i)} + \alpha \sum y_{\hat{\pi}(i)}$$



$$s_i \geq \alpha \sum_{\hat{\pi}(i)} y_{\hat{\pi}(i)}$$

$$r_i \equiv \alpha \sum_{\hat{\pi}(i)} \equiv \alpha^2 \sum_{\hat{\pi}^2(i)} \equiv \dots \equiv \alpha^k \sum_{\hat{\pi}^k(i)} \rightarrow 0$$

$$\Rightarrow r_i \geq 0, \forall i$$

So, if optimality condition holds, we have best $y_i, \forall i$

(ii) Suppose optimality condition does not hold.

$$y_i > c_{ij} + \alpha y_j \quad \text{for some } i, j$$

So define $\hat{\pi}$ by

$$\hat{\pi}(i) = \underset{j}{\operatorname{arg\,min}} \quad c_{ij} + \alpha y_j \quad \& \quad \hat{\pi} \neq \pi$$

which minimizes \leftarrow

$$I = \{i : \pi(i) \neq \hat{\pi}(i)\}$$

To show (i) $\hat{y}_i \leq y_i$ & (ii) $i \in I \Rightarrow \hat{y}_i < y_i$

\Rightarrow no repetition of Π

& eventual termination with optimality condition.

$$\langle y_i, y_i \rangle = c_i \langle \hat{y}_i, \hat{y}_i \rangle + \alpha \langle y_i, \hat{y}_i \rangle$$

$$\langle y_i, y_i \rangle \leq c_i \langle \hat{y}_i, \hat{y}_i \rangle + \alpha \langle y_i, \hat{y}_i \rangle \quad \forall i \in I$$

$\langle y_i, y_i \rangle \leq \langle y_i, y_i \rangle$ then

$$\langle y_i, y_i \rangle \leq \alpha \langle \hat{y}_i, \hat{y}_i \rangle \leq \dots \leq \alpha^k \langle \hat{y}_i, \hat{y}_i \rangle \rightarrow 0 \quad \forall i \in I$$

Solution via Linear Programming.

We have to find y_1, y_2, \dots, y_n

s.t.

$$y_i = \min_j [C_{ij} + \alpha y_j]$$

then $\pi(i) = \operatorname{argmin}$

Maximize

$$y_1 + y_2 + \dots + y_n$$

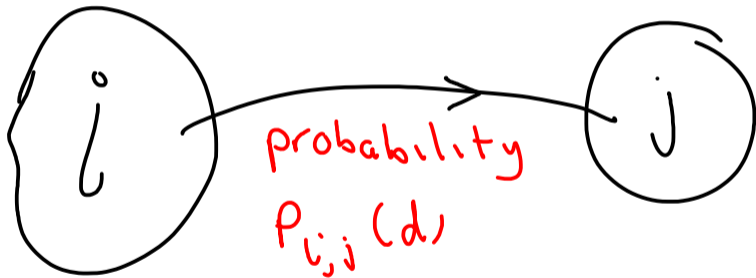
Subject to

$$y_i \leq C_{ij} + \alpha y_j$$

$\forall i, j$

Probabilistic Version

$D_i := \{ \text{possible decisions e.g. how much to produce} \}$



$C_{i:d} = E(\text{cost of move})$

Given d

$\pi: (i) \in D_i, \forall i$

Choose π to

minimize expected
NPV.